

# Supplementary Materials for Complexity of Inference in Latent Dirichlet Allocation

## A Proof of Lemma 2

*Proof of Lemma 2.* Assume there are  $T$  sets each having  $k \geq 3$  elements, and let  $\Phi$  be the optimal LDA objective. Define  $F(n) = \log \Gamma(n + \alpha)$ . Since  $l_{it}$  is constant across all topics, the linear term in Eq. 2 will be a constant  $K$ . First, note that, if there is a perfect matching,

$$\Phi \geq \frac{n}{k} F(k) + \left(T - \frac{n}{k}\right) F(0) + K. \quad (16)$$

The  $F(0)$  term is the contribution of unused topics. Otherwise, assume that the best packing has  $\gamma \leq cn/k$  sets, each with  $k$  elements. Then, by the properties of the log-gamma function,

$$\Phi \leq \gamma F(k) + \frac{n - \gamma k}{k - 1} F(k - 1) + \left(T - \frac{n}{k}\right) F(0) + K, \quad (17)$$

where we assume, conservatively, that all of the remaining words are explained by topics assigned  $(k - 1)$  words. Also, since there was no perfect matching, there were at most  $T - \frac{n}{k}$  unused topics. Using our bound on  $\gamma$ , we have

$$\Phi \leq \frac{cn}{k} F(k) + \frac{n - \frac{cn}{k}k}{k - 1} F(k - 1) + \left(T - \frac{n}{k}\right) F(0) + K \quad (18)$$

$$= \frac{cn}{k} F(k) + \frac{n(1 - c)}{k - 1} F(k - 1) + \left(T - \frac{n}{k}\right) F(0) + K \quad (19)$$

$$= \frac{dn}{k} F(k) + \left(T - \frac{n}{k}\right) F(0) + K, \quad (20)$$

where

$$d := c + (1 - c)\beta, \quad \text{for } \beta := \frac{k}{F(k)} \frac{F(k - 1)}{k - 1}. \quad (21)$$

Note that  $F(k)/k \rightarrow \infty$  as  $k \rightarrow \infty$ . Along with the convexity of  $F$ , it follows that there exists a  $k_0$  such that  $\beta < 1$  for all  $k > k_0$ . Note that  $k > (3 + \alpha)^2$  suffices. This implies that  $d < 1$ , which shows that there is a non-zero gap between the possible values of  $\Phi$ .  $\square$

Note that the maximum concentration objective,  $F(n) = n \log n$ , satisfies the conditions on  $F$  and, in particular, we have  $\beta < 1$  for  $k = 3$ .

## B Derivation of MAP $\theta$ objective

$$\Pr(\theta | \mathbf{w}) \propto \sum_{z_1, \dots, z_N} \Pr(\theta) \Pr(w_1, \dots, w_N, z_1, \dots, z_N | \theta) \quad (22)$$

$$= \Pr(\theta) \sum_{z_1, \dots, z_N} \prod_{i=1}^N \Pr(z_i | \theta) \Pr(w_i | z_i) \quad (23)$$

$$= \Pr(\theta) \prod_{i=1}^N \sum_{z_i} \Pr(z_i | \theta) \Pr(w_i | z_i) \quad (24)$$

$$= \Pr(\theta) \prod_{i=1}^N \sum_{t=1}^T \theta_t \Pr(w_i | z_i = t) \quad (25)$$

$$\propto \prod_{t=1}^T \theta_t^{\alpha_t - 1} \prod_{i=1}^N \sum_{t=1}^T \theta_t \Pr(w_i | z_i = t). \quad (26)$$

### C Proof of Lemma 3

If  $\epsilon \geq K(\alpha, T, N)$  the claim trivially holds. Assume for the purpose of contradiction that there exists a word  $\hat{i}$  such that  $\theta_{\hat{i}}^* < K(\alpha, T, N)$ , where  $\hat{i} = \arg \max_t \psi_{\hat{i}t} \theta_t^*$ .

Let  $Y$  denote the set of topics  $t \neq \hat{i}$  such that  $\theta_t^* \geq 2\epsilon$ . Let  $\beta_1 = \sum_{t \in Y} \theta_t^*$  and  $\beta_2 = \sum_{t \notin Y, t \neq \hat{i}} \theta_t^*$ . Note that  $\beta_2 < 2T\epsilon$ . Consider  $\hat{\theta}$  defined as follows:

$$\hat{\theta}_{\hat{i}} = \frac{1}{N} \quad (27)$$

$$\hat{\theta}_t = \left( \frac{1 - \beta_2 - \frac{1}{N}}{\beta_1} \right) \theta_t^* \text{ for } t \in Y \quad (28)$$

$$\hat{\theta}_t = \theta_t^* \text{ for } t \notin Y, t \neq \hat{i}. \quad (29)$$

Note that this construction implies the bound  $\hat{\theta}_t \geq (1 - 2T\epsilon - \frac{1}{N}) \theta_t^*$  for  $t \in Y$ . Assuming  $n \geq 4$  and  $\epsilon \leq \frac{1}{2TN}$ , we have that  $\hat{\theta}_t \geq \frac{1}{2} \theta_t^* \geq \epsilon$  for  $t \in Y$ , so  $\hat{\theta}$  is feasible.

We will show that  $\Phi(\hat{\theta}) > \Phi(\theta^*)$ , contradicting the optimality of  $\theta^*$ . First we need the following upper bound, which uses the fact that  $\theta_{\hat{i}}^* < \frac{1}{N}$ :

$$\frac{1 - \beta_2 - \frac{1}{N}}{\beta_1} = \frac{1 - \beta_2 - \frac{1}{N}}{1 - \beta_2 - \theta_{\hat{i}}^*} \quad (30)$$

$$< 1. \quad (31)$$

Then, we have:

$$\frac{P(\hat{\theta})}{\alpha - 1} = \sum_{t=1}^T \log(\hat{\theta}_t) \quad (32)$$

$$= \log \frac{1}{N} + \sum_{t \in Y} \log \left( \frac{1 - \beta_2 - \frac{1}{N}}{\beta_1} \right) + \sum_{t \neq \hat{i}} \log \theta_t^* \quad (33)$$

$$\leq \log \frac{1}{N} + \sum_{t \neq \hat{i}} \log \theta_t^*. \quad (34)$$

Thus,

$$\frac{P(\hat{\theta}) - P(\theta^*)}{\alpha - 1} \leq \log \frac{1}{N} - \log \theta_{\hat{i}}^* \quad (35)$$

which when  $\alpha < 1$  gives the inequality:

$$P(\hat{\theta}) - P(\theta^*) \geq (\alpha - 1) \left( \log \frac{1}{N} - \log \theta_{\hat{i}}^* \right) \quad (36)$$

$$= (1 - \alpha) \left( \log N + \log \theta_{\hat{i}}^* \right). \quad (37)$$

Moving on to the second term, we have:

$$L(\hat{\theta}) = \sum_{j \in [N]: j \neq \hat{i}} \log \left( \sum_t \psi_{jt} \hat{\theta}_t \right) + \log \left( \sum_t \psi_{\hat{i}t} \hat{\theta}_t \right) \quad (38)$$

$$\geq (N - 1) \log \left( \frac{1 - \beta_2 - \frac{1}{N}}{\beta_1} \right) + \sum_{j \in [N]: j \neq \hat{i}} \log \left( \sum_t \psi_{jt} \theta_t^* \right) + \log \left( \frac{\psi_{\hat{i}\hat{i}}}{N} \right) \quad (39)$$

$$\geq (N - 1) \log \left( 1 - \frac{2}{N} \right) + \sum_{j \in [N]: j \neq \hat{i}} \log \left( \sum_t \psi_{jt} \theta_t^* \right) + \log \left( \frac{\psi_{\hat{i}\hat{i}}}{N} \right). \quad (40)$$

$$L(\hat{\theta}) - L(\theta^*) \geq (N-1) \log \left(1 - \frac{2}{N}\right) + \log \left(\frac{\psi_{i\hat{i}}}{N}\right) - \log \left(\sum_t \psi_{it} \theta_t^*\right) \quad (41)$$

$$\geq (N-1) \log \left(1 - \frac{2}{N}\right) + \log \left(\frac{\psi_{i\hat{i}}}{N}\right) - \log (T \psi_{i\hat{i}} \theta_{\hat{i}}^*) \quad (42)$$

$$= (N-1) (\log(N-2) - \log N) + \log \left(\frac{1}{N}\right) - \log (T \theta_{\hat{i}}^*) \quad (43)$$

$$\geq (N-1) \left(-\frac{2}{N-2}\right) + \log \left(\frac{1}{N}\right) - \log (T \theta_{\hat{i}}^*) \quad (44)$$

$$\geq -3 + \log \left(\frac{1}{N}\right) - \log (T \theta_{\hat{i}}^*). \quad (45)$$

where we used the lower bound  $\log(N-2) \geq \log N - \frac{2}{N-2}$  that arises from the convexity of the  $\log(x)$  function, and again assumed  $N \geq 4$ .

Finally, putting these two together, we have:

$$\Phi(\hat{\theta}) - \Phi(\theta^*) \geq -3 - \alpha \log N - \log T - \alpha \log \theta_{\hat{i}}^* \quad (46)$$

Plugging in  $\theta_{\hat{i}}^* < K(\alpha, T, N)$  results in  $\Phi(\hat{\theta}) - \Phi(\theta^*) > 0$ , giving the contradiction.

## D Proof of Theorem 8

Here we reduce from the *unique* set cover problem, where we are guaranteed that there is only one minimal size set that covers all elements. It can be shown that Unique Set Cover is NP-hard (under randomized reductions) by using standard reductions from Unique SAT to Vertex Cover, and then from Vertex Cover to Set Cover.

Consider a Unique Set Cover instance and our standard reduction to an LDA instance as described in earlier sections. In particular, let  $\mathbf{w} = (w_1, \dots, w_N)$  denote a Unique Set Cover reduction instance and let  $C \subseteq [T]$  denote the unique minimum cover. Let  $S_i$  be those sets (topics) that cover element  $w_i$ . We will show that for sufficiently small hyperparameters  $\alpha_t$ , we can determine whether a set (topic)  $t \in C$  by testing the value of  $\mathbb{E}[\theta_t | X]$ , thus proving that the computation of the latter is NP-hard.

We have

$$p(\theta | X) \propto \prod_t \theta_t^{\alpha_t - 1} \prod_i \sum_{t' \in S_i} \theta_{t'} \quad (47)$$

$$= \sum_r \prod_t \theta_t^{\alpha_t - 1 + \eta_t(r)}, \quad (48)$$

where the final summation is over elements  $r \in \mathcal{R} := S_1 \times \dots \times S_N$  and  $\eta_t(r) := |\{i : r_i = t\}|$ . For  $r \in \mathcal{R}$ , we write  $|r|$  to denote the number of topics  $t$  such that  $\eta_t(r) \neq 0$ .

Let  $\mathcal{N}$  denote the set of sequences  $n = (n_1, \dots, n_T)$  such that  $n_t \geq 0$  and  $\sum_t n_t = N$ . For  $n \in \mathcal{N}$ , define

$$Z(n) := \int \dots \int \prod_t \theta_t^{\alpha_t - 1 + n_t} d\theta_1 \dots d\theta_T = \frac{\prod_t \Gamma(\alpha_t + n_t)}{\Gamma(\bar{\alpha} + N)}, \quad (49)$$

where  $\bar{\alpha} = \sum_t \alpha_t$ . It follows that

$$\mathbb{E}[\theta_t | X] = \int \dots \int \theta_t \cdot p(\theta | X) d\theta_1 \dots d\theta_T \quad (50)$$

$$= \frac{1}{\sum_r Z(\eta(r))} \int \dots \int \theta_t \sum_r \prod_{\tau} \theta_{\tau}^{\alpha_{\tau} - 1 + n_{\tau}(r)} d\theta_1 \dots d\theta_T \quad (51)$$

$$= \frac{1}{\sum_r Z(\eta(r))} \sum_r Z(\eta(r, t)), \quad (52)$$

where  $\eta(r, t) \in \mathcal{N}$  is given by  $\eta(r, t)_\tau = \eta_r(r) + 1(\tau = t)$ . By the identity  $\Gamma(z + 1) = z\Gamma(z)$ , we have  $Z(\eta(r, t)) = Z(\eta(r)) \frac{\alpha_t + n_t(r)}{\bar{\alpha} + N}$ , and so it follows that

$$\mathbb{E}[\theta_t | X] = \sum_{r \in \mathcal{R}} \bar{Z}(\eta(r)) \frac{\alpha_t + n_t(r)}{\bar{\alpha} + N}, \quad (53)$$

where

$$\bar{Z}(n) := \frac{Z(n)}{\sum_{r' \in \mathcal{R}} Z(\eta(r'))}. \quad (54)$$

For  $c \in [T]$ , let  $\mathcal{R}_c$  denote the set of those  $r \in \mathcal{R}$  such that  $|r| = c$ . Recall that  $C \subseteq [T]$  is the unique minimum set cover associated with the reduction  $X$ . Thus, for  $t \in C$ ,

$$\mathbb{E}[\theta_t | X] = \sum_{r \in \mathcal{R}_{|C|}} \bar{Z}(\eta(r)) \frac{\alpha_t + n_t(r)}{\bar{\alpha} + N} + \sum_{r \in \mathcal{R} \setminus \mathcal{R}_{|C|}} \bar{Z}(\eta(r)) \frac{\alpha_t + n_t(r)}{\bar{\alpha} + N} \quad (55)$$

$$\geq \frac{\alpha_t + 1}{\bar{\alpha} + N} \sum_{r \in \mathcal{R}_{|C|}} \bar{Z}(\eta(r)), \quad (56)$$

where we have used the observation that  $\eta_t(r) \geq 1$  for  $r \in \mathcal{R}_{|C|}$ . Let  $\beta = \sum_{r \in \mathcal{R}_{|C|}} \bar{Z}(\eta(r))$ . If  $t \notin C$ , then  $n_t(r) = 0$  for  $r \in \mathcal{R}_{|C|}$  and so we have  $\mathbb{E}[\theta_t | X] \leq \beta \frac{\alpha_t}{\bar{\alpha} + N} + (1 - \beta)$ . It follows that if

$$\beta > \frac{1}{2} \left( 1 + \frac{\bar{\alpha} + N}{\bar{\alpha} + N + 1} \right) \quad (57)$$

then the minimum cover  $C$  contains a topic  $t$  if and only if  $\mathbb{E}[\theta | X] \geq \frac{1}{4} \left( 1 + 3 \frac{\bar{\alpha} + N}{\bar{\alpha} + N + 1} \right) \frac{\alpha_t + 1}{\bar{\alpha} + N}$ , and, moreover, we can determine the minimum cover from a bound on  $\beta$  and polynomial approximations to the marginal distributions of the components of  $\theta$ . We will take  $\alpha_t = \alpha$  henceforth, and show that for  $\alpha$  small enough, the bound (57) indeed holds.

Let  $n, n' \in \mathcal{N}$  be topic counts associated with the minimal cover and some non-minimal cover, respectively. That is, let  $n = \eta(r)$  for some  $r \in \mathcal{R}_{|C|}$  and let  $n' = \eta(r')$  for some  $r' \in \mathcal{R}_k$  and  $k > |C|$ . We will bound  $Z(n')/Z(n)$  in order to bound  $\beta$ . We have

$$\prod_t \Gamma(\alpha + n_t) \geq \Gamma(\alpha)^{T-|C|} \Gamma(\alpha + 1)^{|C|}, \quad (58)$$

whereas

$$\prod_t \Gamma(\alpha + n'_t) \leq \Gamma(\alpha)^{T-|C|-1} \Gamma(\alpha + 1)^{|C|} \Gamma(\alpha + N - |C|). \quad (59)$$

Therefore,

$$\frac{Z(n')}{Z(n)} \leq \frac{\Gamma(\alpha)^{T-|C|-1} \Gamma(\alpha + 1)^{|C|} \Gamma(\alpha + N - |C|)}{\Gamma(\alpha)^{T-|C|} \Gamma(\alpha + 1)^{|C|}} \quad (60)$$

$$= \frac{\Gamma(\alpha + N - |C|)}{\Gamma(\alpha)}. \quad (61)$$

By the convexity of  $\Gamma(1/c)$  in  $c$ , we have  $\Gamma(\alpha) \geq \alpha^{-1} - \gamma$ , where  $\gamma \approx .577$  is the Euler constant. Therefore,

$$\frac{Z(n')}{Z(n)} \leq \frac{\Gamma(\alpha + N - 1)}{\alpha^{-1} - \gamma} =: \kappa(\alpha). \quad (62)$$

Then by conservatively assuming that there is only one responsibility corresponding to the minimum cover, we have that

$$\beta \geq \frac{Z(n)}{Z(n) + T^N \kappa(\alpha) Z(n)}. \quad (63)$$

Therefore, the bound (57) is achieved when

$$\kappa(\alpha) \leq \frac{1}{T^N(2\bar{\alpha} + 2N + 1)}. \quad (64)$$

In particular, when  $N, T \geq 2$  and

$$\alpha^{-1} > 2T^N \Gamma(N)(2N + 2) \quad (65)$$

the marginal expectations can be used to read off the unique minimal set cover.