Robust Learning against Relational Adversaries

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Abstract

Test-time adversarial attacks have posed serious challenges to the robustness of machine-learning models, and in many settings the adversarial perturbation needs not be bounded by small $\ell_p$-norms. Motivated by attacks in program analysis and security tasks, we investigate relational adversaries, a broad class of attackers who create adversarial examples in a reflexive-transitive closure of a logical relation. We analyze the conditions for robustness against relational adversaries and investigate different levels of robustness-accuracy trade-off due to various patterns in a relation. Inspired by the insights, we propose normalize-and-predict, a learning framework that leverages input normalization to achieve provable robustness. The framework solves the pain points of adversarial training against relational adversaries and can be combined with adversarial training for the benefits of both approaches. Guided by our theoretical findings, we apply our framework to source code authorship attribution and malware detection. Results of both tasks show our learning framework significantly improves the robustness of models against relational adversaries. In the process, it outperforms adversarial training, the most noteworthy defense mechanism, by a wide margin.

1 Introduction

The robustness of machine learning (ML) systems has been challenged by test-time attacks using adversarial examples [Szegedy et al. 2013]. These adversarial examples are intentionally manipulated inputs that preserve the essential characteristics of the original inputs, and thus are expected to have the same test outcome as the originals by human standard; yet they severely affect the performance of many ML models across different domains [Moosavi-Dezfooli et al. 2016, Eykholt et al. 2018, Qin et al. 2019]. As models in high-stake domains such as system security are also undermined by attacks [Grosse et al. 2017, Rosenberg et al. 2018, Hu and Tan 2018, Pierazzi et al. 2020], robust ML in adversarial test environment becomes an imperative task for the ML community.

Existing work on test-time attack and defense evaluation predominately considers $\ell_p$-norm bounded adversarial manipulation [Goodfellow et al. 2014, Carlini and Wagner 2017, Papernot et al. 2018].

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We evaluate API calls, and (2) substitution of equivalent API calls — as logical relations. Next, we propose two Normalize-and-Predict Learning Framework which incorporate model robustness as part of the training objective. However, such regularization do not strictly enforce the robustness, as a result, models are still vulnerable.

In this paper, we consider a general threat model — the relational adversary — in which the attacker can manipulate the original test inputs via transformations specified by a logical relation. Unlike the prior work [Rosenberg et al., 2018; Hu and Tan, 2018; Hosseini et al., 2017; Hosseini and Poovendran, 2018], which investigates specific adversarial settings, our paper extends the scope of attacks to general logical transformation, which can be readily instantiated to incorporate real-world transformations over different domains. Moreover, the relational adversary can apply an arbitrary sequence of transformations to the inputs as long as the essential semantics of the input is preserved.

From the defense perspective, recent work has started to look beyond ℓ_p-norm constraints, including adversarial training [Grosse et al., 2017; Rosenberg et al., 2019; Lei et al., 2019; Laidlaw et al., 2020; Li et al., 2022], verification-loss regularization [Huang et al., 2019] and invariance-induced regularization [Yang et al., 2019]. Adversarial training in principle can achieve high robust accuracy when the adversarial example in the training loop maximizes the loss. However, finding such adversarial examples is in general NP-hard [Katz et al., 2017] and even PSPACE-hard for attacks considered in this paper (Appendix A.1). [Huang et al., 2019] and Yang et al. [2019] add regularizers which incorporate model robustness as part of the training objective. However, such regularization do not strictly enforce the robustness, as a result, models are still vulnerable.

**Normalize-and-Predict Learning Framework** This paper attempts to overcome the limitations of prior work by introducing a learning framework, normalize-and-predict (hereinafter abbreviated as NaP), which guarantees robustness by design. Unlike the prior work which exclusively consider the ℓ_p-norm bounded attacks, we target a relational adversary, whose admissible manipulation is specified by a logical relation. We consider a strong adversary who can apply an arbitrary number of transformations. The key idea underlying NaP is normalization, a powerful concept in computer science — canonicalizing data with multiple representations into a unique form — that is widely applied across various domains (e.g. canonicalization of filenames for computer security, lemmatisation in computational linguistics, etc.). Technically, NaP first converts each data point to a canonical form and subsequently restricts the training and test of models on the normalized data. A potential downside of NaP is the sacrifice of model accuracy in exchange of guaranteed robustness. To cope with this issue, we propose a unified framework that combines NaP with adversarial training in an attempt to achieve the optimal robust-accuracy trade-off. Specifically, the unified framework selectively normalizes relations over which model accuracy can be preserved and adversarially trains on the rest. Our unified framework gets the benefits from both NaP and adversarial training.

We evaluate NaP in the settings of source code authorship attribution and malware detection. For the former, we set out to defend two attribution approaches [Caliskan-Islam et al., 2015; Abuhamad et al., 2018] against the state-of-the-art attack proposed by Quiring et al. [2019]. For the latter, we first formulate two types of common program transformation — (1) addition of redundant libraries and API calls, and (2) substitution of equivalent API calls — as logical relations. Next, we propose two generic relational adversarial attacks to determine the robustness of a model.

The results we obtained in both tasks show that:

1. Models obtained by NaP and the unified framework achieve significantly higher robust accuracy than the vanilla models. In particular, the improvement of robustness far outweighs the drop in accuracy on clean inputs, suggesting a worthwhile trade-off when a sizable portion of the input comes from adversarial sources.
2. NaP achieves higher robust accuracy than adversarial training especially when attackers use a stronger or different attack method than the one used in NaP or adversarial training.
3. Compared to adversarial training, NaP incurs a substantially lower computation overhead when defending against problem space attacks, where adversarial examples are generated through program transformation on the raw code samples rather than gradient manipulation on input features.
Finally, based on our theoretical and empirical results, we conclude that input normalization is vital to robust learning against relational adversaries. We believe techniques that can improve the quality of normalization are promising directions for future work.

2 Related Work.

Test-time attacks using adversarial examples have been extensively studied in the past several years. Research has shown ML models are vulnerable to such attack in a variety of application domains [Moosavi-Dezfooli et al., 2016, Chen et al., 2017, Papernot et al., 2017, Eykholt et al., 2018, Ebrahimi et al., 2018, Qin et al., 2019, Yang et al., 2020a] including system security and program analysis where reliable defense is essential. For instance, Grosse et al. [2017] and Al-Dujaili et al. [2018] evade API/library usage based malware detectors by adding redundant API calls; Rosenberg et al. [2018], Hu and Tan [2018], and Rosenberg et al. [2019] successfully attack running-time behavior based detectors by adding redundant execution traces; Quiring et al. [2019] shows semantic-preserving edits over source code can evade authorship attribution. Pierazzi et al. [2020] extend the attacks from feature-space to problem-space to create realistic, executable attack instances using automated software transplantation.

On the defense end, adversarial training has been the most widely used approach [Grosse et al., 2017, Al-Dujaili et al., 2018, Rosenberg et al., 2019, Li et al., 2022]. In particular, Al-Dujaili et al. [2018] and Li et al. [2022] attempt adversarial training on the same malware detection and authorship attribution tasks in our experiment evaluation, respectively. We show that adversarial training is hard to optimize and ineffective against different attack algorithms or attacks with stronger search parameters than in training. Yang et al. [2019] adds invariance-induced regularizers which indicate the worst-case loss of searching over input transformations. Hendrycks et al. [2019] proposes self-supervised learning in which the learner tries to identify the transformations applied to the training inputs. We extend the scope from specific spatial transformation attacks in image classification to a general adversary based on logic relations. We note that regularization and self-supervised learning may not enforce model robustness on finite samples, nor will they be computationally efficient over arbitrarily long sequence of input transformations. In contrast, NaP enforces robustness by design. Incer et al. [2018, Kouzmenchko [2018] enforce monotonicity over model outputs so that the addition of feature values always increase the maliciousness score. These approaches are limited to guarding against the addition attacks, thus lacks generality.

Normalization is a technique to reduce the number of syntactically distinct instances. First introduced to network security in the early 2000s in the context of intrusion detection systems [Handley et al., 2001], it was later applied to malware detection [Christodorescu et al., 2007, Coogan et al., 2011, Bichsel et al., 2016, Salem and Banesuc, 2016, Baumann et al., 2017]. Our work addresses the open question whether normalization is useful for ML under relational adversary by investigating its impact on both model robustness and accuracy.

The robustness-accuracy trade-off against \(\ell_p\)-attacks has been discussed in literature [Zhang et al., 2019, Tsipras et al., 2018, Fawzi et al., 2018, Raghunathan et al., 2020, Yang et al., 2020c]. The trade-off can either be 1) distributional, where inputs with different Bayes-optimal labels are mixed up in the adversarial feasible set, or 2) algorithmic, where the trade-off is due to low model expressiveness, small sample size, inductive bias in training data or deficiency in the learning algorithm. The former is inevitable because no classifier can achieve both optimal robustness and accuracy. Zhang et al. [2019] construct such an example distribution: the real number line is divided into intervals of size \(\epsilon\), and neighboring segments always have different labels. Our analysis focuses on distributional trade-off because NaP guarantees robustness against the normalized transformations, i.e. there is no further algorithmic trade-off if normalization is exact. The analysis helps strategically choose the set of transformations to normalize in the unified framework. We also note that NaP can use any existing training/testing procedure except the inputs are now normalized. Therefore, existing \(\ell_p\)-based trade-off analysis also applies on the normalized input space for \(\ell_p\)-attacks.

3 Threat Model

We consider a data distribution \(\mathcal{D}\) over an input space \(\mathcal{X}\) and categorical label space \(\mathcal{Y}\). We use bold face letters, e.g. \(\mathbf{x}\), for input vectors and \(y\) for the label. Given a hypothesis class \(\mathcal{H}\), the learner
Attacker’s Capability. We assume the attacker wants to learn a classifier \( f : \mathcal{X} \rightarrow \mathcal{Y} \) in \( \mathcal{H} \) that minimizes the risk over the data distribution. In non-adversarial settings, the learner solves \( \min_{f \in \mathcal{H}} \mathbb{E}_{(x,y) \sim D} \ell(f(x), y) \), where \( \ell \) is a loss function. For classification, \( \ell(f(x), y) = 1(f(x) \neq y) \).

### Logical Relation
A relation \( R \) is a set of input pairs. Each pair \((x, z) \in R\) implies a viable transformation from \( x \) to \( z \). We write \( x \xrightarrow{R} z \) iff \((x, z) \in R\), and \( x \xrightarrow{R_\ell} z \) iff \( x \) can arrive at \( z \) via an arbitrary number of transformations specified by \( R \). In other words, \( \xrightarrow{R_\ell} \) is the reflexive-transitive closure of \( \xrightarrow{R} \). We describe an example relation as follows:

**Example 1** (Semantic-preserving Program Transformations). Let \( P \) denote a set of programs and \( T \) be a set of transformations that preserves program semantics. Then \( T \) induces a relation \( R = \{ (p, t(p)) \mid p \in P, t \in T \} \) on the space \( P \). Figure 1 shows a concrete example. The initial program \( p_0 \) undergoes (1) for-to-while, (2) printf-to-cout, and (3) 0/1-to-false/true transformations to become \( p_3 \). We have \( p_i \xrightarrow{R} p_{i+1} \) and \( p_i \xrightarrow{R_\ell} p_j \) for \( i \in \{0, 1, 2\}, j \geq i \).

Attacker’s Capability. A test-time adversary replaces a clean test input \( x \) with an adversarially manipulated input \( A(x) \), where \( A(\cdot) \) represents the attack algorithm that searches for adversarial examples in a feasible set \( T(x) \). We consider an adversary who wants to maximize the classification error rate: \( \mathbb{E}_{(x,y) \sim D} 1(f(A(x)) \neq y) \). We assume white-box attacks, i.e., the adversary has total access to \( f \), including its structures, model parameters and any defense mechanism in place.

**Definition 1** (relational adversary). An adversary is \( R \)-relational if \( T(x) = \{ z \mid x \xrightarrow{R} z \} \).

In essence, a relational adversary can apply arbitrary composition of transformations specified in \( R \). We assume \( R \) to be accessible to both the learner and the adversary. We believe this is a reasonable assumption to make: \( R \) determines the adversary’s search space like \( \epsilon \) and \( p \) for \( \ell_p \)-norm attacks, and it is impossible to formally evaluate a defense technique against unknown adversarial feasible sets.

Challenges with Relational Adversaries. Relational adversaries are prevalent in program analysis because 1) the attacker often has enough control over the victim input to make substantial changes, and 2) well-defined semantic-preserving transformations exist. We will focus on two such threats — misleading authorship attribution [Quiring et al., 2019] and evading malware detection [Al-Dujaili et al., 2018] — in the empirical evaluation in Sec 5. In the former, the attacker transforms source code samples so as to evade automated authorship attribution while maintains the behavior of the code. In the latter, a malware author manipulates the API usage to evade ML-based detectors while maintains the malicious functionality.

Relational adversaries pose new challenges to robust learning. First, since \( T(x) \) is discrete in nature and the perturbation need not be bounded by \( \ell_p \)-norm, defense mechanisms that leverage local smoothness of model prediction are no longer applicable. Second, adversarial training specifically suffers from efficiency issues against relational adversaries. The reasons are two two-fold. First, the inner maximization procedure needs to search a combinatorially large discrete space. Second, the generation of adversarial examples is based on program transformation in the problem space (rather than gradient manipulation in the feature space), which can not be performed in GPU. These challenges motivate us to study the essence of \( R \) and leverage normalization as a remedy (Sec 4). We evaluate models by the robustness and robust accuracy adapted for relational adversaries as follows,

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1. We consider a strong white-box attacker to avoid interference from security by obscurity, which is shown fragile in various other adversarial settings [Carlini and Wagner, 2017].

2. This is to ensure the syntactic validity of the generated program.
We note that all three patterns are common in program analysis tasks. The second pattern happens when some syntactical feature is strongly correlated with the class label. For example, a malware can be transformed to an adversarial example and the original clean input share the same normal form.

Definition 2 (Robustness and robust accuracy). Let \( Q(R, f, x) \) be the following statement: \( \forall z ((x \rightarrow R) z) \Rightarrow f(x) = f(z) \). Then, a classifier \( f \) is robust at \( x \) if \( Q(R, f, x) \) is true, and the robustness of \( f \) to an \( R \)-relational adversary is: \( \mathbb{E}_{x \sim D_x} 1_{Q(R, f, x)} \), where \( 1_{(.)} \) indicates the truth value of a statement and \( D_x \) is the marginal distribution over inputs. The robust accuracy of \( f \) w.r.t. an \( R \)-relational adversary is then: \( \mathbb{E}_{(x,y) \sim D} 1_{Q(R, f, x) \wedge f(x) = y} \).

4 NaP– A Robust Learning Framework

The Normalize-and-Predict (NaP) framework enhances model robustness by learning and testing over normalized training and test inputs. The framework originates from the following principle:

Suppose we can convert a test input into some canonical form — the normal form — and use the normal form as the model input, then the model prediction is robust if the adversarial example and the original clean input share the same normal form.

We refer the conversion of an input to its normal form as normalization. In this section, we answer the crucial questions of when and how to normalize. We first introduce the framework in Sec 4.1 and then theoretically analyze its performance in light of robustness-accuracy trade-off in Sec 4.2. The analysis shows that a carefully chosen normal form can help achieve the optimal robust accuracy. Following these insights, we show a unified framework of NaP with adversarial training in Sec 4.3 and a computationally efficient heuristic normalizer in Sec 4.4.

4.1 An Overview of the NaP Framework

In NaP, the learner first specifies a normalizer \( \mathcal{N} : \mathcal{X} \rightarrow \mathcal{X} \). We call \( \mathcal{N}(x) \) the ‘normal form’ of input \( x \). The learner then both trains the classifier and predicts the test label over the normal forms instead of the original inputs. Let \( D \) denote the training set. In the empirical risk minimization learning scheme, for example, the learner will now solve the following problem

\[
\min_{f \in \mathcal{H}} \sum_{(x,y) \in D} \ell(f(\mathcal{N}(x)), y),
\]

and use the minimizer \( f^* \) as the classifier. For an actual learning algorithm, the NaP pipeline will replace the training input \( (x, y) \) with \( (\mathcal{N}(x), y) \). At test-time, the model will predict \( f^*(\mathcal{N}(x)) \), i.e., replace the original test input \( x \) with the normal form \( \mathcal{N}(x) \).

4.2 Finding the Normalizer — Trade-off Analysis

The choice of \( \mathcal{N} \) is crucial to NaP’s performance in terms of the robustness-accuracy trade-off. For an input \( x \), let \( S_{\mathcal{N}}(x) = \{ z \in \mathcal{X} \mid \mathcal{N}(x) = \mathcal{N}(z) \} \) denote the set of inputs that share the same normal form as \( x \). For robustness purpose, we want \( S_{\mathcal{N}}(x) \) to be inclusive: if \( S_{\mathcal{N}}(x) \) covers the adversary’s feasible set \( T(x) \), then the model will be robust at \( x \) by design. Meanwhile, for accuracy purpose, we want to restrict the size of \( S_{\mathcal{N}}(x) \): a constant \( \mathcal{N} \) is robust, but has no utility as \( f(\mathcal{N}(\cdot)) \) is also constant. Therefore, we seek an \( \mathcal{N} \) that performs only the necessary normalization for robustness with a minimal impact on accuracy.

Price for Robustness. Like against \( \ell_p \) attacks, robustness-accuracy trade-off also exists for learning against relational attacks. We first examine the intrinsic trade-off due to the structure of the relation. Given a relation \( R \), we have the following condition for a classifier to be robust everywhere.

Proposition 1. A classifier \( f \) is robust at all \( x \in \mathcal{X} \) iff \( x \rightarrow R z \implies f(x) = f(z) \) for all \( x \in \mathcal{X} \).

Under this condition, we characterize three interesting patterns of \( R \) that cause different levels of trade-off as shown in Figure 2. First, if \( x \rightarrow R z \) and \( x, z \) have the same label, then a robust classifier \( f \) comes at no cost of natural accuracy. Second, if \( x, z \) have different labels but can be transformed into each other under \( R \), then a robust \( f \) will have to sacrifice the natural accuracy for either \( x \) or \( z \). Fortunately, as we will explain in Theorem 1 such cost to natural accuracy is indeed necessary for achieving the best robust accuracy. Last, if \( x \) can be transformed to two outputs \( z_1, z_2 \) with different labels, then enforcing robustness may cause additional trade-off to robust accuracy.

We note that all three patterns are common in program analysis tasks. The second pattern happens when some syntactical feature is strongly correlated with the class label. For example, a malware can
call two different APIs — one is open-sourced and the other is a secured version with authentication — for the same function. Malware authors at large predominately use the former for convenience, and thus make its usage a strong signal of malware. However, an advanced attacker may compromise the authentication and subsequently uses the secured API to evade detection. NaP avoids using such features and thus corrects the false sense of security. The third pattern often happens with code injection: x could be an abstract function, while \( z_1, z_2 \) are two instantiations with different purposes. Enforcing the same prediction over \( x, z_1 \) and \( z_2 \) makes no sense in most cases.

**Optimal \( N \) for Robust Accuracy.** Inspired by the robustness-accuracy trade-off analysis, we discover the following normalizer \( N \) that preserves the optimal robust accuracy before and after normalization.

**Definition 3 (Equivalence Group).** A set \( E \subseteq X \) is an equivalence group under relation \( R \) iff 1) \( \forall x, z \in E, \text{ we always have } x \rightarrow_R z \text{ and } z \rightarrow_R x \text{, and } \) 2) \( \forall x, z \in X, (x \rightarrow_R z) \land (z \rightarrow_R x) \) implies \( x, z \in E \). In addition, a single-element set \( \{ x \} \) is an equivalence group if \( x \) does not belong to any equivalence group with more than one element.

**Definition 4 (Normalizer by Equivalence Group).** A normalizer by equivalence group \( N \) picks a deterministic element \( z \) for each equivalence group \( E \), and returns \( N(x) = z \) for all \( x \in E \).

**Theorem 1 (Preservation of Robust Accuracy).** Let \( H \) be the set of all labeling functions and \( N \) be a normalizer by equivalence group, then we must have a classifier \( f \in H \) such that 1) \( f \) has the highest robust accuracy on \( D \) without normalization, 2) \( f(x) = f(z) \) for all \( x, z \) in the same equivalence group, and 3) there exists a classifier \( g \in H \) over the normalized inputs such that \( g(N(x)) = f(x), \forall x \in X \).

Theorem 1 convey an important message. The robustness guaranteed by NaP with a normalizer in Definition 4, i.e. same prediction for all \( x \) in the same equivalence group, is also desired by some model that achieve the highest robustness accuracy on the original unnormalized inputs. Therefore, learning and testing on the normalized inputs incurs no additional robust-accuracy trade-off.

### 4.3 Synergy with Adversarial Training

Despite the difference between the underlying principles of NaP and adversarial training, we can readily unify these two approaches to enjoy the benefits from both worlds. Let \( A(\cdot) \) denote the attack algorithm used in the inner loop of adversarial training. In the unified framework, the learner solves the min-max loss of adversarial training over the normalized inputs, i.e.

\[
\min_{f \in H} \max_{A(\cdot)} \sum_{(x, y) \in D} \ell(f, A(N(x)), y)
\]

(2)

to obtain a model \( f^* \) and predicts \( f^*(N(x)) \) at test-time. In the actual learning algorithm, the learner

- first normalizes all training examples with an \( N \) by equivalence groups, and
- in each iteration, trains over the adversarial variant of the normalized input \((N(x), y)\).

Normalization offers many advantages to the unified approach. First, the unified framework still guarantees robustness for adversarial examples in the same equivalence group as the clean input. Second, normalization significantly reduced the search space of adversarial training. Suppose a program uses \( n \) APIs and each API has \( k \) functionally equivalent substitutes, then the number of variants by API substitution alone grows exponentially to \( k^n \). Normalization removes these variants so that adversarial training can focus on other transformations. Third, normalization can potentially reduce the model capacity needed to learn a robust model as shown in Appendix A.2. In return, adversarial training can deal with transformations for which normalization incurs additional robustness-accuracy trade-off, e.g. against code injection as shown in Sec 4.2.
4.4 Efficient Normalization

Although the problem of exact normalization in its most general form can be computationally hard, we propose a practical framework of generating heuristic normal forms. First, for each equivalence group $E$, we assign an order to all its elements. An ideal normalizer should return the element with the lowest order. Given an input, our heuristic normalizer will iteratively apply the transformation that lowers the order of the input until no such transformations are available. The final form will be the approximate normal form. Taking semantically equivalent source codes in authorship attribution as an example. We can rank code pieces by its syntactical features, e.g. number of for loops. The normalizer, while trying to reduce the order, will apply the for-to-while transformation. To ensure that the approximate normal form still lies in the same equivalence group, every transformation we use in $N$ has a reversible counterpart in $R$, e.g. for-to-while can be negated by while-to-for. Then, $N(x)$ can always be converted back to $x$ by applying the transformations in the reverse direction. The detailed description can be found in Appendix B. The code of our normalizer is also open sourced on github.

5 Experiment

We now evaluate the effectiveness of NaP against relational attacks for real-world attacks. In particular, we seek answers to the following questions.

1. Does NaP deliver the promised robustness improvement under real-world attacks?
2. How much performance edge does NaP provide — both in robust accuracy and computation time — compared to standard adversarial training against relational adversaries?
3. How is the robustness-accuracy trade-off compared to without normalization, and is the trade-off worthwhile?

We investigate these questions over two real-world tasks — authorship attribution and malware detection. Our result shows that normalization in the learning pipeline can significantly boost model robustness against relational adversaries compared to adversarially trained and unprotected models.

5.1 Authorship Attribution

Automated authorship attribution is a classical task in program analysis. Successful identification of the author of a code piece can help catching plagiarism, tracking contributors to shared projects and identifying authors of malicious content [Caliskan-Islam et al., 2015, Abuhamad et al., 2018]. However, Quiring et al. [2019] shows that semantic-preserving transformation over the source code can significantly reduce the performance of ML-based authorship attribution models. Although the attack is hard to defend if the attacker knows the coding style of the imitation target well, we are interested in whether authorship attribution is still possible under common code transformations that does not require any knowledge of the imitation target.

Dataset. We use the dataset provided by Quiring et al. [2019], which is collected from Google Code Jam, a coding platform on which individual programmers compete to solve coding challenges. It consists of 1,632 files of C++ code from 204 authors solving the same 8 programming challenge questions. We follow the same train-test data splits in Quiring et al. [2019]. We create 8 different data splits. Each split uses the codes from one challenge as the test set and the codes from the rest seven challenges for training. We run the experiments over all 8 splits and report the average results.

Relation and Normalization. The attack proposed by Quiring et al. [2019] applies Monte-Carlo tree search to determine the sequence of program transformations for creating adversarial examples. As explained in Quiring et al. [2019], all transformations, 42 in total covering control-flow, variable and function declaration, I/O API, etc., merely change the common, generic program features that do not fundamentally alter the signature of each author, therefore, the generated forgeries should have the same label as the original copies. To keep our engineering workload manageable, we consider 35 transformations that are easier to normalize. Some of the transformations are shown in

\[ \text{https://github.com/Mohannadcse/Normalizer-authorship} \]

\[ \text{https://github.com/EQuiw/code-imitator/tree/master/data/dataset_2017} \]
We consider three attack modes. All three attacks use the Monte-Carlo tree search (MCTS) method in [Quiring et al., 2019] with the same parameters including number of iterations and roll-out per iteration. The Non-Adaptive method attacks without knowing the existence of the normalizer \( \mathcal{N} \), i.e. runs MCTS using prediction scores over the original inputs. The Adaptive attack, in contrast, runs MCTS using the model prediction over the normalized inputs. This means for each roll-out attempt, the attacker will also run the normalizer and using the prediction score over the normalized input to plan its next move. Last, we consider Adaptive+, which runs MCTS over normalized inputs and also attacks with all 42 transformations in [Quiring et al., 2019]. Although no performance is guaranteed against an attacker with a larger adversarial feasible set than in training, we are still curious about how \( \mathbb{N} \mathbb{P} \) compares to the adversarial trained and the unprotected models.

Baseline Models. We consider both the random forest (RF) model and the recurrent neural net model with LSTM units attacked in Quiring et al. [2019] for baseline performance. Hereinafter, we name them Vanilla-RF and Vanilla-LSTM. Our \( \mathbb{N} \mathbb{P} \) framework uses the same model structure, parameters and training procedure as Quiring et al. Hereinafter, we call the normalized models \( \mathbb{N} \mathbb{P} \)-RF and \( \mathbb{N} \mathbb{P} \)-LSTM. For the adversarial training baseline, we train the lone compatible model Vanilla-LSTM into Adv-LSTM as Vanilla-RF, a non-parametric method, has no gradients for adversarially training to exploit [Wang et al., 2018, Yang et al., 2020b]. We note that the standard adversarial training is too computationally expensive for the attack on source code level. We make a number of adaptations that reduce the number of MCTS roll-outs and generate adversarial examples in batch for better parallelism so that the process finishes within a month on a 72-core CPU server. The details of these adaptations can be found in Appendix [B].

\( \mathbb{N} \mathbb{P} \) vs. Vanilla Models. Table 1 shows the test accuracy of all baselines against adversarial examples and clean inputs. \( \mathbb{N} \mathbb{P} \)-LSTM has the highest accuracy against all three attacks followed by \( \mathbb{N} \mathbb{P} \)-RF. Compared to the corresponding vanilla models with the same model structures, the \( \mathbb{N} \mathbb{P} \) models achieves higher accuracy by a wide margin. The accuracy increases by more than 50% for both the non-adaptive and adaptive attack. The results show that \( \mathbb{N} \mathbb{P} \) is highly effective when the attacker uses transformations already considered by the normalizer. Intriguingly, even under the Adaptive+ attack in which some attack transformations are not normalized, \( \mathbb{N} \mathbb{P} \) still achieves nontrivial accuracy – 37.3% for \( \mathbb{N} \mathbb{P} \)-RF and 49.9% for \( \mathbb{N} \mathbb{P} \)-LSTM – which is 36% and 49% higher than Vanilla-RF/Vanilla-LSTM. This is because most of the time the MCTS attack will use the 7 transformations that have not been normalized. \( \mathbb{N} \mathbb{P} \) effectively reduces the attack surface and thus still enhances the accuracy. For clean inputs, the \( \mathbb{N} \mathbb{P} \) models has lower accuracy compared to the vanilla models due to the inevitable robustness-accuracy trade-off analyzed in Sec 4.2. However, the difference (<10% for LSTM and <15% for RF) is much smaller than the accuracy gain in the adversarial setting. The trade-off is worthwhile when a sizable portion of the inputs come from adversarial sources.

\( \mathbb{N} \mathbb{P} \) vs. Adversarial Training. \( \mathbb{N} \mathbb{P} \) models also consistently outperforms the adversarially trained counterparts across all attacks by a significant margin (~40% for Non-Adaptive/Adaptive and up to 24% for Adaptive+). The performance of adversarial training is heavily affected by the strength of the attack used in the training loop: a model adversarially trained with a weak attack in the loop may succumb to a strong attack in test time. Recall that our adversarial training procedure uses an attack with reduced search parameters in order to have reasonable training time. Adv-LSTM has ~60% accuracy against the adaptive attack in the training loop. However, the accuracy drops significantly to...
We introduce two new relational attack algorithms, which are G-Attack Methods. We consider substitution of API calls with functionally equivalent counterparts, i.e., Model and Baselines.

### Dataset

We evaluate the effect of normalization on malware detection using N which is 40x more than the vanilla training time. As Table 2 shows, relational attacks are overwhelmingly effective to detectors that are oblivious to potential transformations. Adversarial examples almost always (>99% FNR) evade the detection. In contrast, adversarial training requires much longer training time. Even with reduced search parameters, adversarial training still takes more than 20 days to finish on the same infrastructure, which is 40x more than the vanilla training time. N&P shows a clear advantage in running time.

#### 5.2 Malware Detection

**Dataset.** We evaluate the effect of normalization on malware detection using Sleipnir, a data set containing Windows binary API usage features of 34,995 malware and 19,696 benign software, extracted from their Portable Executable (PE) files using LIEF [Thomas, 2017]. The dataset was created by Al-Dujaili et al. [2018] and used to evaluate the effectiveness of their adversarial training against API injection attacks. The detection is exclusively based on the API usage of a malware. There are 22,761 unique API calls in the data set, so each PE file is represented by a binary indicator vector \( x \) ∈ \( \{0, 1\}^m \), where \( m = 22,761 \). We sample 19,000 benign PEs and 19,000 malicious PEs to construct the training (60%), validation (20%), and test (20%) sets.

**Relation and Normalization.** Al-Dujaili et al. [2018] considers adding redundant API calls, i.e., (x, z) ∈ R iff z is obtained by flipping some x’s feature values from 0 to 1. On top of API addition, We consider substitution of API calls with functionally equivalent counterparts, i.e., (x, z) ∈ R iff z is obtained by changing some of x’s feature values from 1 to 0 in conjunction with some other feature values changed from 0 to 1. With expert knowledge, we extract nearly 2,000 equivalent API groups described in Appendix C.3. We normalizes API substitutions by condensing the features of equivalent APIs into one whose value indicates if any API in the group is used.

**Attack Methods.** We introduce two new relational attack algorithms, which are GreedyByGroup and GreedyByGrad. GreedyByGroup searches the combination of API usage within each equivalence group that maximizes the test loss, and then combine the adversarial perturbations from all equivalence groups. GreedyByGrad makes a first-order approximation of the change in test loss caused by potential transformations, applies the transformations with top \( m \) approximated increases and repeats this procedure for \( K \) iterations. Their detailed algorithm descriptions are in Appendix C.1.

In model evaluation, we use our attacks together with the rfgsm_k attack in Al-Dujaili et al. [2018] and call an adversarial attack successful if any of the attack algorithms evades the detection.

**Model and Baselines.** We compare four ML-based malware detectors. The Unified detector uses our unified framework in Sec.4.3 we normalize over equivalent API groups, and then adversarially trains over API addition. The Adv-Trained detector is adversarially trained with the best adversarial example generated using GreedyByGrad and rfgsm_k additive attack. We also include the model proposed by Al-Dujaili et al. [2018], which is adversarially trained against only rfgsm_k additive attack, and lastly a Natural model with no defense. We use the same network architecture as Al-Dujaili et al. [2018], a fully-connected neural net with three hidden layers, each with 300 ReLU nodes, to set up a fair comparison. We train each model to minimize the negative log-likelihood loss as Al-Dujaili et al. [2018], a fully-connected neural net with three hidden layers, each with 300 ReLU nodes, to set up a fair comparison. We train each model to minimize the negative log-likelihood loss as Al-Dujaili et al. [2018], a fully-connected neural net with three hidden layers, each with 300 ReLU nodes, to set up a fair comparison. We train each model to minimize the negative log-likelihood loss as Al-Dujaili et al. [2018], a fully-connected neural net with three hidden layers, each with 300 ReLU nodes, to set up a fair comparison. We train each model to minimize the negative log-likelihood loss as Al-Dujaili et al. [2018], a fully-connected neural net with three hidden layers, each with 300 ReLU nodes, to set up a fair comparison. We train each model to minimize the negative log-likelihood loss as Al-Dujaili et al. [2018], a fully-connected neural net with three hidden layers, each with 300 ReLU nodes, to set up a fair comparison.

**Results.** As Table 2 shows, relational attacks are overwhelmingly effective to detectors that are oblivious to potential transformations. Adversarial examples almost always (>99% FNR) evade the naturally trained model, and also evade the detector in Al-Dujaili et al. [2018] most of the time (>89% FNR) as it does not consider API substitution. On the defense end, Unified achieves the highest robust accuracy: the evasion rate (FNR) only increases by 0.5% on average. Adv-Trained comes second
but the evasion rate is $> 20\%$ higher. The evasion is mostly caused by GREEDYBYGROUP, the attack that is too computationally expensive to be included in the training loop. This result corroborates with the theoretical advantage of NaP: its robustness guarantee is independent of training algorithms. Last, all detectors using robust learning techniques have higher FPR compared to Natural, which is expected because of the inevitable robustness-accuracy trade-off. However, the difference is much smaller compared to the cost due to attacks, and thus the trade-off is worthwhile.

6 Conclusion and Future Work

In this work, we set the first step towards robust learning against relational adversaries: we theoretically characterize the conditions for robustness and the sources of robustness-accuracy trade-off, and propose a provably robust learning framework. Our empirical evaluation shows that input normalization can significantly enhance model robustness. For future work, we see automatic detection of semantics-preserving transformation as a promising addition to our current expert knowledge approach, and plan to extend the normalization approach to model explainability, fairness and security problems beyond relational adversaries.

7 Societal Impact

We propose input normalization as a principle approach to enhance ML model robustness against relational adversaries. In most cases, the extra robustness in security-critical tasks will reduce the loss caused by malicious behaviors and thus bring positive societal impact. We do acknowledge that one of our specific empirical evaluation setting — automated authorship attribution — can cause privacy issues when used for censorship; an enhanced authorship attribution technique may allow the censoring agent to better identify the author. We note that NaP in our experiment only normalizes over the most basic set of program features to investigate ML models’ stability in common use cases. For more privacy-sensitive cases, an author can still use more advanced techniques, such as randomize variable/function names, encrypt its code or even normalize its code with a more comprehensive relation before code submission, to remove the idiosyncrasies in its codes. Our NaP framework over generic program features will not conflict with these anonymization techniques.
References


**Checklist**

1. For all authors...
   (a) Do the main claims made in the abstract and introduction accurately reflect the paper’s contributions and scope?
      A: Yes.
   (b) Have you read the ethics review guidelines and ensured that your paper conforms to them?
      A: Yes.
   (c) Did you discuss any potential negative societal impacts of your work?
      A: Yes, this is done in Section 7, Page 9.
   (d) Did you describe the limitations of your work?
      A: Yes. In fact, a main contribution of our paper is to discuss the robustness-accuracy trade-off caused by input normalization, i.e. the fundamental limitation of learning against relational adversaries.

2. If you are including theoretical results...
   (a) Did you state the full set of assumptions of all theoretical results?
      A: Yes.
   (b) Did you include complete proofs of all theoretical results?
      A: Yes. The full proofs and the formal version of the statements can all be found in Appendix A.

3. If you ran experiments...
   (a) Did you include the code, data, and instructions needed to reproduce the main experimental results (either in the supplemental material or as a URL)?
      A: The normalizer in our paper is now open sourced on [https://github.com/Mohannadcse/Normalizer-authorship](https://github.com/Mohannadcse/Normalizer-authorship).
   (b) Did you specify all the training details (e.g., data splits, hyperparameters, how they were chosen)?
      A: Yes.
   (c) Did you report error bars (e.g., with respect to the random seed after running experiments multiple times)?
      A: Yes.
   (d) Did you include the amount of compute and the type of resources used (e.g., type of GPUs, internal cluster, or cloud provider)?
      A: Yes.

4. If you are using existing assets (e.g., code, data, models) or curating/releasing new assets...
   (a) If your work uses existing assets, did you cite the creators?
      A: Yes.
(b) Did you mention the license of the assets?
A: Yes. The source code used in [Al-Dujaili et al., 2018] is under MIT license. The repo can be found at [https://github.com/ALFA-group/robust-adv-malware-detection](https://github.com/ALFA-group/robust-adv-malware-detection) We directly requested the authors to obtain the Sleipnir dataset. The dataset and code used in [Quiring et al., 2019] is under GPL-3.0 license. We got the code from its github page at [https://github.com/EQuiw/code-imitator](https://github.com/EQuiw/code-imitator).

(c) Did you include any new assets either in the supplemental material or as a URL?
A: At this moment, we are waiting for the internal legal inspection over proprietary information. We can release the code as soon as the process is complete.

(d) Did you discuss whether and how consent was obtained from people whose data you’re using/curating?
A: Yes. We obtain the code directly from the collector of the datasets. The datasets contain no personal biometric information.

5. If you used crowdsourcing or conducted research with human subjects...
A: No, we did not.
A  Proofs and Explanation for Theoretical Results

In this section, we present the omitted proofs for theorems and observations due to page limit of the main body.

A.1  Proof for Computational Hardness of Adversarial Training

We first write the full statement of hardness in the following theorem.

**Theorem 2.** Let \( \mathcal{R} \subseteq \{0, 1\}^d \times \{0, 1\}^d \) be a relation. Given a function \( f \), an input \( \mathbf{x} \in \{0, 1\}^d \) and a feasible set \( T(\mathbf{x}) = \{ \mathbf{z} : \mathbf{x} \rightarrow_{\mathcal{R}} \mathbf{z} \} \), solving the following maximization problem:

\[
\max_{\mathbf{z} \in T(\mathbf{x})} l(f, \mathbf{z}, y)
\]

is PSPACE-hard when \( l(f, \mathbf{x}, y) \) is the 0-1 classification loss.

**Proof.** Let \( \alpha : \{0, 1\}^d \rightarrow \{0, 1\} \) be a predicate. Define a loss function \( l(f, \mathbf{z}, y) \) as follows: \( l(f, \mathbf{z}, y) = \alpha(\mathbf{z}) \) (the loss function is essentially the value of the predicate). Note that \( \max_{\mathbf{x} \in T(\mathbf{x})} l(f, \mathbf{z}, y) \) is equal to 1 iff there exists a \( \mathbf{z} \in T(\mathbf{x}) \) such that \( \alpha(\mathbf{z}) = 1 \). This is a well-known problem in model checking called reachability analysis, which is well known to be PSPACE-complete (the reduction is from the problem of checking emptiness for a set of DFAs, which is known to be PSPACE-complete [Kozlenko, 1977]). \( \square \)

Recall that the maximization problem \( \max_{\mathbf{x} \in B_{\epsilon}(x, r)} l(f, \mathbf{z}, y) \) used in adversarial training for the image modality was proven to be NP-hard [Katz et al., 2017]. Hence it seems that the robust optimization problem in our context is in a higher complexity class than in the image domain.

A.2  Model Capacity Requirement

In this section, we illustrate how the \( NA \)P framework can potentially help reduce the model complexity for learning a robustly accurate classifier. We start with the following proposition.

**Proposition 2** (Model Capacity Requirement). For some hypothesis class \( \mathcal{H} \) and relation \( \mathcal{R} \), there exists \( f \in \mathcal{H} \) such that \( f(\mathcal{N}_\mathcal{R}(\cdot)) \) is robustly accurate, but no \( f \in \mathcal{H} \) can be robustly accurate on the original inputs. In other words, robustly accurate classifier can only be obtained after normalization.

We first define an equivalence relation induced by equivalent coordinates over binary inputs, and then write the formal statement of the observation in the following claim.

**Definition 5** (Equivalence relation induced by equivalent coordinates). Let \( \mathbf{x} = (x_1, \ldots, x_d) \) be a binary input vector on \( \{0, 1\}^d \), where each \( x_i, i \in \{1, \ldots, d\} \) is a coordinate. Let \( I = \{1, \ldots, d\} \) be the set of coordinate indices for inputs in \( \mathcal{X} \) and \( U = \{i_1, \ldots, i_m\} \subseteq I \). In an equivalence relation \( \mathcal{R} \) induced by \( U \), \( \mathbf{x} \rightarrow_{\mathcal{R}} \mathbf{z} \) iff 1) \( x_i = z_i \) for all \( i \in I \setminus U \), and 2) \( \bigvee_{i \in U} x_i = \bigvee_{i \in U} z_i \). Notice that \( \mathbf{x} \rightarrow_{\mathcal{R}} \mathbf{z} \) iff \( \mathbf{z} \rightarrow_{\mathcal{R}} \mathbf{x} \).

The notation \( \bigvee_{i \in U} x_i \) means taking a logic or operation over all \( x_i \)s for \( i \in U \). Intuitively, having an equivalence relation induced by coordinates with indices in \( U \) means the presence of any combination of such coordinates is equivalent to any other combination. Usage of interchangeable APIs in malware implementation is an example of equivalence relation: the attacker can choose any combination from a set of equivalent APIs to implement the same functionality.

In Definition 5, we use \( U \) to represent the set of indices of the equivalent coordinates. In the following theorems and proofs, we overload \( U \) to also represent the set of equivalent coordinates directly, and the \( \bigvee \) operation will be taken over all coordinates in \( U \).

**Claim 1.** Consider \( \mathcal{X} = \{0, 1\}^5 \) and \( \mathcal{Y} = \{0, 1\} \). Let the coordinates of an input \( \mathbf{x} \in \mathcal{X} \) be \( \{x_1, x_1', x_2, x_3, x_4\} \). Suppose we have an equivalence relation induced by \( U = \{x_1, x_1'\} \). Meanwhile, the true label \( y \) of an input \( \mathbf{x} \) is 1 iff any of the following clauses is true: 1) \( (x_2 = 1) \land (x_3 = 1) \), 2) \( (x_1 \lor x_1' = 1) \land (x_2 = 1) \), 3) \( (x_1 \lor x_1' = 1) \land (x_3 = 1) \land (x_4 = 1) \). Then

1. no linear model can classify the inputs with perfect robust accuracy, but
2. a robust and accurate linear model exists under normalize-and-predict.
Proof. Let $\mathcal{H} = \{ f_{w,b} : \text{sgn}(\langle w, x \rangle + b) \}$. Let $w_1, w_1', w_2, w_3, w_4$ denote the coordinates in $w$ that corresponds to $x_1, x_1', x_2, x_3, x_4$.

We know $y = 1$ if $x_1 = 1, x_2 = 1$ and the other coordinates are zero because the second clause in the labeling rule is satisfied. Therefore, in order to classify this input instance correctly, we must have $w_2 + w_1 + b > 0$. Since $x_1$ and $x_1'$ are equivalent, we should also have $w_2 + w_1' + b > 0$.

Similarly, we know $y = 0$ if $x_3 = 1, x_4 = 1$ and the other coordinates are zero because none of the clauses are satisfied. Therefore, we must have $w_3 + w_4 + b < 0$.

In order to classify all possible $x$ correctly, the classifier $f_{w,b}$ must satisfy

\begin{align*}
w_2 + w_4 + b &< 0 \quad (3) \\
w_2 + w_1 + b &> 0 \quad (4) \\
w_2 + w_1' + b &> 0 \quad (5) \\
w_3 + w_1 + w_1' + b &< 0 \quad (6) \\
w_3 + w_4 + w_1' + b &> 0 \quad (7)
\end{align*}

First, by Formula 4 and 5, we have $w_1 > w_4$ and $w_1' > w_4$. However, by Formula 6 and 7, we have $w_1 + w_1' < w_4 + w_1$, which implies $w_1' < w_4$. Contradiction. Therefore, no linear classifier can satisfy all the equations.

On the other hand, if we perform normalization by letting $x_1 = x_1 \lor x_1'$ and removing $x_1'$, then a classifier $f_{w,b}$ with $w_1 = 0.4, w_2 = 0.7, w_3 = 0.5, w_4 = 0.2, b = -1$ can perfectly classify $x$.

A.3 Proof for Theorem

Proof. We prove the statement by contradiction: suppose all classifiers $f \in \mathcal{H}$ with the highest robust accuracy has $f(x) \neq f(z)$ for some $x, z$ in the same equivalence group $\mathcal{E}$, then we must be able to find a $f'$ such that $f'$ is at least as robust accurate as $f$ and $f'(x) = f'(z)$.

We construct a classifier $f'$ as follows: 1) $f'(x) = f'(z)$, and 2) $f'(z') = f(z')$ for all $z' \in \mathcal{X}, z' \neq z$, i.e. $f'$ and $f$ agree at all inputs except at $z$. We discuss the robust accuracy of $f$ and $f'$ in the following two complementary cases.

Case 1. If $x$ and $z$ are not in the adversarial feasible set $\mathcal{T}(x')$ of an input $x'$, then $f$ is robust accurate at $x'$ iff $f'$ is robust accurate at $x'$. This is obvious as $f(x'') = f'(x'')$ for all $x'' \in \mathcal{T}(x')$.

Case 2. If $x$ or $z$ is in $\mathcal{T}(x')$, then both $x$ and $z$ must be in $\mathcal{T}(x')$. Suppose $x \in \mathcal{T}(x')$, then by the definition of relational adversary, we have $x' \rightarrow^* \mathcal{R} x$ and by the definition of equivalence group, we have $x \rightarrow^*_\mathcal{R} z$. Therefore, we must have $x' \rightarrow^* \mathcal{R} z$, i.e., $z \in \mathcal{T}(x')$. Similarly, $x$ must be in $\mathcal{T}(x')$ if $z$ is in $\mathcal{T}(x')$. In this case, $f$ cannot be robust accurate at $x'$: both $x$ and $z$ are in $\mathcal{T}(x')$, and $f(x) \neq f(z)$; if $f(x') = f(x)$, then $z$ will become a successful adversarial example; similarly, if $f(x') = f(z)$, then $x$ will be the adversarial example. Since $f$ is already not robust accurate at $x'$, the robust accuracy of $f'$ cannot be lower than $f$ in this case.

Combining the two cases, we can conclude that $f'$ is at least as robust accurate as $f$ over the underlying data distribution $\mathcal{D}$. The existence of $q$ naturally follows the definition of normalizer by equivalence group: the normalizer $N$ guarantees the consistency of prediction within each $\mathcal{E}$, and $q$ just need to match $f$'s prediction on all normalized inputs. \qed
B Authorship Attribution – Algorithms and Implementation Details

B.1 Implementation and Infrastructure

We implement source code normalizer on top of Clang\cite{clang2016}, an open-source C/C++ frontend for the LLVM compiler framework. For fair comparison of running time, we run all experiment series on an Amazon EC2 c5.18xlarge instance with 72 cores and 144GB memory. We train the model using the \texttt{sklearn} and \texttt{keras} APIs with \texttt{tensorflow} backend.

B.2 Attack Transformations to be Normalized

The evasion and impersonation attack in Quiring et al.\cite{Quiring2019} uses a total of 42 transformation options. These transformations can be divided into two categories based on the information needed. The general transformations, such as changes in control statement, can be applied to any code without prior knowledge of the target author. In contrast, the template-based transformations, which aim to mimic the target author’s function/variable naming and type-def habits, will require samples of the target author’s code pieces. We consider 35 of the 42 transformations as listed in Table 3\footnote{We follow the counting method in Quiring et al.\cite{Quiring2019} and count by the number of options instead of number of transformers. For example, the input interface transformer has two options – stdin and file; therefore, we count the input interface transformer as having two transformation options.}.

The selection covers most transformations including API usage, variable declaration, I/O style and control statement. The transformations are also common across attack papers Quiring et al.\cite{Quiring2019}, Liu et al.\cite{Liu2021}. The only general transformation not considered by us is function in-lining, which removes declared functions in the code and move the commands in-line to the caller function. While this may be done for simple online judge submissions, the boundary of usage of a function is in general hard to track in large code projects. In-lining function in one script may raise problems in other scripts that import and call the function. We do not normalize the transformation as it is hardly semantics-preserving Stucki et al.\cite{Stucki2020}. The template-based attack transformations not considered by our normalizer is function/variable renaming. The function/variable transformer peeks into the codes written by the target of imitation, and then rename the function and variables to match the target author’s habit. This attack action requires extensive knowledge to the target author’s coding habit. In the extreme case, the attacker may be more familiar to the target author than the learner, and thus makes the evasion inevitable. We do not normalize this transformation as it can significantly obscure the ground truth.
<table>
<thead>
<tr>
<th>Transformer</th>
<th>Family</th>
<th>Transformer Description</th>
<th>Syntactic Feature of Interests (SFoI)</th>
<th>Normalizing Action in $\mathcal{T}_N$</th>
</tr>
</thead>
<tbody>
<tr>
<td>For statement transformer</td>
<td>Control</td>
<td>Replaces a for-statement by an equivalent while-statement</td>
<td>No. of for loops</td>
<td>Transform for loops to while</td>
</tr>
<tr>
<td>While statement transformer</td>
<td></td>
<td>Replaces a while-statement by an equivalent for-statement</td>
<td></td>
<td></td>
</tr>
<tr>
<td>If statement transformer</td>
<td></td>
<td>Split the condition of a single if-statement at logical operands (e.g., &amp;&amp; or</td>
<td></td>
<td>) to create a cascade or a sequence of two if-statements depending on the logical operand</td>
</tr>
<tr>
<td>Array transformer</td>
<td>Declaration</td>
<td>Converts a static or dynamically allocated array into a C++ vector object</td>
<td>No. of array that can potentially be converted to C++ vector object</td>
<td>Convert ALL static or dynamically allocated array into a C++ vector object.</td>
</tr>
<tr>
<td>String transformer</td>
<td></td>
<td>Array option: Converts a char array (C-style string) into a C++ string object. The transformer adapts all usages in the respective scope, for instance, it replaces all calls to strlen by calling the instance methods size. String option: Converts a C++ string object into a char array (C-style string). The transformer adapts all usages in the respective scope, for instance, it deletes all calls to cstr()</td>
<td>No. of char array</td>
<td>Convert all string to C++ string object.</td>
</tr>
<tr>
<td>Integral type transformer</td>
<td></td>
<td>Promotes integral types (char, short, int, long, long long) to the next higher type, e.g., int is replaced by long.</td>
<td>Number of integral type declaration lower than long long</td>
<td>Replace all integral type with long long.</td>
</tr>
<tr>
<td>Floating-point type transformer</td>
<td></td>
<td>Converts float to double as next higher type.</td>
<td>No. of float type declaration lower than double</td>
<td>Convert all float to double.</td>
</tr>
<tr>
<td>Boolean transformer</td>
<td></td>
<td>Bool option: Converts true or false by an integer representation to exploit the implicit casting. Int option: Converts an integer type into a boolean type if the integer is used as boolean value only</td>
<td>No. of Boolean values</td>
<td>Use int representation in all cases</td>
</tr>
<tr>
<td>Init-Decl transformer</td>
<td></td>
<td>Move into option: Moves a declaration for a control statement if defined outside into the control statement. For instance, int i; ;; for(i = 0; i &lt; N; i++) becomes for(int i = 0; i &lt; N; i++). Move out option: Moves the declaration of a control statement’s initialization variable out of the control statement.</td>
<td>No. of looping variable declared outside the control statement</td>
<td>Use move-in option for all scenarios.</td>
</tr>
<tr>
<td>Typedef transformer</td>
<td></td>
<td>Convert option: Convert a type from source file to a new type via typedef, and adapt all locations where the new type can be used. Delete option: Deletes a type definition (typedef) and replace all usages by the original data type.</td>
<td>No. of user-defined types</td>
<td>Apply the delete option to all typedef.</td>
</tr>
<tr>
<td>Transformer</td>
<td>Family</td>
<td>Transformer Description</td>
<td>Syntactic Feature of Interests (SFoI)</td>
<td>Normalizing Action in $T_N$</td>
</tr>
<tr>
<td>--------------------------</td>
<td>-----------------</td>
<td>----------------------------------------------------------------------------------------</td>
<td>--------------------------------------</td>
<td>----------------------------</td>
</tr>
<tr>
<td>Include typedef transformer</td>
<td>Template</td>
<td>Inserts a type using typedef, and updates all locations where the new type can be used. Defaults are extracted from the 2016 Code Jam Competition.</td>
<td>No. of user-defined types</td>
<td>Apply the delete option to all typedef.</td>
</tr>
<tr>
<td>Unused code transformer</td>
<td>Declaration</td>
<td>Function option: Removes functions that are never called. Variable option: Removes global variables that are never used.</td>
<td>No. of user-defined types &amp; functions</td>
<td>Remove unused variables and functions.</td>
</tr>
<tr>
<td>Input interface transformer</td>
<td>API</td>
<td>Stdin option: Instead of reading the input from a file (e.g. by using the API ifstream or freopen), the input to the program is read from stdin directly (e.g. cin or scanf). File option: Instead of reading the input from stdin, the input is retrieved from a file.</td>
<td>No. of file I/O</td>
<td>Use stdin always.</td>
</tr>
<tr>
<td>Output interface transformer</td>
<td>API</td>
<td>Stdout option: Instead of printing the output to a file (e.g. by ofstream or freopen), the output is written directly to stdout (e.g. cout or printf). File option: Instead of writing the output directly to stdout, the output is written to a file.</td>
<td>No. of file I/O</td>
<td>Use stdout always.</td>
</tr>
<tr>
<td>Input API transformer</td>
<td></td>
<td>C++-Style option: Substitutes C APIs used for reading input (e.g., scanf) by C++ APIs (e.g., usage of cin). C-Style option: Substitutes C++ APIs used for reading input (e.g., usage of cin) by C APIs (e.g., scanf).</td>
<td>No. of C-style I/O API</td>
<td>Use C++ API always.</td>
</tr>
<tr>
<td>Output API transformer</td>
<td></td>
<td>C++-Style option: Substitutes C APIs used for writing output (e.g., printf) by C++ APIs (e.g., usage of cout). C-Style option: Substitutes C++ APIs used for writing output (e.g., usage cout) by C APIs (e.g., printf).</td>
<td>No. of C++ Style always.</td>
<td>Use C++ Style always.</td>
</tr>
<tr>
<td>Sync-with_stdio transformer</td>
<td></td>
<td>Enable or remove the synchronization of C++ streams and C streams if possible.</td>
<td>No. of potential synchronization sites</td>
<td>Enable all synchronization.</td>
</tr>
<tr>
<td>Compound statement transformer</td>
<td>Misc</td>
<td>Insert option: Adds a compound statement ({...}). The transformer adds a new compound statement to a control statement (if, while, etc.) given their body is not already wrapped in a compound statement. Delete option: Deletes a compound statement ({...}). The transformer deletes compound statements that have no effect, i.e., compound statements containing only a single statement.</td>
<td>No. of compound statements</td>
<td>Apply the delete option to all locations.</td>
</tr>
<tr>
<td>Return statement transformer</td>
<td></td>
<td>Adds a return statement. The transformer adds a return statement to the main function to explicitly return 0 (meaning success). Note that main is a non-void function and is required to return an exit code. If the execution reaches the end of main without encountering a return statement, zero is returned implicitly.</td>
<td>No. of implicit return sites</td>
<td>Adds the return statement to all applicable locations.</td>
</tr>
</tbody>
</table>
We identify the Syntactic Features of Interests (SFoI) corresponding to the transformers as shown in Table 3. Each SFoI’s value can be changed by one or multiple transformers. For example, we use the number of for loops in the code as one SFoI, and its value can be changed by two transformers – the for-to-while and the while-to-for transformer over the control flow.

We take the order function $\lambda$ as the sum of values of the SFoIs, and the normal form of a code piece is the variant that has the smallest sum of SFoIs. For example, if the SFoIs are 1) number of for loops and 2) number of printf statements, then the normal form will be the variant that has the least number of for loops and printf statements in total. In Table 3, we identify a total of 16 SFoIs and use the sum of all of them as our order function $\lambda$.

We select a subset of the attack transformations as the set of normalizing transformations $\mathcal{T}_N$ so as to respect the equivalence groups. Our $\mathcal{T}_N$ contains and only contains the attack transformations that strictly decrease the value of SFoIs. For example, when the number of for loops is used as an SFoI, we keep the for-to-while transformer in $\mathcal{T}_N$ and discard the while-to-for. The right-most column of Table 3 shows all the transformations we keep in $\mathcal{T}_N$. Notice that the SFoIs all take non-negative integer values, and the normalizing transformations reduce the value of $\lambda$ by 1 once applied. Therefore, the number of iterations in normalization for an input $x$ is bounded by the value of $\lambda(x)$.

**B.4 Normalization in Action**

Figure 3 shows an example of our normalizer in action. The left-most box contains an code snippet originally written by Author A. The subsequent code boxes in the top row illustrate a sequence of transformations applied to the original code. The attacker first converts the C-style printf statement to the C++ style cout statement, then changes the for loop to a while loop, and eventually changes the value of a boolean variable from 0 to False. The final variant in right-most box is a successful adversarial example that misled the model to predict a different author.

In this code example, three syntactical features of interests are involved: 1) the number of for loops, 2) the number of C-style I/O statements and 3) the number of Boolean values that can be cast into integers. The normalizer applies the normalizing actions in a iterative manner, reducing the number of for loops, C-style IO statements and Boolean values until no more action is applicable. All four code pieces – the original input, the final adversarial example and the two intermediate variants – will be normalized into the same normal form as depicted in the bottom box in Figure 3.

**B.5 Adaptation for Adversarial Training**

In existing adversarial training literature, the $\ell_p$-norm based adversarial examples are created directly in the feature-space using gradient ascent; the attack can be readily computed in GPU in a similar manner as model updates. The code transformations, however, are performed in the problem-space; the MCTS computation is CPU-intensive and thus takes much longer. In addition, the validity check of adversarial examples further increases the computational load. We make the following adjustments to speed up the adversarial training process. First, we use Vanilla-LSTM as the initial model and fine-tune it using adversarial inputs. The models show improved robust accuracy (~60%) to the attacks in the training loops after 10 epochs. Second, instead of generating adversarial examples at every training step, we generate the adversarial training inputs for all .cpp files with respect to the model at beginning of an epoch. This change allows us to generate the adversarial training inputs in parallel. To further speed up the adversarial training procedure, we also reduce the max-depth from 25 to 5 as well as the number of random play-outs at each node from 50 to 10 in the Monte-Carlo tree-search. With these modification, we finally manage to finish adversarial training within a month.
C Malware Detection – Algorithms and Implementation Details

In this section, we present the omitted algorithm descriptions and experiment implementation details for the malware detection experiment.

C.1 Generic Relational Attack Algorithms

In Sec 5, we introduce two generic relational attack algorithms – GreedyByGroup and GreedyByGrad. The algorithm boxes below shows the exact description of both algorithms.

**Algorithm 1 GreedyByGroup**

\( x^{adv} = x, k = 0 \)

Partition \( \mathcal{R} \) into \( m \) groups \( \{\mathcal{R}_1, \ldots, \mathcal{R}_m\} \).

while \( k < K \) do

\( k = k + 1 \)

for \( \mathcal{R}_i \in \{\mathcal{R}_1, \ldots, \mathcal{R}_m\} \) do

\( x_i = \arg \max_{z: x^{adv} \rightarrow \mathcal{R}_i} \ell(f, z, y) \).

end

Combine \( x_i \)’s to obtain the new \( x^{adv} \).

end

return \( x^{adv} \)

**Algorithm 2 GreedyByGrad**

\( x^{adv} = x, k = 0 \)

while \( k < K \) do

\( k = k + 1 \)

\( g = \nabla_x \ell(f, x^{adv}, y) \)

for \( (x^{adv}, z) \in \mathcal{R} \) do

\( c(x^{adv}, z) = \sum_{i=1}^d g_i(z_i - x^{adv}_i) \)

end

Apply the transformations with top \( m \) largest positive \( c(x^{adv}, z) \) to obtain the new \( x^{adv} \).

end

return \( x^{adv} \)

C.2 GreedyByGrad and GreedyByGroup for Malware Detection on Sleipnir

We instantiate the two attacks on our malware detection task as follows. For GreedyByGroup, we divide the relation by equivalent API groups. In each iteration, the attacker searches the best combination of API in each equivalence group that causes the most increase in test loss. These combinations are concatenated using a logical OR operation. For GreedyByGrad, the attacker tries all single API substitutions and additions allowed by the relation, and use the first-order approximation of test loss to determine the \( m \) transformations to be applied in each iteration. We run both attacks for various \( K \) and \( m \). We find that \( K = 10 \) and \( m = 10 \) suffices to bypass detection for the vanilla model.

C.3 Extract API Substitution Rules

In Sec 5, we consider malware authors who can substitute API calls with equivalent API calls to evade ML-based malware detector. We now explain how we extract the equivalent APIs. We identify four types of patterns for extracting equivalent APIs:

- API with the same name but located in different Dynamically Linkable Libraries (DLLs). For example, memcpy, a standard C library function, is shipped in libraries with different names, including crrd11.dll, msvcr90.dll, and msvcr110.dll.
- API with and without the Ex suffix. The Ex suffix represents an extension to the same API without the suffix.
- API with and without the A or W suffixes. The A suffix represents the single character version. The W suffix represents the wide character version.
- API with/without _s suffix. The _s suffix represents the secure version of an API.

Using these four patterns, we extracted about 2,000 equivalent API groups. About 500 of the groups have more than 2 APIs and the maximal group has 23 APIs.