# A Quantitative Geometric Approach to Neural-Network Smoothness 

Zi Wang<br>University of Wisconsin-Madison<br>zw@cs.wisc.edu

Gautam Prakriya<br>The Chinese University of Hong Kong<br>gautamprakriya@gmail.com

Somesh Jha<br>University of Wisconsin-Madison<br>jha@cs.wisc.edu


#### Abstract

Fast and precise Lipschitz constant estimation of neural networks is an important task for deep learning. Researchers have recently found an intrinsic trade-off between the accuracy and smoothness of neural networks, so training a network with a loose Lipschitz constant estimation imposes a strong regularization, and can hurt the model accuracy significantly. In this work, we provide a unified theoretical framework, a quantitative geometric approach, to address the Lipschitz constant estimation. By adopting this framework, we can immediately obtain several theoretical results, including the computational hardness of Lipschitz constant estimation and its approximability. We implement the algorithms induced from this quantitative geometric approach, which are based on semidefinite programming (SDP) $\left.\right|^{1}$ Our empirical evaluation demonstrates that they are more scalable and precise than existing tools on Lipschitz constant estimation for $\ell_{\infty}$-perturbations. Furthermore, we also show their intricate relations with other recent SDP-based techniques, both theoretically and empirically. We believe that this unified quantitative geometric perspective can bring new insights and theoretical tools to the investigation of neural-network smoothness and robustness.


## 1 Introduction

The past decade has witnessed the unprecedented success of deep learning in many machine learning tasks (Krizhevsky et al., 2017, Mikolov et al., 2013). Despite the growing popularity of deep learning, researchers have also found that neural networks are very vulnerable to adversarial attacks (Szegedy) et al. 2014, Goodfellow et al. 2015; Papernot et al., 2016). As a result, it is important to train neural networks that are robust against those attacks (Madry et al., 2018). In recent years, the deep learning community starts to focus on certifiably robust neural networks (Albarghouthi, 2021, Hein and Andriushchenko, 2017, Katz et al., 2017; Cohen et al., 2019, Raghunathan et al., 2018; Wang et al., 2022; Leino et al., 2021). One way to achieve certified robustness is to estimate the smoothness of neural networks, where the smoothness is measured by the Lipschitz constant of the neural network. Recent works have found that to achieve both high accuracy and low Lipschitzness, the network has to significantly increase the model capacity (Bubeck and Sellke, 2021). This implies that there is an intrinsic tension between the accuracy and smoothness of neural networks.
Commonly considered adversarial attacks are the $\ell_{\infty}$ and $\ell_{2}$-perturbations in the input space. Leino et al. (2021); Cohen et al. (2019) have successfully trained networks with low $\ell_{2}$-Lipschitz constant,

[^0]and Huang et al. (2021) trained networks with low local Lipschitzness for $\ell_{\infty}$-perturbations. There are few well-established techniques to train neural networks with low global Lipschitzness for $\ell_{\infty^{-}}$ perturbations. The techniques for the $\ell_{2}$-perturbation do not easily transfer to the $\ell_{\infty}$-case (Leino et al., 2021). One critical step is to measure the Lipschitz constant more precisely and efficiently. Jordan and Dimakis (2020) showed that for ReLU networks, it is NP-hard to approximate the Lipschitz constant for $\ell_{\infty}$-perturbations within a constant factor, and proposed an exponential-time algorithm to compute the exact Lipschitz constant. However, researchers are interested in more scalable approaches to certify and train networks. In this work, we consider the Formal Global Lipschitz constant (FGL) (See Equation (3)), which is roughly the maximum of the gradient operator norm, assuming all activation patterns on hidden layers are independent and possible. FGL is an upper bound of the exact Lipschitz constant and has been used in Raghunathan et al. (2018); Fazlyab et al. (2019); Latorre et al. (2020).

We address the Lipschitz constant estimation from the quantitative geometric perspective. Quantitative geometry aims to understand geometric structures via quantitative parameters, and has connections with many mathematical disciplines, such as functional analysis and combinatorics (Naor, 2013). In computer science, quantitative geometry plays a central role in understanding the approximability of discrete optimization problems. We approach those hard discrete optimization problems by considering the efficiently solvable continuous counterparts, and analyze the precision loss due to relaxation, which is often the SDP relaxation (Goemans and Williamson, 1995; Nesterov, 1998; Alon and Naor, 2004). The natural SDP relaxations for the intractable problems usually induce the optimal known polynomial time algorithms (Bhattiprolu et al. 2022). By adopting the quantitative geometric approach, we can immediately understand the computational hardness and approximability of FGL estimations. Our algorithms on two-layer networks are the natural SDP relaxations from the quantitative geometric perspective.
Latorre et al. (2020) employed polynomial optimization methods on the FGL estimation for $\ell_{\infty^{-}}$ perturbations. Polynomial optimization is a very general framework, and many problems can be cast in this framework (Motzkin and Straus, 1965, Goemans and Williamson, 1995). Therefore, we argue that this is not a precise characterization of the FGL-estimation problem. On the other hand, there are also several SDP-based techniques for FGL estimations. Raghunathan et al. (2018) proposed an SDP algorithm to estimate the FGL for $\ell_{\infty}$-perturbations of two layer networks, and Fazlyab et al. (2019) devised an SDP algorithm to estimate the $\ell_{2}$-FGL. We will demonstrate the intricate relationships between our algorithms and these existing SDPs on two-layer networks.
Several empirical studies have found that techniques on one $\ell_{p}$-perturbations often do not transfer to another ones, even though the authors claim that in theory these techniques should transfer (Fazlyab et al. 2019, Leino et al. 2021). This in-theory claim usually comes from a qualitative perspective. In finite-dimensional space, one can always bound one $\ell_{p}$-norm from another one, so techniques for one $\ell_{p}$-perturbations can also provide another bound for a different $\ell_{p}$-perturbations. However, this bound is loose and in practice not useful (Latorre et al., 2020). Instead, we believe that when transferring techniques from one norm to another one, we should consider the quantitative geometric principle: we should separate the geometry-dependent component from the geometry-independent one in those techniques, and modify the geometry-dependent component accordingly for a different normed space. As we will demonstrate, our whole work is guided by this principle. We hope that our unified quantitative geometric framework can bring insights to the empirical hard-to-transfer observations, and new tools to address these issues.

Contributions. To summarize, we have made the following contributions:

1. We provide a unified theoretical framework to address FGL estimations, which immediately yields the computational hardness and SDP-based approximation algorithms on two-layer networks (Section 3).
2. We demonstrate the relations between our algorithms and other SDP-based tools, which in return inspires us to design the algorithms for multi-layer networks. This provides more insightful and compositional interpretations of existing works, and makes them easier-togeneralize (Sections 4 and 5).
3. We implement the algorithms and name the tool GeoLIP. We empirically validate our theoretical claims, and compare GeoLIP with existing methods to estimate FGL for $\ell_{\infty}$-perturbations. The result shows that GeoLIP provides a tighter bound ( $20 \%-60 \%$ improvements) than
existing tools on small networks, and much better results than the naive matrix-norm-product method for deep networks, which existing tools cannot handle (Section6).

## 2 Preliminaries

Notation. Let $[n]=\{1, \ldots, n\}$. For two functions $f$ and $g, f \circ g(x)=f(g(x))$. A 0-1 cube is $\{0,1\}^{n}$, and a norm- 1 cube is $\{-1,1\}^{n}$ for some integer $n>0 . \mathbb{R}_{+}=[0, \infty)$. For any vector $v \in \mathbb{R}^{n}, \operatorname{diag}(v)$ is an $n \times n$ diagonal matrix, with diagonal values $v$. Let $e_{n}=(1, \ldots, 1) \in \mathbb{R}^{n}$ be an $n$-dimensional vector of all 1 's; and $I_{n}=\operatorname{diag}\left(e_{n}\right)$, the identity matrix. Let $\|v\|_{p}$ denote the $\ell_{p}$ norm of $v$. We use $q$ to denote the Hölder conjugate of $p$ as a convention, i.e., $\frac{1}{p}+\frac{1}{q}=1$. If $v$ is an operator in the $\ell_{p}$-space, the operator norm of $v$ is then $\|v\|_{q}$. Throughout the paper, we consider the $\ell_{p}$-norm of the input's perturbation, and therefore, the $\ell_{q}$-norm of the gradient, which acts as an operator on the perturbation. A square matrix $X \succeq 0$ means that $X$ is positive semidefinite (PSD). Let $\operatorname{tr}(X)$ be the trace of a square matrix $X$. If $a, b \in \mathbb{R}^{n}$, let $\langle a, b\rangle$ be the inner product of $a$ and $b$.

Lipschitz function. Given two metric spaces $\left(X, d_{X}\right)$ and $\left(Y, d_{Y}\right)$, a function $f: X \rightarrow Y$ is Lipschitz continuous if there exists $K>0$ such that for all $x_{1}, x_{2} \in X$,

$$
\begin{equation*}
d_{Y}\left(f\left(x_{2}\right), f\left(x_{1}\right)\right) \leq K d_{X}\left(x_{2}, x_{1}\right) \tag{1}
\end{equation*}
$$

The smallest such $K$ satisfying Equation (1), denoted by $K_{f}$, is called the Lipschitz constant of $f$. For neural networks, $X$ is in general $\mathbb{R}^{m}$ equipped with the $\ell_{p}$-norm. We will only consider the case when $Y=\mathbb{R}$. In actual applications such as a classification task, a neural network has multiple outputs. The prediction is the class with the maximum score. One can then use the margin between each pair of class predictions and its Lipschitz constant to certify the robustness of a given prediction (Raghunathan et al. 2018, Leino et al. 2021). From Rademacher's theorem, if $f$ is Lipschitz continuous, then $f$ is is almost everywhere differentiable, and $K_{f}=\sup _{x}\|\nabla f(x)\|_{q}$.

Neural network as function. A neural network $f: \mathbb{R}^{m} \rightarrow \mathbb{R}$ is characterized by:

$$
f_{1}(x)=W^{1} x+b_{1} ; f_{i}(x)=W^{i} \sigma\left(f_{i-1}(x)\right)+b_{i}, i=2, \ldots, d
$$

where $W^{i} \in \mathbb{R}^{n_{i+1} \times n_{i}}$ is the weight matrix between the layers, $n_{1}=m, d$ is the depth of the neural network, $\sigma$ denotes an activation function, $b_{i} \in \mathbb{R}^{n_{i+1}}$ is the bias term, and $f=f_{d} \circ \cdots \circ f_{1}$. Because we only consider the $\mathbb{R}$ as the codomain of $f, W^{d}=u \in \mathbb{R}^{1 \times n_{d}}$ is a vector. From chain rule, the gradient of this function is

$$
\begin{equation*}
\nabla f(x)=\left(W^{1}\right)^{T}\left[\operatorname{diag}\left(\sigma^{\prime}\left(f_{1}(x)\right)\right)\left(W^{2}\right)^{T} \cdots \operatorname{diag}\left(\sigma^{\prime}\left(f_{d-1}(x)\right)\right)\left(W^{d}\right)^{T}\right] \tag{2}
\end{equation*}
$$

Common activation functions, including ReLU (Nair and Hinton, 2010), sigmoid functions, and ELU (Clevert et al. 2016) are almost everywhere differentiable. As a result, we are interested in the supremum operator norm of Equation (2).
However, checking all possible inputs $x$ is infeasible, and common activation functions have bounded derivative values, say $[a, b]$. We are then interested in the following value instead:

$$
\begin{equation*}
\max _{v^{i} \in[a, b]^{n_{i}}}\left\|\left(W^{1}\right)^{T} \cdot \operatorname{diag}\left(v^{2}\right) \cdots \cdot \operatorname{diag}\left(v^{d}\right)\left(W^{d}\right)^{T}\right\|_{q}, \tag{3}
\end{equation*}
$$

where $n_{i}$ is the dimension of each $\operatorname{diag}\left(v^{i}\right)$. We call this value the formal global Lipschitz constant (FGL) because we treat all activation functions independent but in reality not all activation patterns are feasible. Therefore, this is an upper bound of the true global Lipschitz constant of the neural network. However, it is the value studied in most global Lipschitzness literature (Scaman and Virmaux, 2018, Fazlyab et al., 2019, Latorre et al., 2020), and also turns out useful in certifying the robustness of neural networks (Raghunathan et al. 2018; Leino et al. 2021; Pauli et al. 2022). We use $\ell_{p}$-FGL to denote the FGL for $\ell_{p}$-perturbations.
In this paper, we focus on the $\ell_{p}$-perturbation on the input, where $p=\infty$ or $p=2$. Because $q$ is the Hölder conjugate of $p$, we are interested in the value of Equation (3), when $q=1$ (for $p=\infty$ ), and $q=2$ (for $p=2$ ). Notice that in the ReLU-network case, $[a, b]$ is $[0,1]$. We will use ReLU-networks as the illustration for the rest of the paper because of the popularity of ReLU in practice and the easy presentation of the 0-1-cube. However, the algorithms presented in this work can be adapted with minor adjustments to other common activation functions.

Remark 2.1. FGL considers all possible activation patterns on the hidden layers, while some of the activation patterns might be unachievable in reality. Therefore, FGL is an upper bound of the true Lipschitz constant. Notice that the activation pattern induced from an input is also decided by the bias term. Therefore, to find the true Lipschitz constant, one has to incorporate the information from the bias term.

## 3 Two-layer neural networks

In this section, we consider the two-layer neural network case. We reduce the FGL estimation to the matrix mixed-norm problem. This immediately yields the computational complexity and approximation algorithms for FGL estimations. In Appendix A.1, we show that we can consider $\{0,1\}$ instead of $[0,1]$ in Equation (3) for two-layer networks.

Problem description. For a two-layer network where $W^{1}=W \in \mathbb{R}^{n \times m}$ and $W^{d}=u \in$ $\mathbb{R}^{1 \times n}$, its FGL (as in Equation (3)) is $\max _{y \in\{0,1\}^{n}}\left\|W^{T} \operatorname{diag}(y) u^{T}\right\|_{q}$, where we use $y$ to denote $v^{1}$ in this case. If we expand the matrix multiplication, it is easy to check that this equals to $\max _{y \in\{0,1\}^{n}}\left\|W^{T} \operatorname{diag}(u) y\right\|_{q}$. Let $A=W^{T} \operatorname{diag}(u)$, then the $\ell_{p}$-FGL is

$$
\begin{equation*}
\max _{y \in\{0,1\}^{n}}\|A y\|_{q} \tag{4}
\end{equation*}
$$

## $3.1 \ell_{\infty}$-FGL estimation

We consider a natural SDP relaxation to Equation (4) when $q=1$, and analyze the result using the celebrated Grothendieck Inequality, which is a fundamental tool in functional analysis.

Mixed-norm problem. The $\infty \rightarrow 1$ mixed-norm of a matrix is defined as

$$
\|A\|_{\infty \rightarrow 1}=\max _{\|x\|_{\infty}=1}\|A x\|_{1}
$$

The mixed-norm problem appears similar to Equation (4) when $q=1$, except for that instead of a norm-1-cube, the cube in Equation (4) is a 0-1-cube. Alon and Naor (2004) showed that it is NP-hard, specifically MAXSNP-hard, to compute the $\infty \rightarrow 1$ mixed-norm of a matrix $A$, via a reduction to the graph Max-Cut problem. Moreover, Alon and Naor (2004) constructed a natural SDP relaxation for the mixed-norm problem:

$$
\begin{gather*}
\max \operatorname{tr}(B X) \\
\text { s.t. } X \succeq 0, X_{i i}=1, i \in[n+m] \tag{5}
\end{gather*}
$$

where $A$ is a submatrix of $B$. We provide the detailed derivation of this relaxation in Appendix A. 3 In fact, this relaxation admits a constant approximation factor. Grothendieck (1956) developed the local theory of Banach spaces, and showed that there exists an absolute value $K_{G}$ such that
Theorem 3.1. For any $m, n \geq 1, A \in \mathbb{R}^{n \times m}$, and any Hilbert space $H$, the following holds:

$$
\max _{u_{i}, v_{j} \in B(H)} \sum_{i, j} A_{i j}\left\langle u_{i}, v_{j}\right\rangle_{H} \leq K_{G}\|A\|_{\infty \rightarrow 1}
$$

where $B(H)$ denotes the unit ball of the Hilbert space.
The precise value of $K_{G}$ is still an outstanding open problem, and it is known that $K_{G}<1.783$ (Krivine, 1979; Braverman et al. 2011). The approximation factor of the SDP relaxation in Equation (5) is $K_{G}$. Similar to the mixed-norm problem, we show that the $\ell_{\infty}$-FGL estimation is MAXSNP-hard and provide an SDP relaxation, which also admits the $K_{G}$-approximation ratio. We provide a detailed explanation on why $K_{G}$ is the approximation ratio and how we can view the SDP relaxation as a geometric transformation in Appendix A. 4
Theorem 3.2. $\ell_{\infty}-F G L$ estimation is MAXSNP-hard.
From 0-1 cube to norm -1 cube. If we can transform the $0-1$ cube in Equation (4) to a norm1 cube, and formulate an equivalent optimization problem, then one can apply the SDP program in Equation (5) to compute an upper bound of the FGL. Indeed, we provide a cube rescaling technique,
and it allows us to construct the SDP for the $\ell_{\infty}$-FGL estimation. We provide the full detail of this technique in Appendix A.5 and the result SDP for the $\ell_{\infty}$-FGL estimation is

$$
\begin{gather*}
\max \frac{1}{2} \operatorname{tr}(B X)  \tag{6}\\
\text { s.t. } X \succeq 0, X_{i i}=1, i \in[n+m+1]
\end{gather*}
$$

where $B$ is a $(n+1+m) \times(n+1+m)$ matrix, and $B=\left(\begin{array}{ccc}0 & 0 & 0 \\ A & A e_{n} & 0\end{array}\right)$. As a result, we have:
Theorem 3.3. There exists a polynomial-time approximation algorithm to estimate the $\ell_{\infty}-F G L$ of two-layer neural networks, moreover, the approximation ratio is $K_{G}$.

## $3.2 \quad \ell_{2}$-FGL estimation

Scaman and Virmaux (2018) showed that the $\ell_{2}$-FGL estimation is NP-hard. If $q=2$ in Equation (4), the objective becomes similar to the $\infty \rightarrow 2$ mixed-norm problem. This is a quadratic optimization problem with a PSD weight matrix over a cube, and can be viewed as a generalization of the graph Max-Cut problem. In the quadratic-optimization formulation of Max-Cut, the weight matrix is the Laplacian of the graph, a special PSD matrix (Goemans and Williamson, 1995). Nesterov (1998) generalized Goemans-Williamson's technique and analyzed the case when the weight matrix is PSD, showing that the natural SDP relaxation in this case has a $\frac{\pi}{2}$-approximation ratio. This provides a $\sqrt{\frac{\pi}{2}}$-approximation algorithm for the $\ell_{2}$-FGL estimation. The approximation ratio comes from a similar inequality to the one in Theorem 3.1, known as the Little Grothendieck Inequality. The SDP for $\ell_{2}$-FGL estimation is:

$$
\begin{align*}
& \max \frac{1}{2} \sqrt{\operatorname{tr}\left(\left(\begin{array}{cc}
A^{T} A & A^{T} A e_{n} \\
e_{n}^{T} A^{T} A & e_{n}^{T} A^{T} A e_{n}
\end{array}\right) X\right)}  \tag{7}\\
& \text { s.t. } X \succeq 0, X_{i i}=1, i \in[n+1] .
\end{align*}
$$

The full derivation is provided in Appendix A.8, and we have the following theorem:
Theorem 3.4. There exists a polynomial-time approximation algorithm to estimate the $\ell_{2}-F G L$ of two-layer neural networks with an approximation factor $\sqrt{\frac{\pi}{2}}$.
Remark 3.5. As we have discussed, for two-layer networks, the $\ell_{p}$-FGL estimation is essentially the $\infty \rightarrow q$ mixed-norm problem. Indeed the mixed-norm problem is an outstanding topic in theoretical computer science. As discussed in Bhattiprolu et al. (2018), the $\infty \rightarrow q$ mixed norm problem has constant approximation algorithms if $q \leq 2$, and is hard to approximate within almost polynomial factors when $q>2$. Because when $q>2$, its Hölder conjugate $p<2$. This implies that for two-layer networks, the FGL estimation can be much harder for $\ell_{p}$-perturbations when $p<2$.
Briët et al. (2017) showed that it is NP-hard to approximate the $\infty \rightarrow 2$ mixed-norm problem better than $\sqrt{ } \frac{\pi}{2}$. Raghavendra and Steurer (2009) proved that assuming the unique games conjecture (Khot 2002), it is INP-hard to approximate the $\infty \rightarrow 1$ mixed-norm problem better than $K_{G}$. These optimal approximation ratios match our SDP relaxations for FGL estimations accordingly.

## 4 Relations to existing SDP works

Before introducing our approach for multi-layer networks, we first examine some existing SDP works on FGL estimations, and discuss their relationships with our natural SDP relaxations in Section 3 .
$\ell_{\infty}$-FGL estimation. Raghunathan et al. (2018) formulated an SDP that only works for two-layer networks. Theirs is essentially the same as ours in Equation (6) (See the detailed comparison in Appendix A.10). However, we provide a rigorous derivation and simpler formulation, and also a sound theoretical analysis of the bound, which illustrate more insights to this problem. Raghunathan et al. (2018) treated the SDP relaxation as a heuristic to a hard quadratic programming problem. We prove that this relaxation is not only a heuristic, but in fact induces an approximation algorithm with a tight bound.
$\ell_{2}$-FGL estimation. Fazlyab et al. (2019) proposed LipSDP, another SDP-based algorithm for the $\ell_{2}$-FGL estimation problem. Fazlyab et al. (2019) provided several variants of LipSDP to balance the
precision and scalability. Pauli et al. (2022) demonstrated that the most precise version of LipSDP, LipSDP-Network, fails to produce an upper bound for $\ell_{2}$-FGL. In this paper, all the references of LipSDP are to LipSDP-Neuron, the less precise version. Surprisingly, even though the approach in LipSDP appears quite different from Equation (7), we show that LipSDP is dual of Equation (7) to estimate the $\ell_{2}$-FGL on two-layer networks. LipSDP for two-layer networks is:

$$
\min _{\zeta, \lambda}\left\{\sqrt{\zeta}:\left(\begin{array}{cc}
-2 a b W^{T} W T-\zeta I_{m} & (a+b) W^{T} T \\
(a+b) T W & -2 T+u^{T} u
\end{array}\right) \preceq 0, \lambda_{i} \geq 0\right\}
$$

where $T=\operatorname{diag}(\lambda)$ for $\lambda \in \mathbb{R}_{+}^{n} ; a$ and $b$ are the lower and upper bounds of the activation's derivative. We will construct a new quadratic program, which we show is equivalent to Equation (4) when $q=2$, and LipSDP is its dual SDP relaxation.

Let the input of the $i$-th activation node on $\operatorname{diag}(y)$ be $y_{i}$, and $w_{i}$ be the row vector of $W$. Hence, $y_{i}=w_{i} x$. Let $\Delta x \in \mathbb{R}^{m}$ be a perturbation on $x$, so $\Delta y_{i}=w_{i} \Delta x$. Let $\Delta \sigma(y) \in \mathbb{R}^{n}$ denote the induced perturbation on $\operatorname{diag}(y)$. The constraint from the activation function is $\frac{\Delta \sigma(y)_{i}}{\Delta y_{i}} \in[a, b]$, in other words, $\frac{\Delta \sigma(y)_{i}}{w_{i} \Delta x} \in[a, b]$. One can write the range constraint as

$$
\left(\Delta \sigma(y)_{i}-a \cdot w_{i} \Delta x\right)\left(\Delta \sigma(y)_{i}-b \cdot w_{i} \Delta x\right) \leq 0
$$

This can be written in the quadratic form:

$$
\binom{w_{i} \Delta x}{\Delta \sigma(y)_{i}}^{T}\left(\begin{array}{cc}
-2 a b & a+b  \tag{8}\\
a+b & -2
\end{array}\right)\binom{w_{i} \Delta x}{\Delta \sigma(y)_{i}} \geq 0, \quad \forall i \in[n] .
$$

Since for the two layer network, $f(x)=u \sigma(y)$, then $\Delta f(x)=u \Delta \sigma(y)$. The objective for $\ell_{2}$-FGL estimation is $\max _{\Delta x, \Delta \sigma(y)} \sqrt{\frac{(u \Delta \sigma(y))^{2}}{(\Delta x)^{2}}}$. The equivalence between this program and Equation $\sqrt[4]{4}$ when $q=2$ is presented in Appendix A. 11 .
Remark 4.1. Another interpretation for the quadratic program is that we want to quantify how the output changes given a data-independent input change, i.e., $\Delta x$. In other words, we want to analyze the effect of $\Delta x$ propagating from the input to the output, with symbolic values rather than actual inputs. The idea is similar to symbolic execution from program analysis (Baldoni et al., 2018).

Duality to LipSDP. Now we will show that LipSDP is the dual SDP to the program formulated above. The dual SDP derivation is of similar form in Ben-Tal and Nemirovski (2001, Ch.4.3.1). Let us introduce a variable $\zeta$ such that $\zeta-\frac{(u \Delta \sigma(y))^{2}}{(\Delta x)^{2}} \geq 0$. In other words,

$$
\begin{equation*}
\zeta(\Delta x)^{2}-(u \Delta \sigma(y))^{2} \geq 0 \tag{9}
\end{equation*}
$$

For each constraint in Equation (8), let us introduce a dual variable $\lambda_{i} \geq 0$. Multiply each constraint with $\lambda_{i}$, then $\binom{\Delta x}{\Delta \sigma(y)_{i}}^{T}\left(\begin{array}{cc}-2 a b \lambda_{i} w_{i}^{T} w_{i} & (a+b) \lambda_{i} w_{i}^{T} \\ (a+b) \lambda_{i} w_{i} & -2 \lambda_{i}\end{array}\right)\binom{\Delta x}{\Delta \sigma(y)_{i}} \geq 0, \quad \forall i \in[n]$. Sum all of them, then we have $\binom{\Delta x}{\Delta \sigma(y)}^{T}\left(\begin{array}{cc}-2 a b W^{T} W T & (a+b) W^{T} T \\ (a+b) T W & -2 T\end{array}\right)\binom{\Delta x}{\Delta \sigma(y)} \geq 0$, where $T=\operatorname{diag}(\lambda)$ is the $n \times n$ diagonal matrix of dual variables $\lambda_{1}, \ldots, \lambda_{n}$.

Equation $\left\{9\right.$ can be rewritten as: $\binom{\Delta x}{\Delta \sigma(y)}^{T}\left(\begin{array}{cc}\zeta I_{m} & 0 \\ 0 & -u^{T} u\end{array}\right)\binom{\Delta x}{\Delta \sigma(y)} \geq 0$. As a result, the dual program for the new optimization program is

$$
\min _{\zeta, \lambda}\left\{\sqrt{\zeta}:\left(\begin{array}{cc}
-2 a b W^{T} W T-\zeta I_{m} & (a+b) W^{T} T \\
(a+b) T W & -2 T+u^{T} u
\end{array}\right) \preceq 0, \lambda_{i} \geq 0\right\} .
$$

Remark 4.2. In Remark 3.5 , we mention that $\sqrt{\frac{\pi}{2}}$ is the optimal approximation ratio for the $\infty \rightarrow 2$ mixed-norm problem, which matches the approximation ratio in Theorem 3.4 Hence, improving the natural SDP relaxation in Equation (7) can be very hard. The duality provides another evidence of LipSDP-Neuron's correctness, and hints that LipSDP-Network, the improved variant, may be wrong.

## $5 \ell_{\infty}$-FGL estimation for multi-layer networks

For a multi-layer neural network, the formal gradient becomes a high-degree polynomial, and its $\ell_{q}$-norm estimation becomes a high-degree polynomial optimization problem over a cube, which is in general a hard problem (Lasserre, 2015). We provide a discussion of the polynomial optimization approach of FGL estimation in Appendix B. Here we provide an SDP dual program of the $\ell_{\infty}$-FGL estimation inspired by the dual SDP approach in Section 4. The difference is that now we consider $\ell_{\infty}$-perturbations to the input space instead of $\ell_{2}$. Hence, the objective becomes

$$
\max _{\Delta x, \Delta \sigma(y)} \frac{|u \Delta \sigma(y)|}{\|\Delta x\|_{\infty}}
$$

If we add an extra constraint $\|\Delta x\|_{\infty}=1$, the above objective becomes

$$
\begin{equation*}
\max _{\Delta x, \Delta \sigma(y)} \frac{1}{2}(u \Delta \sigma(y)+u \Delta \sigma(y)) . \tag{10}
\end{equation*}
$$

The constraints are

$$
\binom{w_{i} \Delta x}{\Delta \sigma(y)_{i}}^{T}\left(\begin{array}{cc}
2 a b & -(a+b) \\
-(a+b) & 2
\end{array}\right)\binom{w_{i} \Delta x}{\Delta \sigma(y)_{i}} \leq 0, \Delta(x)_{j}^{2} \leq 1, \forall i \in[n], j \in[m]
$$

We can write $\Delta(x)_{j}^{2} \leq 1$ as $\binom{1}{\Delta x_{j}}^{T}\left(\begin{array}{cc}1 & 0 \\ 0 & -1\end{array}\right)\binom{1}{\Delta x_{j}} \geq 0$.
Now let us introduce $n+m$ non-negative dual variables $(\tau, \lambda)$, where $\tau \in \mathbb{R}_{+}^{m}$ and $\lambda \in \mathbb{R}_{+}^{n}$. If we multiply each dual variable with the constraint and add all the constraints together, we will have

$$
\left(\begin{array}{c}
1 \\
\Delta x \\
\Delta \sigma(y)
\end{array}\right)^{T}\left(\begin{array}{ccc}
\sum_{j=1}^{m} \tau_{j} & 0 & 0 \\
0 & -2 a b W^{T} W T_{2}-T_{1} & (a+b) W^{T} T_{2} \\
0 & (a+b) T_{2} W & -2 T_{2}
\end{array}\right)\left(\begin{array}{c}
1 \\
\Delta x \\
\Delta \sigma(y)
\end{array}\right) \geq 0
$$

where $T_{1}=\operatorname{diag}(\tau)$ and $T_{2}=\operatorname{diag}(\lambda)$. As a result, we can incorporate the objective Equation 10, and obtain the dual SDP for the $\ell_{\infty}$-FGL estimation:

$$
\min _{\zeta, \lambda, \tau}\left\{\frac{\zeta}{2}:\left(\begin{array}{ccc}
\sum_{j=1}^{m} \tau_{j}-\zeta & 0 & u  \tag{11}\\
0 & -2 a b W^{T} W T_{2}-T_{1} & (a+b) W^{T} T_{2} \\
u^{T} & (a+b) T_{2} W & -2 T_{2}
\end{array}\right) \preceq 0, \lambda_{i}, \tau_{j} \geq 0\right\} .
$$

Remark 5.1. The SDP programs in Section 3 are strictly feasible because the identity matrix is a positive definite solution. Hence, Slater's condition is satisfied and strong duality holds.

Multi-layer extension. We can simply extend the dual program to multiple-layer networks. We first vectorize all the units in the input layer and hidden layers, and then constrain them using layer-wise inequalities to formulate an optimization problem. Let us consider a general $d$-layer multi-layer network, where $W^{i} \in \mathbb{R}^{n_{i+1} \times n_{i}}$ for $i \in[d-1]$, and $W^{d}=u \in \mathbb{R}^{1 \times n_{d}}$. Let $\Delta x$ denote the perturbation on the input layer, $\Delta z^{i}$ be the perturbation on the $i$-th hidden layer, and $w_{j}^{i}$ be the $j$-th row vector of $W^{i}$. The only difference between two layer networks and multi-layer networks is that we have the additional constraints:

$$
\binom{\Delta z^{i}}{\Delta z_{j}^{i+1}}^{T}\left(\begin{array}{cc}
-2 a b\left(w_{j}^{i+1}\right)^{T} w_{j}^{i+1} & (a+b)\left(w_{j}^{i+1}\right)^{T} \\
(a+b) w_{j}^{i+1} & -2
\end{array}\right)\binom{\Delta z^{i}}{\Delta z_{j}^{i+1}} \geq 0
$$

Let $\Lambda_{i} \in \mathbb{R}_{+}^{n_{i}}$ and $T_{i}=\operatorname{diag}\left(\Lambda_{i}\right)$ for $i \in[d]$. Following the similar SDP dual approach, we can add all the constraints together and formulate the following SDP program:

$$
\begin{equation*}
\min _{\zeta, \Lambda_{i}}\left\{\frac{\zeta}{2}:(L+N) \preceq 0, i \in[d]\right\}, \tag{12}
\end{equation*}
$$

where

$$
L=\left(\begin{array}{cccccc}
0 & 0 & 0 & \cdots & 0 & 0  \tag{13}\\
0 & -2 a b\left(W^{1}\right)^{T} W^{1} T_{2} & (a+b)\left(W^{1}\right)^{T} T_{2} & \cdots & 0 & 0 \\
0 & (a+b) T_{2} W^{1} & -2 T_{2}-2 a b\left(W^{2}\right)^{T} W^{2} T_{3} & \cdots & 0 & 0 \\
\vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\
0 & 0 & 0 & \cdots & (a+b) T_{d} W^{d-1} & -2 T_{d}
\end{array}\right)
$$

Table 1: $\ell_{\infty}$-FGL estimation of various methods: DGeoLIP and NGeoLIP induce the same values on two layer networks. DGeoLIP always produces tighter estimations than LiPopt and MP do.

| Network | DGeoLIP | NGeoLIP | LiPopt | MP | Sample | BruF |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2-layer/16 units | 185.18 | 185.18 | 259.44 | 578.54 | 175.24 | 175.24 |
| 2-layer/256 units | 425.04 | 425.04 | 1011.65 | 2697.38 | 306.98 | N/A |
| 8-layer/64 units per layer | 8327.2 | - | N/A | $8.237 * 10^{7}$ | 1130.6 | N/A |

Table 2: Running time (in seconds) of various tools on $\ell_{\infty}$-FGL estimation: DGeoLIP and NGeoLIP are faster than LiPopt. Notice that the running time is implementation and solver-dependent.

| Network | DGeoLIP | NGeoLIP | LiPopt | BruF |
| :---: | :---: | :---: | :---: | :---: |
| 2-layer/16 units | 28.1 | 22.3 | 1572 | 4.8 |
| 2-layer/256 units | 976.0 | 70.9 | 2690 | N/A |
| 8-layer/64 units | 329.5 | - | N/A | N/A |

$$
N=\left(\begin{array}{cccc}
\sum_{k=1}^{n_{1}} \Lambda_{1 k}-\zeta & 0 & \ldots & u \\
0 & -T_{1} & \ldots & 0 \\
\vdots & \vdots & \ddots & \vdots \\
u^{T} & 0 & \cdots & 0
\end{array}\right)
$$

Remark 5.2. If we expand the matrix inequality derived from the compact neural-network representation in Fazlyab et al. (2019, Theorem 2), we will have exactly the same matrix for network constraints as $L$ (Equation (13)) in the dual program formulation. In other words, we provide a compositional optimization interpretation to the compact neural-network representation in LipSDP. With this interpretation, one can extend the SDP to beyond feed-forward structures, such as skip connections (He et al. 2016). Notice that if we apply the similar reasoning to the multi-layer network $\ell_{2}$-FGL estimation, we will obtain LipSDP-Neuron.

## 6 Evaluation and discussion

The primary goal of our work is to provide a theoretical framework, and also algorithms for $\ell_{\infty}$-FGL estimations on practically used networks. The $\ell_{2}$-FGL can be computed using LipSDP. We have implemented the algorithms using MATLAB (MATLAB, 2021), the CVX toolbox (CVX Research, 2020) and MOSEK solver (ApS, 2019), and name the tool GeoLIP. To validate our theory and the applicability of our algorithms, we want to empirically answer the following research questions:

RQ1: Is GeoLIP better than existing methods in terms of precision and scalability?
RQ2: Are the dual SDP programs devised throughout the paper valid?

As we shall see, GeoLIP is indeed better than existing methods in terms of precision and scalability; and the dual SDP programs produce the same values as their natural-SDP-relaxation counterparts.

### 6.1 Experimental design

To answer RQ1, we will run GeoLIP and existing tools that measure the $\ell_{\infty}$-FGL on various feedforward neural networks trained with the MNIST dataset (LeCun and Cortes, 2010). We will record the computed $\ell_{\infty}$-FGL to compare the precision, and the computation time to compare the scalability.
To answer RQ2, we will run the natural SDP relaxations for $\ell_{p}$-FGL estimations proposed in Section 3 , LipSDP for $\ell_{2}$-FGL estimation, and the dual program Equation for $\ell_{\infty}$-FGL on two-layer neural networks, and compare their computed FGLs.

Measurements. Our main baseline tool is LiPopt (Latorre et al., 2020), which is an $\ell_{\infty}$-FGL estimation tool $]^{2}$ Notice that LiPopt is based on the Python Gurobi solver (Gurobi Optimization, LLC 2022), while we use the MATLAB CVX and MOSEK solver. LiPopt relies on a linear programming (LP) hierarchy for the polynomial optimization problem. We use LiPopt-k to denote the $k$-th degree of the LP hierarchy. BruF stands for an brute-force exhaustive enumeration of all possible activation patterns. This is the ground truth for FGL estimations. However, this is an exponential-time search, so we can only run it on networks with a few hidden units. Sample means that we randomly sample 200,000 points in the input space and compute the gradient norm at those points. Notice that this is a lower bound of the true Lipschitz constant, and thus a lower bound of the FGL. MP stands for the weight-matrix-norm-product method. This is a naive upper bound of FGL. We use NGeoLIP to denote the natural SDP relaxations devised in Section 3, and DGeoLIP to denote the dual SDP Equation (12) for $\ell_{\infty}$-FGL estimation. Notice that NGeoLIP only applies to two-layer networks.

We use "-" in the result tables to denote that the experimental setting is not in the scope of the tool's application, and "N/A" to denote the computation takes too much time ( $>20$ hours).

Network setting. We run the experiments on fully-connected feed-forward neural networks, trained with the MNIST dataset for 10 epochs using the ADAM optimizer (Kingma and Ba, 2015). All the trained networks have accuracy greater than $92 \%$ on the test data. For two-layer networks, we use 8 , $16,64,128,256$ hidden nodes. For multiple-layer networks, we consider 3, 7, 8-layer networks, and each hidden layer has 64 ReLU units. Because MNIST has 10 classes, we report the estimated FGL with respect to label 8 as in Latorre et al. (2020), and the average running time per class: we record the total computation time for all 10 classes from each tool, and report the average time per class.

### 6.2 Discussion

We present selected results in Tables 1 and 2, and related major discussions here. The full results, more experimental setup and additional discussions can be found in Appendix C

RQ1. In the experiments of LiPopt, we only used LiPopt-2. In theory, if one can go higher in the LP hierarchy in LiPopt, the result becomes more precise. However, in the case of fully-connected neural networks, using degree-3 in LiPopt is already impractical. For example, on the simplest network that we used, i.e., the single-hidden-layer neural network with 8 hidden units, using LiPopt-3, one $\ell_{\infty}$-FGL computation needs at least 200 hours projected by LiPopt. As a result, for all the LiPopt-related experiments, we were only able to run LiPopt-2. As Latorre et al. (2020) pointed out, the degree has to be at least the depth of the network to compute a valid bound, so we have to use at least LiPopt-k for $k$-layer networks. LiPopt is unable to handle neural networks with more than two layers because this requires LiPopt with degrees beyond 2. Even if we only consider LiPopt-2 on two-layer networks, the running time is still much higher compared to GeoLIP. This demonstrates the great advantage of GeoLIP in terms of scalability compared with LiPopt. If we compare LiPopt-2 with GeoLIP on two-layer networks from Table 1, it is clear that GeoLIP produces more precise results. For networks with depth greater than 2 , we can only compare GeoLIP with the matrix-norm-product method. As we can see from all experiments, GeoLIP's estimation on the FGL is always much lower than MP.

We have also shown that the two-layer network $\ell_{\infty}$-FGL estimation from GeoLIP has a theoretical guarantee with the approximation factor $K_{G}<1.783$ (Theorem 3.3). If we compare the two-layer network results from GeoLIP and Sampling in Table 1, which is a lower bound of true Lipschitz constant, the ratio is within 1.783. This validates our theoretical claim.

RQ2. In Section 4, we have demonstrated the duality between NGeoLIP and LipSDP for the $\ell_{2}$-FGL estimation on two-layer networks, even though the approaches appear drastically different. The experiments show that on two-layer networks, LipSDP and NGeoLIP for $\ell_{2}$-FGL estimations (Table 7 in the appendix), and DGeoLIP and NGeoLIP for $\ell_{\infty}$-FGL estimations produce the same values. These results empirically validate the duality arguments, and also all the related SDP programs.

SDP relaxation. Applying SDP on intractable combinatorial optimization problem was pioneered by the seminal Goemans-Williamson algorithm for the Max-Cut problem (Goemans and Williamson 1995). For two-layer networks, we have reduced the FGL estimation to the mixed-norm problem, and

[^1]provide approximation algorithms with ratios compatible with the known optimal constants in the corresponding mixed-norm problems. Improving them can be a very hard task. We also provide a compositional SDP interpretation of LipSDP-Neuron. Although Pauli et al. (2022) demonstrated the flaw in LipSDP-Network, our compositional SDP interpretation shows that LipSDP-Neuron is correct. In fact, from the compositional SDP interpretation, the program is only constrained by the underlying perturbation geometry and the layer-wise restriction from each hidden unit, so the constraints and objective exactly encode the FGL-estimation problem without additional assumptions. Because often the SDP relaxation for intractable problems gives the optimal known algorithms, we conjecture that GeoLIP and LipSDP are also hard to improve on FGL estimations.
Latorre et al. (2020) used polynomial optimization to address the $\ell_{\infty}$-FGL estimation. We argue that approaching FGL-estimations from the perspective of polynomial optimization loses the accurate characterization of this problem. For example, for two-layer networks, we have provided constant approximation algorithms to estimate FGLs in both $\ell_{\infty}$ and $\ell_{2}$ cases. However, for a general polynomial optimization problem on a cube, we cannot achieve constant approximation. For example, the maximum independent set of a graph can be encoded as a polynomial optimization problem over a cube (Motzkin and Straus, 1965), but the maximum independent set problem cannot be approximated within a constant factor in polynomial time unless $\mathrm{P}=$ NP (Trevisan, 2004).

## 7 Related work

Chen et al. (2020) employed polynomial optimization to compute the true Lipschitz constant of ReLU-networks for $\ell_{\infty}$-perturbations, and used Lasserre's hierarchy of SDPs (Lasserre, 2001) to solve the polynomial optimization problem. However, their approach is highly tailored to ReLU networks, while ours, like LipSDP, can handle common activations, such as sigmoid and ELU.

Latorre et al. (2020) also proposed to use LiPopt to estimate the local Lipschitz constant. However, estimating this quantity is not the problem studied in our work, and there are tools specifically designed for local perturbations and the Lipschitz constant (Laurel et al., 2022; Zhang et al., 2019).
Lipschitz regularization of neural networks is an important task, and recent works (Aziznejad et al. 2020; Bungert et al. 2021; Gouk et al. 2021; Krishnan et al. 2020; Terjék, 2020) have investigated this problem. However, here we study a related but different problem, i.e., Lipschitzness measurement of neural networks. Our work can motivate new Lipschitz regularization techniques.

## 8 Conclusion

In this work, we have provided a quantitative geometric framework for FGL estimations, and also algorithms for the $\ell_{\infty}$-FGL estimation. One important lesson is that when transferring techniques from one perturbations to another ones, we should also transfer the underlying geometry. One future work is to train smooth neural networks using the SDPs proposed in this paper.

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## Checklist

1. For all authors...
(a) Do the main claims made in the abstract and introduction accurately reflect the paper's contributions and scope? [Yes] Our main claims in the abstract and introduction accurately reflect the paper's contributions and scope. In particular, we listed our contributions in Section 1, and forward referenced each contribution to the corresponding section in the paper.
(b) Did you describe the limitations of your work? [Yes] We clearly defined what is the quantity to measure and what the network is in Section2, and what the assumptions are for each theorem in Section 3 .
(c) Did you discuss any potential negative societal impacts of your work? [Yes] We discussed them in Appendix $D$
(d) Have you read the ethics review guidelines and ensured that your paper conforms to them? [Yes] We have read the ethics review guidelines. Because our work is to measure the smoothness of neural networks (see Sections 1 and 2, and we only used the standard MNIST dataset (see Section 6.1), the paper conforms to guidelines.
2. If you are including theoretical results...
(a) Did you state the full set of assumptions of all theoretical results? [Yes] We clearly defined the quantity to measure and the network structures in Sections 2 to 5 .
(b) Did you include complete proofs of all theoretical results? [Yes] We provided important intuition and ideas in the main paper (see Sections 3 to 5], and included complete proofs in Appendix A
3. If you ran experiments...
(a) Did you include the code, data, and instructions needed to reproduce the main experimental results (either in the supplemental material or as a URL)? [Yes] We included the code, data, and instructions as a URL.
(b) Did you specify all the training details (e.g., data splits, hyperparameters, how they were chosen)? [Yes] We included major experimental setting in Section 6.1, and detailed specification in Appendix C. 1 .
(c) Did you report error bars (e.g., with respect to the random seed after running experiments multiple times)? [N/A] Our experiments are deterministic. Given a neural network, our algorithm always returns the same result.
(d) Did you include the total amount of compute and the type of resources used (e.g., type of GPUs, internal cluster, or cloud provider)? [Yes] This was provided in Appendix C. 1
4. If you are using existing assets (e.g., code, data, models) or curating/releasing new assets...
(a) If your work uses existing assets, did you cite the creators? [Yes] In terms of dataset, we only used the standard MNIST, and cited the creators. All the tools used in the paper were properly cited. See Section 6.1 and Appendix C. 1 .
(b) Did you mention the license of the assets? [Yes] See Appendix C. 1 .
(c) Did you include any new assets either in the supplemental material or as a URL? [Yes] We included our code as a URL.
(d) Did you discuss whether and how consent was obtained from people whose data you're using/curating? [N/A]
(e) Did you discuss whether the data you are using/curating contains personally identifiable information or offensive content? [N/A]
5. If you used crowdsourcing or conducted research with human subjects...
(a) Did you include the full text of instructions given to participants and screenshots, if applicable? [N/A]
(b) Did you describe any potential participant risks, with links to Institutional Review Board (IRB) approvals, if applicable? [N/A]
(c) Did you include the estimated hourly wage paid to participants and the total amount spent on participant compensation? [N/A]

## A Elided background, derivations and proofs

## A. 1 Additional analysis background

Gradient as operator. If a function $g: \mathbb{R}^{m} \rightarrow \mathbb{R}$ is a differentiable function at $a \in \mathbb{R}^{m}$, then the total derivative of $g$ at $a$ is

$$
D g(a)=\left[\frac{\partial g}{\partial x_{1}}(a), \ldots, \frac{\partial g}{\partial x_{m}}(a)\right]
$$

and the gradient of $g$ at $a$ is $\nabla g(a)$ is the transpose of $D g(a)$. The linear approximation of $g$ at $a$ is $\langle D g(a), d x\rangle$. Equivalently, we can view the change of a function with respect to an infinitesimal
perturbation as the inner product of $\nabla g(a)$ and $d x$. In this sense, the gradient acts as an operator on the perturbation.

Differentiable activation. Because we want to upper bound the true Lipschitz constant, we only need to show that the quantity considered in the paper indeed upper bounds the true Lipschitz constant considered in the paper. If the activation function is differentiable, then the neural network $f$ is also differentiable, so Equation (3) is trivially true, as proved and applied in Latorre et al. (2020. Theorem 1 and Equation 4).

ReLU activation. For ReLU networks, it is true if we have $[a, b]=[0,1]$. One can consider the (Clarke) generalized Jacobian as in Jordan and Dimakis (2020). At each input point, the Clarke Jacobian is contained in $\left\{\left(W^{1}\right)^{T} \cdot \operatorname{diag}\left(v^{2}\right) \cdots \cdot \operatorname{diag}\left(v^{d}\right)\left(W^{d}\right)^{T} \mid v^{i} \in[a, b]^{n_{i}}\right\}$. Alternatively, we can also use the perturbation propagation argument in Section 4 to see this upper bound. Note that Raghunathan et al. (2018) used this interval representation for ReLU's derivative.

Maximum over hypercube. Now we want to show that the optimization problems over hypercubes considered in this work attain the maximum at the vertices. Without of loss of generality, let us assume the hypercube is $[-1,1]^{n}$. Otherwise, we can transform the hypercube to $[-1,1]^{n}$. Let $A \in \mathbb{R}^{m \times n}, x \in \mathbb{R}^{n}, y \in \mathbb{R}^{m}$, and $z \in \mathbb{R}^{n}$.

We will use the following facts

1. $\|x\|_{1}=\max _{z \in\{-1,1\}^{n}}\langle x, z\rangle$;
2. $\ell_{\infty}$ is the dual of $\ell_{1}$;
3. Let $U \subseteq \mathbb{R}^{n}$. When $\max _{x, z \in U}\langle A x, A z\rangle$ is well-defined, we have $\max _{x \in U}\langle A x, A x\rangle=$ $\max _{x, z \in U}\langle A x, A z\rangle$.

The first fact is from $\|x\|_{1}=\left|x_{1}\right|+\ldots+\left|x_{n}\right|=\max _{z \in\{-1,1\}^{n}}\langle x, z\rangle$. The second fact is from Hölder's inequality for finite-dimensional vector space. For the third one, $\langle A x, A z\rangle$ is maximized only when $A x=A z$. Now we can show that the maximization problems considered in this paper attain the maximum at the hypercube vertices.

$$
\begin{align*}
\max _{\|x\|_{\infty}=1}\|A x\|_{1} & =\max _{\|x\|_{\infty}=1, y \in\{-1,1\}^{m}}\langle A x, y\rangle \quad \text { (From fact (1)) } \\
& =\max _{\|x\|_{\infty}=1, y \in\{-1,1\}^{m}}\left\langle x, A^{T} y\right\rangle \\
& =\max _{y \in\{-1,1\}^{m}}\left\|A^{T} y\right\|_{1}  \tag{14}\\
& =\max _{x \in\{-1,1\}^{n}, y \in\{-1,1\}^{m}}\left\langle x, A^{T} y\right\rangle \\
& =\max _{x \in\{-1,1\}^{n}, y \in\{-1,1\}^{m}}\langle A x, y\rangle . \\
\max _{\|x\|_{\infty}=1}\|A x\|_{2}^{2} & =\max _{\|x\|_{\infty}=1}\langle A x, A x\rangle \\
& =\max _{\|x\|_{\infty}=1,\|z\|_{\infty}=1}\langle A x, A z\rangle \quad \text { (From fact (2)) } \\
& =\max _{\|x\|_{\infty}=1,\|z\|_{\infty}=1}\left\langle A^{T} A x, z\right\rangle .
\end{align*}
$$

Using the similar idea in Equation (14), we have

$$
\max _{\|x\|_{\infty}=1}\|A x\|_{2}^{2}=\max _{x \in\{-1,1\}^{n}, z \in\{-1,1\}^{n}}\langle A x, A z\rangle=\max _{x \in\{-1,1\}^{n}}\langle A x, A x\rangle .
$$

More generally, in the bilinear forms considered above, if $x=x_{1} \otimes \cdots \otimes x_{d}$ is generated by the tensor product of variables over cubes, we can fix one variable and write $x$ as a matrix product, and then move the fixed variable to the hypercube vertices. We can repeat this process to move all variables to the vertices.

## A. 2 Additional definitions

For any neural network $f$, let $O P T(f)$ be the optimal value of Equation (3). We say an algorithm $\mathcal{A}$ is an approximation algorithm for Equation (3) with approximation ratio $\alpha>1$, if $O P T(f) \leq$ $\mathcal{A}(f) \leq \alpha O P T(f)$.

## A. 3 SDP for the $\infty \rightarrow 1$ mixed-norm problem

Recall that for $v \in \mathbb{R}^{m},\|v\|_{1}=\max _{\|u\|_{\infty}=1}\langle u, v\rangle$. We can reformulate the mixed-norm problem as follows:

$$
\max _{x \in\{-1,1\}^{n}}\|A x\|_{1}=\max _{(x, y) \in\{-1,1\}^{n+m}}\langle A x, y\rangle .
$$

If we let $z=\left(\begin{array}{ll}x^{T} & y^{T}\end{array}\right)$, we can have

$$
\max _{(x, y) \in\{-1,1\}^{n+m}}\langle A x, y\rangle=\max _{z \in\{-1,1\}^{n+m}} z \cdot B \cdot z^{T},
$$

where $B$ is a $(m+n) \times(m+n)$ matrix. The last $m$ rows and first $n$ columns of $B$ is $A$, and the rest are 0: $B=\left(\begin{array}{cc}0 & 0 \\ A & 0\end{array}\right)$.
The natural SDP relaxation of the $\infty \rightarrow 1$ mixed-norm problem is:

$$
\begin{gathered}
\max \operatorname{tr}(B X) \\
\text { s.t. } X \succeq 0, X_{i i}=1, i \in[n+m] .
\end{gathered}
$$

In other words, we treat $X$ as the SDP matrix relaxed from the rank-1 matrix $z^{T} \cdot z$.

## A. 4 SDP relaxation and Grothendieck inequalities

In this work, we used the Grothendieck inequality as in Theorem 3.1 .

$$
\begin{equation*}
\max _{u_{i}, v_{j} \in B(H)} \sum_{i, j} A_{i j}\left\langle u_{i}, v_{j}\right\rangle_{H} \leq K_{G}\|A\|_{\infty \rightarrow 1} \tag{15}
\end{equation*}
$$

for any $A \in \mathbb{R}^{n \times m}$; and the little Grothendieck inequality:

$$
\begin{equation*}
\max _{u_{i}, v_{j} \in B(H)} \sum_{i, j}\left(A^{T} A\right)_{i j}\left\langle u_{i}, v_{j}\right\rangle_{H} \leq \frac{\pi}{2}\|A\|_{\infty \rightarrow 2}^{2} \tag{16}
\end{equation*}
$$

Notice that $A^{T} A$ is a PSD matrix.
As discussed in Appendix A.3.

$$
\|A\|_{\infty \rightarrow 1}=\max _{z \in\{-1,1\}^{n+m}} z \cdot B \cdot z^{T}
$$

The natural SDP relaxation is

$$
\begin{gathered}
\max \operatorname{tr}(B X) \\
\text { s.t. } X \succeq 0, X_{i i}=1, i \in[n+m] .
\end{gathered}
$$

Because $X \succeq 0, X=M M^{T}$ for some $M \in \mathbb{R}^{(m+n) \times d}$, where $d \geq 1$. Let $M_{i}$ be the $i$-th row vector of $M . X_{i j}=\left\langle M_{i}, M_{j}\right\rangle$, and $X_{i i}=1$ means $\left\langle M_{i}, M_{i}\right\rangle=1$. As a result, $\operatorname{tr}(B X)=\sum_{i, j} A_{i j} X_{i j}=$ $\sum_{i, j} A_{i j}\left\langle M_{i}, M_{j}\right\rangle_{H}$, where $H$ is the Hilbert space of $\mathbb{R}^{d}$ equipped with the canonical inner product. Thus, Equation (15) implies that $K_{G}$ is the approximation ratio in the SDP relaxation for the mixednorm problem.
In contrast, in the mixed-norm problem, the variable to $B_{i j}$ is $z_{i} z_{j}$, the product of two scalars. If $d=1$ in the SDP relaxation, $M$ is a column vector, and $X$ is a rank- 1 matrix. In this case, the SDP coincides with the combinatorial problem, because the inner product degenerates to the multiplication of two scalars. Hence, the SDP relaxation can be viewed as a continuous relaxation of a discrete problem, and Equation (15) quantifies this geometric transformation. Another interpretation for the SDP relaxation is that SDP drops the rank-1 constraint in the quadratic formulation of the mixed-norm problem.

## A. 5 Rescaling from $0-1$ cube to norm- 1 cube

Now let us show how we transform the $0-1$ cube in Equation (4) to a norm-1 cube, and formulate an equivalent optimization problem. As a result, we can apply the SDP program in Equation (5) to compute an upper bound of the $\ell_{\infty}$-FGL.

Let $x_{i}=\left(t_{i}+1\right) / 2$, where $t_{i} \in\{-1,1\}$. We have

$$
\begin{align*}
& \max _{x \in\{0,1\}^{n}}\|A x\|_{1} \\
= & \max _{x \in\{0,1\}^{n}, y \in\{-1,1\}^{m}} y^{T} A x  \tag{17}\\
= & \max _{t \in\{-1,1\}^{n}, y \in\{-1,1\}^{m}} \frac{1}{2} y^{T} A\left(t+e_{n}\right) .
\end{align*}
$$

Let $O P T_{1}$ be the optimal value of

$$
\max _{(t, y) \in\{-1,1\}^{n+m}} y^{T} A\left(t+e_{n}\right)
$$

Introduce another variable $\tau \in\{-1,1\}$, and let $O P T_{2}$ be the optimal value of

$$
\begin{equation*}
\max _{(t, y, \tau) \in\{-1,1\}^{n+m+1}} y^{T} A\left(t+\tau e_{n}\right) \tag{18}
\end{equation*}
$$

Lemma A.1. $O P T_{1}=O P T_{2}$.
Proof. Clearly $O P T_{2} \geq O P T_{1}$.
Now if $(\hat{t}, \hat{y}, \tau=-1)$ is an optimal solution to Equation $\sqrt{18}$, then $(-\hat{t},-\hat{y}, \tau=1)$ is also an optimal solution, so $O P T_{2} \leq O P T_{1}$.

Now let $z=(t, \tau)$, and we can verify that $y^{T} A\left(t+\tau e_{n}\right)=y^{T} B z$, where $B=\left(\begin{array}{ll}A & A e_{n}\end{array}\right)$.
As a result, the semidefinite program to the $\ell_{\infty}$-FGL constant is

$$
\begin{gathered}
\max \frac{1}{2} \operatorname{tr}(B X) \\
\text { s.t. } X \succeq 0, X_{i i}=1, i \in[n+m+1],
\end{gathered}
$$

where $B$ is a $(n+1+m) \times(n+1+m)$ matrix, and $B=\left(\begin{array}{ccc}0 & 0 & 0 \\ A & A e_{n} & 0\end{array}\right)$.

## A. 6 Proof of Theorem 3.2

Proof. We will use the cube rescaling techniques introduced in Appendix A.5. Alon and Naor (2004) showed that matrix cut-norm is MAXSNP-hard. We will show that if one can solve the FGL estimation problem, then one can find the cut norm of a matrix.
Given a matrix $A$, the cut norm of a matrix $A \in \mathbb{R}^{m \times n}$ is defined as

$$
C N(A)=\max _{x \in\{0,1\}^{n}, y \in\{0,1\}^{m}}\langle A x, y\rangle .
$$

We need to transform $y$ from $0-1$ cube to norm- 1 cube, so similarly let $y_{i}=\left(t_{i}+1\right) / 2$, where $t_{i} \in\{-1,1\}$. The we will have

$$
C N(A)=\max _{x \in\{0,1\}^{n}, y \in\{0,1\}^{m}}\langle A x, y\rangle=\frac{1}{2} \max _{x \in\{0,1\}^{n}, t \in\{-1,1\}^{m}}\left\langle A x,\left(t+e_{m}\right)\right\rangle
$$

Let $B=\binom{A}{e_{m}^{T} A}$. From above we know that

$$
\max _{x \in\{0,1\}^{n}, t \in\{-1,1\}^{m}}\left\langle A x,\left(t+e_{m}\right)\right\rangle=\max _{x \in\{0,1\}^{n},(t, \tau) \in\{-1,1\}^{m+1}}\langle B x,(t, \tau)\rangle .
$$

One can then construct a two layer neural network, where the first weight matrix is $B^{T}$, and the second weight matrix is $(1, \ldots, 1) \in \mathbb{R}^{n}$. Because the network we consider has only one output, the second weight matrix is only a vector. The FGL of this network is exactly twice of the cut norm of $A$.

## A. 7 Proof of Theorem 3.3

Proof. Let $B=\left(\begin{array}{ccc}0 & 0 & 0 \\ A & A e_{n} & 0\end{array}\right)$. Combing Equations 5 , 17 , and 18), the approximation algorithm for Equation (4) where $q=1$ is induced by the following SDP program:

$$
\begin{gathered}
\max \frac{1}{2} \operatorname{tr}(B X) \\
\text { s.t. } X \succeq 0, X_{i i}=1, i \in[n+m+1] .
\end{gathered}
$$

## A. 8 Natural SDP relaxation of $\ell_{2}$-FGL estimation

Now let $q=2$ in Equation (4), we will have:

$$
\max _{y \in\{0,1\}^{n}}\|A y\|_{2}
$$

In other words, we only need to solve the following program:

$$
\begin{equation*}
\max _{z \in\{0,1\}^{n}} z^{T}\left(A^{T} A\right) z \tag{19}
\end{equation*}
$$

Let $M=A^{T} A$, then $M$ is a PSD matrix. We have demonstrated the scaling techniques in Appendix A. 5 Let $x \in\{-1,1\}^{n+1}$, one can verify that

$$
\begin{equation*}
\max _{z \in\{0,1\}^{n}} z^{T} M z=\frac{1}{4} \max _{x \in\{-1,1\}^{n+1}} x^{T} \hat{M} x \tag{20}
\end{equation*}
$$

where $\hat{M}=\left(\begin{array}{cc}M & M e_{n} \\ e_{n}^{T} M & e_{n}^{T} M e_{n}\end{array}\right)$.
It is easy to verify that if $M$ is PSD, $\hat{M}$ is also PSD. Because $M=A^{T} A, \hat{M}=\left(A, A e_{n}\right)^{T} \cdot\left(A, A e_{n}\right)$. Now we can consider the following natural SDP relaxation to $\max _{x \in\{-1,1\}^{n+1}} x^{T} \hat{M} x$ :

$$
\begin{gather*}
\max \operatorname{tr}(\hat{M} X) \\
\text { s.t. } X \succeq 0, X_{i i}=1, i \in[n+1] . \tag{21}
\end{gather*}
$$

This SDP relaxation admits a $\frac{\pi}{2}$-approximation factor from Equation 16 (Rietz, 1974, Nesterov 1998).

## A. 9 Proof of Theorem 3.4

Proof. Let $\hat{M}=\left(\begin{array}{cc}M & M e_{n} \\ e_{n}^{T} M & e_{n}^{T} M e_{n}\end{array}\right)$, where $M=A^{T} A$. Combining Equations 19 , to 21\}, the approximation algorithm for Equation (4) where $q=2$ is induced by the following SDP program:

$$
\begin{gather*}
\max \frac{1}{2} \sqrt{\operatorname{tr}(\hat{M} X)}  \tag{22}\\
\text { s.t. } X \succeq 0, X_{i i}=1, i \in[n+1] .
\end{gather*}
$$

## A. 10 Comparison with Raghunathan et al. (2018)

Raghunathan et al. (2018) formulated the following SDP to upper bound the $\ell_{\infty}$-FGL on two-layer neural networks:

$$
\begin{gather*}
\max \frac{1}{4} \operatorname{tr}(C X)  \tag{23}\\
\text { s.t. } X \succeq 0, X_{i i}=1, i \in[n+m+1]
\end{gather*}
$$

where $C$ is a $(m+n+1) \times(m+n+1)$ matrix, and $C=\left(\begin{array}{ccc}0 & 0 & A^{T} \\ 0 & 0 & e_{n}^{T} A^{T} \\ A & A e_{n} & 0\end{array}\right)$.
If we compare Equations and 23, $C=B+B^{T}$. Because $X$ is symmetric, $\operatorname{tr}(C X)=2 \operatorname{tr}(B X)$. Therefore, Equations (6) and 23 produce the same result.

## A. 11 Equivalence between the new optimization program and Equation 19

Notice that because $u \in \mathbb{R}^{1 \times n}, u \Delta \sigma(x)$ is a scalar. We can view each $z_{i}$ in Equation $\sqrt{19}$ as $\frac{\Delta \sigma(x)_{i}}{\Delta y_{i}}$, the derivative of $\sigma(x)_{i}$ without the limit. Therefore, $\Delta \sigma(x)_{i}=z_{i} \Delta y_{i}$. Recall that from Section 4 , $\Delta y_{i}=w_{i} \Delta x$, so $\Delta \sigma(x)_{i}=w_{i} z_{i} \Delta x$, then $u \Delta \sigma(x)=\Delta x \sum_{i}^{n} u_{i} z_{i} w_{i}=\Delta x(A z)$, where $A=$ $W^{T} \operatorname{diag}(u)$ as defined in Equation (4).
As a result, from Cauchy-Schwarz inequality, the above objective is

$$
\max _{\Delta x, \Delta \sigma(x)} \frac{(u \Delta \sigma(x))^{2}}{(\Delta x)^{2}}=\max _{z}(A z)^{2}, ~ 子, ~ s . t . ~ z \in[a, b]^{n} .
$$

This demonstrates the equivalence between the new optimization program and Equation (19) when $[a, b]=[0,1]$ for $\sigma=$ ReLU .

## B Polynomial optimization approach to the FGL estimation

We briefly discuss the gradient approach to estimate the FGL. Let us use a three layer network as an example:

$$
f(x)=u \sigma\left(V \sigma\left(W x+b_{1}\right)+b_{2}\right)
$$

where $x \in \mathbb{R}^{l \times 1}, W \in \mathbb{R}^{n \times l}, b_{1} \in \mathbb{R}^{n}, V \in \mathbb{R}^{m \times n}, b_{2} \in \mathbb{R}^{m}$ and $u \in \mathbb{R}^{1 \times m}$.
The formal gradient vector of this network is

$$
W^{T} \operatorname{diag}(y) V^{T} \operatorname{diag}(z) u^{T}
$$

where $\operatorname{diag}(y) \in \mathbb{R}^{n \times n}$ and $\operatorname{diag}(z) \in \mathbb{R}^{m \times m}$. The $i$-th component of this vector is then

$$
\sum_{k=1}^{m} \sum_{j=1}^{n}\left(u_{k} V_{k j} W_{j i}\right) \cdot\left(y_{j} z_{k}\right)
$$

Therefore, the $\ell_{p}$-norm estimation of the formal gradient ends up being a polynomial optimization problem over a cube. For example, the $\ell_{1}$-norm (corresponding to $\ell_{\infty}$-perturbations) of the gradient is

$$
\begin{equation*}
\max _{x_{i} \in\{-1,1\}, y_{j} \in\{0,1\}, z_{k} \in\{0,1\}} \sum_{i, j, k=1}^{l, n, m} T_{i j k} \cdot x_{i} y_{j} z_{k} \tag{24}
\end{equation*}
$$

where $T_{i j k}=W_{j i} V_{k j} u_{k}$.
This is essentially a tensor cut-norm problem, and it is an open problem whether there exists an approximation algorithm within a constant factor to the general tensor cut-norm problem (Kannan 2010). Notice that Equation (24) is not a general tensor-cut-norm problem, because the tensor is generated from the weight matrices. For example, if we fix $j$, the projected matrices of $T$ are of rank-1. Each vector in $T_{:, j,:}$ is the product of $V_{k j} u_{k}$ with the vector $W_{j:}$ :

$$
\forall k: T_{:, j, k}=W_{j:} V_{k j} u_{k}
$$

However, we do not have the theoretical technique to exploit the low-rank structure of these special polynomial optimization problems. The perturbation analysis in Sections 4 and 5 can be viewed as exploiting this structure in practice.

Table 3: $\ell_{\infty}$-FGL estimations of different methods for two-layer networks: DGeoLIP and NGeoLIP induce the same estimations, and they are also close to the sampled lower bounds. In the meantime, the result from GeoLIP is tighter than LiPopt's result.

| \#Units | DGEoLIP | NGEoLIP | LIPopt-2 | MP | SAMPLE | BRUF |
| :--- | :---: | :---: | :---: | :---: | :---: | ---: |
| 8 | 142.19 | 142.19 | 180.38 | 411.90 | 134.76 | 134.76 |
| 16 | 185.18 | 185.18 | 259.44 | 578.54 | 175.24 | 175.24 |
| 64 | 287.60 | 287.60 | 510.00 | 1207.70 | 253.89 | N/A |
| 128 | 346.27 | 346.27 | 780.46 | 2004.34 | 266.22 | N/A |
| 256 | 425.04 | 425.04 | 1011.65 | 2697.38 | 306.98 | N/A |

Table 4: Average running time (in seconds) of different methods for two-layer-network $\ell_{\infty}$-FGL estimations: GeoLIPs are faster than LiPopt.

| \# Hidden Units | DGeoLiP | NGEoLIP | LiPopt-2 | BruF |
| :--- | :---: | :---: | :---: | :---: |
| 8 | 23.1 | 21.5 | 1533 | $<0.1$ |
| 16 | 28.1 | 22.3 | 1572 | 4.8 |
| 64 | 93.4 | 31.7 | 1831 | N/A |
| 128 | 292.5 | 42.2 | 2055 | N/A |
| 256 | 976.0 | 70.9 | 2690 | N/A |

## C Complete experimental specifications and results

GeoLIP is available at https://github.com/z1w/GeoLIP. To accommodate users who do not have access to MATLAB, we also implement a version based on CVXPY (Diamond and Boyd 2016). However, the MATLAB implementation works more efficiently in terms of memory and speed, and we encourage users to work with the MATLAB version when possible. We conducted all the GeoLIP-related experiments with the MATLAB version.

## C. 1 Experimental specifications

Tools. We obtain the LiPopt implementation from https://github.com/latorrefabian/ lipopt, under the MIT License.

Server specification. All the experiments are run on a workstation with forty-eight Intel ${ }^{\circledR}$ Xeon ${ }^{\circledR}$ Silver 4214 CPUs running at 2.20 GHz , and 258 GB of memory, and eight Nvidia GeForce RTX 2080 Ti GPUs. Each GPU has 4352 CUDA cores and 11 GB of GDDR6 memory.

Dataset and split. We used the standard MNIST dataset from the PyTorch package (Paszke et al. 2019). We used the "train" parameter in the MNIST function to split training and testing data.

## C. 2 Experimental results

Single hidden layer. We consider the $\ell_{\infty}$-FGL estimation on two layer neural networks with different numbers of hidden units. The results are summarized in Tables 3 and 4 ,

Multiple hidden layers. We consider the $\ell_{\infty}$-FGL estimation on 3, 7, 8-layer neural networks. Each hidden layer in the network has 64 ReLU units. The results are summarized in Tables 5 and 6
$\ell_{2}$-FGL estimation. We measure the $\ell_{2}$-FGL on two-layer networks mainly to compare whether Equation (7) and LipSDP produce the same result. Additionally, we also want to empirically examine the approximation guarantee from Theorem 3.4 Still, we consider networks with 8 , $16,64,128,256$ hidden nodes. The results are summarized in Tables 7 and 8 .

Table 5: $\ell_{\infty}$-FGL estimations of different methods for multi-layer networks: GeoLIP's result is much tighter than the matrix-product method, and LiPopt is unable to handle these networks.

| \# Layers | GEoLIP | Matrix Product | Sample | LIPtopt |
| :--- | :---: | :---: | :---: | :---: |
| 3 | 529.42 | 9023.65 | 311.88 | N/A |
| 7 | 5156.5 | $1.423 * 10^{7}$ | 1168.8 | N/A |
| 8 | 8327.2 | $8.237 * 10^{7}$ | 1130.6 | N/A |

Table 6: Average running time (in seconds) of GeoLIP for multi-layer-network $\ell_{\infty}$-FGL estimations.

| 3-Layer Net | 7-Layer Net | 8-Layer Net |
| :--- | :---: | :---: |
| 120.9 | 284.3 | 329.5 |

## C. 3 Additional discussions

Duality. The results in Table 7 show that the results of LipSDP and GeoLIP on two-layer-network $\ell_{2}$-FGL estimation are exactly the same, which empirically demonstrates the duality between LipSDP and GeoLIP, as discussed in Section 4. Though Pauli et al. (2022) showed that the most precise version of LipSDP is invalid for estimating an upper bound of $\ell_{2}-\mathrm{FGL}$, our dual-program argument shows that the less precise version of LipSDP is correct.

Precision. We showed that GeoLIP's approximation factor for the $\ell_{2}$-FGL estimation on two layer networks is $\sqrt{\frac{\pi}{2}} \approx 1.253$ in Theorem 3.4. The $\ell_{2}$-FGL from GeoLIP is very close to the sampled lower bound of true Lipschitz constant in Table 7 On the other hand, because the result from GeoLIP is an upper bound of FGL, and this result is not much greater than the sampled lower bound of true Lipschitz constant, this empirically demonstrates that the true Lipschitz constant is not very different from the FGL on two-layer networks.

Running time. If we compare the running time in Tables 4 and 8 the dual program takes more time to solve than the natural relaxation. This is particularly true when the number of hidden neurons increases. From the reported numbers of variables and equality constraints by CVX, the dual program and natural relaxation have similar numbers. It is also observed that the CPU usage is higher when the natural relaxation is being solved. We want to point out that the running time and optimization algorithm are solver-dependent, and efficiently solving SDP is beyond the scope of this work. It is an interesting future direction to exploit the block structure of the dual programs, and develop algorithms that are compatible with those programs, because training smooth networks is a critical task, and it is promising to incorporate the SDP programs.
$\ell_{2}$ versus $\ell_{\infty}$ FGLs. If we compare results from Tables 3 and 7 , we can also find that the discrepancy between matrix product method and sampled lower bound is much smaller in the $\ell_{2}$ case. This could also explain why Gloro works for $\ell_{2}$-perturbations but not the $\ell_{\infty}$ case in practice, where Leino et al. (2021) used matrix-norm product to upper bound the Lipschitz constant of the network in Gloro.

Sampling. Sampling can only give a lower bound of the true Lipschitz constant, while we are trying to estimate an upper bound. We use sampling as a sanity check to ensure that the SDP method is at least sound and indeed provides an upper bound of the FGL. It is interesting to see that in networks where we can brute-force enumerate all the activation patterns, sampling provides very close results to the ground-truth ones. Notice that for those networks, there are only a few hidden units (8 or 16), while we sample many $(200,000)$ inputs, which might activate all or most of the patterns. However, for networks with many activation nodes, it is infeasible to have a brute-force enumeration of all the activation patterns, so we do not have the ground-truth information. Sampling has no guarantee whether it can activate all patterns unless we have sampled all possible inputs, which is also impractical.

Multi-layer network guarantees. The discrepancy between the results from sampling and GeoLIP is relatively large for multi-layer networks. The approximation guarantee of GeoLIP is in terms of

Table 7: $\ell_{2}$-FGL estimations of different methods for two-layer networks: LipSDP and NGeoLIP induce the same estimations, and these results are also close to the sampled lower bounds.

| \#Units | NGEOLIP | LIPSDP | MP | SAMPLE | BRUF |
| :--- | :---: | :---: | :---: | :---: | :---: |
| 8 | 6.531 | 6.531 | 11.035 | 6.527 | 6.527 |
| 16 | 8.801 | 8.801 | 13.936 | 8.795 | 8.799 |
| 64 | 12.573 | 12.573 | 22.501 | 11.901 | N/A |
| 128 | 15.205 | 15.205 | 30.972 | 13.030 | N/A |
| 256 | 18.590 | 18.590 | 35.716 | 14.610 | N/A |

Table 8: Average running time (in seconds) of LipSDP and NGeoLIP for two-layer-network $\ell_{2}$-FGL estimations.

| \# Hidden Units | LipSDP | NGEoLIP | BruF |
| :--- | :---: | :---: | :---: |
| 8 | 11.5 | 1.2 | $<0.1$ |
| 16 | 15.7 | 1.2 | 5.1 |
| 64 | 64.2 | 1.3 | N/A |
| 128 | 216.1 | 1.7 | N/A |
| 256 | 758.1 | 4.1 | N/A |

the FGL, rather than true Lipschitz constant. It is unclear how large the gap between true Lipschitz constant and the FGL is for multi-layer networks. Narrowing this gap is an interesting research direction and beyond the scope of this work. We do not know whether for multi-layer networks, GeoLIP has an approximation guarantee that is independent of the network. We leave this as an open problem.

## D Negative societal impacts

Our work is mainly theoretical and to measure an intrinsic mathematical property of neural networks, and can benefit the verification of deep-learning systems. A misuse of our work can give a false sense of safety, so the practical use of our work should be careful.


[^0]:    ${ }^{1}$ Our code is available at https://github.com/z1w/GeoLIP

[^1]:    ${ }^{2}$ Another method was proposed by Chen et al. 2020 , however, the code is not available and we are not able to compare it with GeoLIP.

