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# Supplement to Robust Portfolio Optimization

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## A Matrix Projection

In this section, we summarize the algorithm proposed in [1] for solving the matrix projection problem (3.3). Let

$$\Omega_1 := \left\{ \mathbf{x} = \text{vec}(\mathbf{X}) : \mathbf{X} \in S_\lambda \right\}$$

$$\Omega_2 := \left\{ \mathbf{z} = \text{vec}(\mathbf{Z}) : \mathbf{Z} \in \mathbb{R}^{d \times d}, \mathbf{Z} = \mathbf{Z}^\top, \sum_{i,j=1}^d |Z_{ij}| \leq 1 \right\}.$$

For any symmetric matrix  $\mathbf{V} \in \mathbb{R}^{d \times d}$  and  $\mathbf{v} = \text{vec}(\mathbf{V})$ , define the projection of  $\mathbf{v}$  onto  $\Omega_i$  as

$$P_{\Omega_i}(\mathbf{v}) = \arg \min_{\mathbf{x} \in \Omega_i} \|\mathbf{x} - \mathbf{v}\|_2^2, \quad (\text{A.1})$$

for  $i = 1, 2$ . The algorithm for solving (3.3) builds on solutions to the problems in (A.1). Solving for  $P_{\Omega_1}(\mathbf{v})$  is straightforward. It's well known that

$$P_{\Omega_1}(\mathbf{v}) = \text{vec}(\mathbf{U}\tilde{\mathbf{\Lambda}}\mathbf{U}^\top), \quad (\text{A.2})$$

where  $\mathbf{V} = \mathbf{U}\mathbf{\Lambda}\mathbf{U}^\top$  is a spectral decomposition of  $\mathbf{V}$ ,  $\tilde{\mathbf{\Lambda}} = \text{diag}(\tilde{\Lambda}_{11}, \dots, \tilde{\Lambda}_{dd})$  and  $\tilde{\Lambda}_{ii} = \min\{\max\{\Lambda_{ii}, \lambda_{\min}\}, \lambda_{\max}\}$  for  $i = 1, \dots, d$ .

Next we solve for  $P_{\Omega_2}(\mathbf{v})$ . Let  $\text{sign}(\mathbf{v}) = \{\text{sign}(v_1), \dots, \text{sign}(v_d)\}^\top$  be a vector of the signs of  $\mathbf{v}$ 's entries. Denote  $|\mathbf{v}| = \text{sign}(\mathbf{v}) \circ \mathbf{v}$  and  $\tilde{\mathbf{v}} = T_{|\mathbf{v}|}(|\mathbf{v}|)$ , where  $T_{|\mathbf{v}|}$  is a permutation transformation that sorts the elements of  $|\mathbf{v}|$  in descending order. Now, if  $\mathbf{1}^\top \tilde{\mathbf{v}} \leq 1$ , we set  $(\tilde{\mathbf{x}}, \tilde{\mathbf{y}}) = (\tilde{\mathbf{v}}, 0)$ . If  $\mathbf{1}^\top \tilde{\mathbf{v}} > 1$ , let  $\Delta \mathbf{v} := (\tilde{v}_1 - \tilde{v}_2, \dots, \tilde{v}_{d-1} - \tilde{v}_d, \tilde{v}_d)^\top \in \mathbb{R}^d$ . Note that  $\Delta v_i \geq 0$  for  $i = 1, \dots, d$  and  $\sum_{i=1}^d i \Delta v_i = \mathbf{1}^\top \tilde{\mathbf{v}} > 1$ . Thus, there exists a smallest integer  $K$  such that  $\sum_{i=1}^K i \Delta v_i \geq 1$ . In this case, we set

$$\tilde{\mathbf{y}} = \frac{1}{K} \left( \sum_{i=1}^K \tilde{v}_i - 1 \right) \text{ and } \tilde{\mathbf{x}} = (\tilde{v}_1 - \tilde{\mathbf{y}}, \dots, \tilde{v}_K - \tilde{\mathbf{y}}, 0, \dots, 0)^\top \in \mathbb{R}^d.$$

Now we can express  $P_{\Omega_2}(\mathbf{v})$  as

$$P_{\Omega_2}(\mathbf{v}) = \text{sign}(\mathbf{v}) \circ T_{|\mathbf{v}|}^{-1}(\tilde{\mathbf{x}}). \quad (\text{A.3})$$

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**Algorithm 1** Solving matrix projection problem (3.3)

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 $\tilde{\mathbf{R}}^Q \leftarrow \text{MatrixProjection}(\hat{\mathbf{R}}^Q, \lambda_{\min}, \lambda_{\max}, \mathbf{x}^0, \mathbf{z}^0, \gamma, \epsilon, N)$ 
 $\mathbf{r} \leftarrow \text{vec}(\hat{\mathbf{R}}^Q)$ 
for  $k = 0, \dots, N$  do
   $\mathbf{e}_x^k \leftarrow \mathbf{x}^k - P_{\Omega_1}(\mathbf{x}^k - \mathbf{z}^k)$ 
   $\mathbf{e}_z^k \leftarrow \mathbf{z}^k - P_{\Omega_2}(\mathbf{z}^k + \mathbf{x}^k - \mathbf{r})$ 
   $\mathbf{e}_k \leftarrow (\mathbf{e}_x^k, \mathbf{e}_z^k)^\top$ 
  if  $\|\mathbf{e}_k\|_{\max} < \epsilon$ , then
    break
  else
     $\mathbf{x}^{k+1} \leftarrow \mathbf{x}^k - \gamma(\mathbf{e}_x^k - \mathbf{e}_z^k)/2$ 
     $\mathbf{z}^{k+1} \leftarrow \mathbf{z}^k - \gamma(\mathbf{e}_x^k + \mathbf{e}_z^k)/2$ 
  end if
end for
return  $\tilde{\mathbf{R}}^Q = \text{mat}(\mathbf{x}^k)$ 
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Next we solve the matrix projection problem in (3.3). Recall that  $\hat{\mathbf{R}}^Q$  is the matrix to be projected to  $S_\lambda$ . Since for any vector  $\mathbf{y} \in \mathbb{R}^d$ , we have  $\|\mathbf{y}\|_{\max} = \max_{\mathbf{c} \in \mathbb{R}^d, \|\mathbf{c}\|_1 \leq 1} \mathbf{c}^\top \mathbf{y}$ , it follows that problem (3.3) can be reformulated as the following mini-max problem:

$$\min_{\mathbf{x} \in \Omega_1} \max_{\mathbf{z} \in \Omega_2} \mathbf{z}^\top \left\{ \mathbf{x} - \text{vec}(\hat{\mathbf{R}}^Q) \right\}. \quad (\text{A.4})$$

If  $(\mathbf{x}^{\text{opt}}, \mathbf{z}^{\text{opt}})$  is a solution to problem (A.4), then  $\text{mat}(\mathbf{x}^{\text{opt}})$  is a solution to problem (3.3). Algorithm 1 gives the pseudo code for solving problem (A.4), and thus (3.3). Recall that  $0 \leq \lambda_{\min} < \lambda_{\max} \leq \infty$  are the lower and upper bounds of the eigenvalues of the projection.  $\mathbf{x}^0 \in \Omega_1$  and  $\mathbf{z}^0 \in \Omega_2$  are arbitrary initial points.  $\gamma \in (0, 2)$  is a parameter controlling the step lengths of every iteration.  $\epsilon > 0$  is a prespecified tolerance level.  $N \in \mathbb{N}$  is the maximum number of iterations desired. The convergence of Algorithm 1 is guaranteed by the following theorem.

**Theorem A.1** ([1]). *Let  $\mathbf{u}^{\text{opt}} := (\mathbf{x}^{\text{opt}}, \mathbf{z}^{\text{opt}})$  be a solution to (A.4). Denote  $\mathbf{u}^k := (\mathbf{x}^{k^\top}, \mathbf{z}^{k^\top})^\top$  and  $\mathbf{e}_u^k := (\mathbf{e}_x^{k^\top}, \mathbf{e}_z^{k^\top})^\top$ . Then Algorithm 1 produces a sequence  $\{\mathbf{u}^k\}$  satisfying*

$$\|\mathbf{u}^{k+1} - \mathbf{u}^{\text{opt}}\|^2 \leq \|\mathbf{u}^k - \mathbf{u}^{\text{opt}}\|^2 + \frac{\gamma(2-\gamma)}{2} \|\mathbf{e}_u^k\|^2.$$

## References

- [1] M. H. Xu and H. Shao. Solving the matrix nearness problem in the maximum norm by applying a projection and contraction method. *Advances in Operations Research*, 2012:1–15, 2012.