

439 Appendices

440 The appendices that follow provide the proofs of the results in the body of the paper. Throughout
 441 the proofs in the appendix we use the following notation to denote the hitting and movement costs
 442 of the online learner: $H_t := f_t(x_t)$ and $M_t := c(x_t, x_{t-1})$, where x_t is the point chosen by the
 443 online algorithm at time t . Similarly, we denote the hitting and movement costs of the offline optimal
 444 (adversary) as $H_t^* := f_t(x_t^*)$ and $M_t^* := c(x_t^*, x_{t-1}^*)$, where x_t^* is the point chosen by the offline
 445 optimal at time t .

446 Before moving to the proofs, we summarize a few standard definitions that are used throughout the
 447 paper.

448 **Definition 1.** A function $f : \mathcal{X} \rightarrow \mathbb{R}$ is α -strongly convex with respect to a norm $\|\cdot\|$ if for all x, y in
 449 the relative interior of the domain of f and $\lambda \in (0, 1)$, we have

$$f(\lambda x + (1 - \lambda)y) \leq \lambda f(x) + (1 - \lambda)f(y) - \frac{\alpha}{2}\lambda(1 - \lambda)\|x - y\|^2.$$

450 **Definition 2.** A function $f : \mathcal{X} \rightarrow \mathbb{R}$ is β -strongly smooth with respect to a norm $\|\cdot\|$ if f is
 451 everywhere differentiable and if for all x, y we have

$$f(y) \leq f(x) + \langle \nabla f(x), y - x \rangle + \frac{\beta}{2}\|y - x\|^2.$$

452 **Definition 3.** A function $f : \mathbb{R}^d \rightarrow \mathbb{R}$ is quasiconvex if its domain \mathcal{X} and all its sublevel sets

$$S_\alpha = \{x \in \mathcal{X} \mid f(x) \leq \alpha\},$$

453 for $\alpha \in \mathbb{R}$, is convex.

454 **Definition 4.** For a norm $\|\cdot\|$ in \mathcal{X} , its dual norm (on \mathcal{X}) $\|\cdot\|_*$ is defined to be

$$\|y\|_* = \sup\{\langle x, y \rangle \mid \|x\| \leq 1\}.$$

455 **Definition 5.** For a convex function $f : \mathcal{X} \rightarrow \mathbb{R}$, its Fenchel Conjugate f^* is defined to be

$$f^*(y) = \sup\{\langle x, y \rangle - f(x) \mid x \in \mathcal{X}\}.$$

456 Next, we introduce a few technical lemmas that are important throughout our analysis.

457 The first technical lemma is a characterization of strongly convex functions.

458 **Lemma 1.** Suppose f is α -strongly convex for some $\alpha > 0$ with respect to some norm $\|\cdot\|$ and
 459 both f and f^* are differentiable, then the first condition implies the second condition and the third
 460 condition:

- 461 1. $\forall x, y, f(y) \leq f(x) + \langle \nabla f(x), y - x \rangle + \frac{\beta}{2}\|x - y\|^2$;
- 462 2. $\forall x, y, f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle + \frac{1}{2\beta}\|\nabla f(x) - \nabla f(y)\|_*^2$;
- 463 3. $\forall x, y, \|\nabla f(x) - \nabla f(y)\|_* \leq \beta\|x - y\|$.

464 To prove Lemma 1, we use Lemma 2, Lemma 3, and Lemma 4 below.

465 The following lemma is Theorem 6 in [25].

466 **Lemma 2.** If f is convex and closed, the following two conditions are equivalent:

- 467 1. $\forall x, y, f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle + \frac{\beta}{2}\|x - y\|^2$;
- 468 2. $\forall x, y, f^*(y) \leq f^*(x) + \langle \nabla f^*(x), y - x \rangle + \frac{1}{2\beta}\|x - y\|_*^2$

469 i.e. f is β -strongly convex w.r.t some norm $\|\cdot\|$ if and only if f^* is $\frac{1}{\beta}$ -strongly smooth w.r.t the dual
 470 norm $\|\cdot\|_*$.

471 The next lemma is a special case of Lemma 17 in [34].

472 **Lemma 3.** Let f be a closed, convex, and differentiable function. Then we have

$$f^*(\nabla f(x)) + f(x) = \langle \nabla f(x), x \rangle.$$

473 Now we prove a technical result that describes a property of the gradient of the Fenchel Conjugate.

474 **Lemma 4.** Suppose f is α -strongly convex for some $\alpha > 0$ with respect to some norm $\|\cdot\|$ and both
475 f and f^* are differentiable. Then we have

$$x = \nabla f^*(\nabla f(x)), \forall x.$$

476 *Proof.* For convenience, we define $y = \nabla f(x)$ and $x' = \nabla f^*(y)$. It suffices to prove that $x' = x$.

477 By Lemma 3, we obtain

$$f^*(y) + f(x) = \langle y, x \rangle = \langle x, y \rangle. \quad (1)$$

478 Again by Lemma 3, we have

$$f(x') + f^*(y) = f^{**}(x') + f^*(y) = \langle x', y \rangle, \quad (2)$$

479 where we use the fact that $f^{**} = f$.

480 Combining inequalities (1) and (2), we obtain

$$0 = f(x) - f(x') - \langle x - x', y \rangle = f(x) + \langle x' - x, \nabla f(x) \rangle - f(x') \leq -\frac{\alpha}{2} \|x - x'\|^2,$$

481 where in the last inequality we use the definition of α -strongly convex. Therefore we have proved
482 that $x = x'$. \square

483 Using the three lemmas above, we now prove Lemma 1.

484 *Proof of Lemma 1.* By the first condition and Lemma 2, we know f^* is $\frac{1}{\beta}$ -strongly convex with
485 respect to $\|\cdot\|_*$. Therefore we see

$$f^*(\nabla f(y)) \geq f^*(\nabla f(x)) + \langle \nabla f^*(\nabla f(x)), \nabla f(y) - \nabla f(x) \rangle + \frac{1}{2\beta} \|\nabla f(x) - \nabla f(y)\|_*^2.$$

486 Using Lemma 3 and Lemma 4, we obtain

$$\langle y, \nabla f(y) \rangle - f(y) \geq (\langle x, \nabla f(x) \rangle - f(x)) + \langle x, \nabla f(y) - \nabla f(x) \rangle + \frac{1}{2\beta} \|\nabla f(x) - \nabla f(y)\|_*^2.$$

487 Rearranging the terms, we get

$$f(x) \geq f(y) + \langle x - y, \nabla f(y) \rangle + \frac{1}{2\beta} \|\nabla f(x) - \nabla f(y)\|_*^2,$$

488 which is the second condition.

489 The third condition follows from subtracting the second condition from the first condition. \square

490 Finally, before moving the the proofs of our main results, we prove two properties of the Bregman
491 Divergence that play an important role in the analysis.

492 **Lemma 5.** $\forall a, b, c \in \mathbb{R}^d$ and potential h , we have

$$\langle \nabla h(b) - \nabla h(c), c - a \rangle = D_h(a||b) - D_h(a||c) - D_h(c||b).$$

493 *Proof.* By the definition of Bregman Divergence, we obtain

$$\begin{aligned} & D_h(a||b) - D_h(a||c) - D_h(c||b) \\ &= (h(a) - h(b) - \langle \nabla h(b), a - b \rangle) - (h(a) - h(c) - \langle \nabla h(c), a - c \rangle) \\ &\quad - (h(c) - h(b) - \langle \nabla h(b), c - b \rangle) \\ &= -\langle \nabla h(b), a - b \rangle + \langle \nabla h(c), a - c \rangle + \langle \nabla h(b), c - b \rangle \\ &= (-\langle \nabla h(b), a - b \rangle + \langle \nabla h(b), c - b \rangle) + \langle \nabla h(c), a - c \rangle \\ &= \langle \nabla h(b), c - a \rangle + \langle \nabla h(c), a - c \rangle \\ &= \langle \nabla h(b) - \nabla h(c), c - a \rangle. \end{aligned}$$

494 \square

Lemma 6. For all $a, b, c \in \mathbb{R}^d$, we have

$$D_h(c||a) - D_h(c||b) = D_h(0||a) - D_h(0||b) + \langle \nabla h(b) - \nabla h(a), c \rangle.$$

495 *Proof.* Using the definition of Bregman divergence, we obtain

$$\begin{aligned} D_h(c||a) - D_h(c||b) &= h(c) - h(a) - \langle \nabla h(a), c - a \rangle - h(c) + h(b) + \langle \nabla h(b), c - b \rangle \\ &= (h(b) - \langle \nabla h(b), b \rangle) - (h(a) - \langle \nabla h(a), a \rangle) + \langle \nabla h(b) - \nabla h(a), c \rangle \\ &= D_h(0||a) - D_h(0||b) + \langle \nabla h(b) - \nabla h(a), c \rangle. \end{aligned}$$

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□

497 A Proof of Theorem 1

498 We consider a sequence of hitting cost functions on the real line such that the algorithm stays at the
499 starting point through time steps $t = 1, 2, \dots, n$ and is forced to incur a huge movement cost at time
500 step $t = n + 1$, whereas the offline adversary can pay relatively little cost by dividing the long trek
501 between x_0 and v_{n+1} into multiple small steps through time steps $t = 1, 2, \dots, n + 1$.

502 Specifically, suppose the starting point of the algorithm and the offline adversary is $x_0 = x_0^* = 0$,
503 and the hitting cost functions are

$$f_t(x) = \begin{cases} \frac{m}{2}x^2 & t \in \{1, 2, \dots, n\} \\ \frac{m'}{2}(x-1)^2 & t = n+1 \end{cases}$$

504 for some large parameter m' that we choose later.

505 Suppose the algorithm first moves at time step t_0 . If $t_0 < n + 1$, we stop the game at time step t_0 and
506 compare the algorithm with an offline adversary which always stays at $x = 0$. The total cost of offline
507 adversary is 0, but the total cost of the algorithm is non-zero. So, the competitive ratio is unbounded.

508 Next we consider the case where $t_0 \geq n + 1$. This implies that $x_1, \dots, x_n = 0$ and x_{n+1} is some
509 non-zero point, say x . We see that the cost incurred by the online algorithm is

$$\text{cost}(ALG) \geq \min_{x_{n+1}} (M_{n+1} + H_{n+1}) = \min_x \left(\frac{1}{2}x^2 + \frac{m'}{2}(x-1)^2 \right).$$

510 Notice that the right hand side tends to $\frac{1}{2}$ as m' tends to infinity; specifically, we have

$$\text{cost}(ALG) \geq \min_x \left(\frac{1}{2}x^2 + \frac{m'}{2}(x-1)^2 \right) = \frac{1}{2(1 + \frac{1}{m'})}. \quad (3)$$

511 Now let us consider the offline optimal. Notice that, in the limit as m' tends to infinity, the offline
512 optimal must satisfy $x_0^* = 0$ and $x_{n+1}^* = 1$; otherwise it would incur unbounded cost. Our lower
513 bound is derived by considering the case when $m' \rightarrow \infty$ and so we constrain the adversary to satisfy
514 the above, knowing that the adversary is not optimal for finite m' , i.e., $\text{cost}(ADV) \geq \text{cost}(OPT)$
515 with $\text{cost}(ADV) \rightarrow \text{cost}(OPT)$ as $m' \rightarrow \infty$.

516 Let the sequence of points the adversary chooses as $x^* = (x_0^*, x_1^*, \dots, x_{n+1}^*) \in \mathbb{R}^{n+2}$. We compute
517 the cost incurred by the adversary as follows where, to simplify presentation, we define $\mathcal{K}(n, y)$ to be
518 the set $\{x \in \mathbb{R}^{n+2} \mid x_i \leq x_{i+1}, x_0 = 0, x_{n+1} = y\}$.

$$\begin{aligned} a_n &= 2 \min_{x^* \in \mathcal{K}(n, 1)} \sum_{i=1}^{n+1} (H_i^* + M_i^*) \\ &= 2 \min_{x^* \in \mathcal{K}(n, 1)} \left(\sum_{i=1}^n \frac{m}{2} (x_i^*)^2 + \sum_{i=1}^{n+1} \frac{1}{2} (x_i^* - x_{i-1}^*)^2 \right). \end{aligned}$$

519 In words, a_n is twice the minimal offline cost subject to the constraints $x_0^* = 0, x_{n+1}^* = 1$. We derive
520 the limiting behavior of the offline costs as $n \rightarrow \infty$ in the following lemma.

521 **Lemma 7.** For $m > 0$, define

$$a_n = 2 \min_{x^* \in \mathcal{K}(n,1)} \left(\sum_{i=1}^n \frac{m}{2} (x_i^*)^2 + \sum_{i=1}^{n+1} \frac{1}{2} (x_i^* - x_{i-1}^*)^2 \right).$$

522 Then we have $\lim_{n \rightarrow \infty} a_n = \frac{-m + \sqrt{m^2 + 4m}}{2}$.

523 Given the lemma, the total cost of the offline adversary will be $\frac{a_n}{2}$. Finally, applying (3), we know
524 $\forall n$ and $\forall m' > 0$,

$$\frac{\text{cost}(\text{ALG})}{\text{cost}(\text{ADV})} \geq \frac{\frac{1}{2(1+\frac{1}{m'})}}{\frac{a_n}{2}} = \frac{1}{(1 + \frac{1}{m'})a_n}.$$

525 By taking the limit $n \rightarrow \infty$ and $m' \rightarrow \infty$ and using Lemma 7, we obtain

$$\frac{\text{cost}(\text{ALG})}{\text{cost}(\text{OPT})} = \lim_{n, m' \rightarrow \infty} \frac{\text{cost}(\text{ALG})}{\text{cost}(\text{ADV})} \geq \left(\frac{-m + \sqrt{m^2 + 4m}}{2} \right)^{-1} = \frac{1 + \sqrt{1 + \frac{4}{m}}}{2}.$$

526 All that remains is to prove Lemma 7, which describes the cost of the offline adversary in the limit as
527 n tends to infinity.

528 *Proof of Lemma 7.* Using the fact that the costs are all homogeneous of degree 2, we see that for all
529 $y \in [0, 1]$, we have

$$\begin{aligned} & \min_{x^* \in \mathcal{K}(n,y)} \left(\sum_{i=1}^n \frac{m}{2} (x_i^*)^2 + \sum_{i=1}^{n+1} \frac{1}{2} (x_i^* - x_{i-1}^*)^2 \right) \\ &= y^2 \min_{x^* \in \mathcal{K}(n,1)} \left(\sum_{i=1}^n \frac{m}{2} (x_i^*)^2 + \sum_{i=1}^{n+1} \frac{1}{2} (x_i^* - x_{i-1}^*)^2 \right). \end{aligned} \quad (4)$$

530 The sequence $\{a_n\}$, $n \geq 0$ has a recursive relationship as follows:

$$\begin{aligned} a_{n+1} &= 2 \min_{x^* \in \mathcal{K}(n+1,1)} \left(\sum_{i=1}^{n+1} \frac{m}{2} (x_i^*)^2 + \sum_{i=1}^{n+2} \frac{1}{2} (x_i^* - x_{i-1}^*)^2 \right) \\ &= 2 \min_{0 \leq x \leq 1} \left(\min_{x^* \in \mathcal{K}(n,x)} \left(\sum_{i=1}^n \frac{m}{2} (x_i^*)^2 + \sum_{i=1}^{n+1} \frac{1}{2} (x_i^* - x_{i-1}^*)^2 \right) \right. \\ &\quad \left. + \frac{m}{2} x^2 + \frac{1}{2} (1-x)^2 \right) \\ &= 2 \min_{0 \leq x \leq 1} \left(x^2 \min_{x^* \in \mathcal{K}(n,1)} \left(\sum_{i=1}^n \frac{m}{2} (x_i^*)^2 + \sum_{i=1}^{n+1} \frac{1}{2} (x_i^* - x_{i-1}^*)^2 \right) \right. \\ &\quad \left. + \frac{m}{2} x^2 + \frac{1}{2} (1-x)^2 \right) \\ &= 2 \min_{0 \leq x \leq 1} \left(\frac{a_n}{2} x^2 + \frac{m}{2} x^2 + \frac{1}{2} (1-x)^2 \right) \\ &= \frac{a_n + m}{a_n + m + 1}. \end{aligned} \quad (5)$$

531 Solving the equation $x = \frac{x+m}{x+m+1}$, we find the two fixed points of the recursive relationship $a_{n+1} =$
532 $\frac{a_n + m}{a_n + m + 1}$ are

$$x_1 = \frac{-m + \sqrt{m^2 + 4m}}{2},$$

533 and

$$x_2 = \frac{-m - \sqrt{m^2 + 4m}}{2}.$$

534 Notice that for $i = 1, 2$, we have

$$m - (m + 1)x_i = -(1 - x_i)x_i.$$

535 Using this property, we obtain

$$a_{n+1} - x_1 = \frac{a_n + m}{a_n + m + 1} - x_1 = \frac{(1 - x_1)a_n + m - (m + 1)x_1}{a_n + m + 1} = \frac{(1 - x_1)(a_n - x_1)}{a_n + m + 1}, \quad (6)$$

536 and

$$a_{n+1} - x_2 = \frac{a_n + m}{a_n + m + 1} - x_2 = \frac{(1 - x_2)a_n + m - (m + 1)x_2}{a_n + m + 1} = \frac{(1 - x_2)(a_n - x_2)}{a_n + m + 1}. \quad (7)$$

Notice that $a_{n+1} - x_2 > 0$. By dividing equations (6) and (7), we obtain

$$\left(\frac{a_{n+1} - x_1}{a_{n+1} - x_2} \right) = \frac{1 - x_1}{1 - x_2} \cdot \left(\frac{a_n - x_1}{a_n - x_2} \right), \forall n \geq 0.$$

Remember that $a_0 = 1$. Therefore we have

$$\left(\frac{a_n - x_1}{a_n - x_2} \right) = \left(\frac{1 - x_1}{1 - x_2} \right)^n \left(\frac{a_0 - x_1}{a_0 - x_2} \right) = \left(\frac{1 - x_1}{1 - x_2} \right)^{n+1}.$$

Rearranging this equation, we get

$$a_n = \left(1 - \left(\frac{1 - x_1}{1 - x_2} \right)^{n+1} \right)^{-1} \left(x_1 - x_2 \cdot \left(\frac{1 - x_1}{1 - x_2} \right)^{n+1} \right).$$

537 Since $0 < \left(\frac{1 - x_1}{1 - x_2} \right) < 1$, we have

$$\lim_{n \rightarrow \infty} a_n = x_1 = \frac{-m + \sqrt{m^2 + 4m}}{2}. \quad (8)$$

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□

539 B Proof of Theorem 2

540 Our proof of Theorem 2 relies on a set of technical lemmas, which follow. Lemma 8 and Lemma
541 10 work together to establish a lower bound on the competitive ratio as m tends to zero when the
542 balance parameter γ is set to be $o(1/m)$, while Lemma 11 lower bound on the competitive ratio as
543 m tends to zero when the balance parameter γ is set to be $\Omega(1/m)$.

544 **Lemma 8.** *If $\gamma = o(1/m)$, the competitive ratio of OBD is $\Omega(1/(\gamma m))$ when $m \rightarrow 0^+$.*

545 *Proof.* Our approach is to construct a scenario where OBD is forced to move along the circumference
546 of a large circle, but the offline adversary moves along the circumference of a much smaller circle (see
547 Figure 1). The adversary is hence able to pay much smaller movements costs, forcing the competitive
548 ratio to be large.

549 We propose a series of costs which force OBD to move in a circle. The idea is to construct a cost
550 function so that, at the end of every round, the relative positions of the OBD algorithm, the offline
551 adversary, and the minimizer are fixed. Since OBD is memoryless, we can simply input this function
552 arbitrarily many times and the positions of OBD and the offline adversary will trace out a pair of
553 concentric circles (see Figure 1).

554 Suppose that, at the start of a round, OBD is at the point A . Let ℓ be the distance between OBD and
555 the adversary. Consider a right triangle ABC such that $|AB| = h = \sqrt{\gamma m} \ell$, the offline adversary is
556 at some point D on the hypotenuse AC and $|AD| = |BC| = \ell$ (see Figure 2). Let us introduce a
557 coordinate system such that the origin lies at C , the x -axis contains BC and the y -axis is parallel

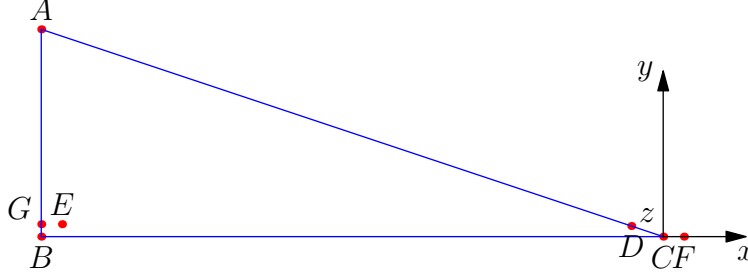


Figure 2: In the right triangle $\triangle ABC$, $\angle ABC = 90^\circ$, $|BC| = \ell$, $|AB| = h = \sqrt{\gamma m} \ell$. Point D is on the line segment AC such that $|AD| = \ell$. OBD starts at point A and selects point E . The offline adversary starts at point D and selects point F . G is the projection point of E on line segment AB .

to AB , such that the positive part of the axis lies on the same side of BC as the segment AC . Our goal is to construct a cost function which forces OBD towards B . This will preserve the relative positions of OBD and the adversary, since we assumed that they were a distance ℓ away at the start of the round. Consider the costs $g(u) = \frac{m}{2} \|u - C\|^2$, $h(u) = \alpha \cdot d(u, BC)$ where $d(u, BC)$ is the distance from the point u to the line passing through B and C and $\alpha > 0$ is a parameter we will pick later. Define $f_t(u) = h(u) + g(u)$. Notice that f_t is m -strongly convex because it is the sum of an m -strongly convex function and a convex function. Intuitively, when α is large, the function f_t is infinity outside of the line BC but is equal to $g(u) = \frac{m}{2} \|u - C\|^2$ when restricted to points u on the line. After observing the cost f_t , OBD will pick some new point E .

The following lemma highlights that E can be driven arbitrarily close to B by taking α to be sufficiently large.

Lemma 9. Let $\varepsilon > 0$, and suppose α is picked to that $\alpha > \frac{hm\ell^2}{\varepsilon^2}$. Then the point E picked by OBD satisfies $|EB| < \varepsilon$.

We instruct the adversary to pick the point F on the line BC (the x -axis) such that $EF = \ell$ (see Figure 2). Notice that $|CF| = |BF| - |BC| \leq |BE| + |EF| - |BC| = |EB| + \ell - \ell < \varepsilon$, where we used the triangle inequality. Let $z = |DC|$. We see that the total cost incurred by the offline adversary is

$$M_t^* + H_t^* = \frac{1}{2} |DF|^2 + \frac{m}{2} |CF|^2 \leq \frac{1}{2} (|DC| + |CF|)^2 + \frac{m}{2} |CF|^2 \leq \frac{1}{2} (z + \varepsilon)^2 + \frac{m\varepsilon^2}{2},$$

where we applied the triangle inequality.

Notice that $h = |AB| = \sqrt{|AC|^2 - |BC|^2}$ by the Pythagorean theorem (recall that ABC is a right triangle). Since $|AC| = \ell + z$ and $|BC| = \ell$, we see that $h = \sqrt{2z\ell + z^2}$. Hence the movement cost incurred by the OBD is

$$M_t \geq \frac{1}{2} (h - \varepsilon)^2 = \frac{1}{2} (\sqrt{2z\ell + z^2} - \varepsilon)^2.$$

Hence the ratio of the costs is

$$\frac{M_t + H_t}{M_t^* + H_t^*} \geq \frac{M_t}{M_t^* + H_t^*} \geq \frac{\frac{1}{2} (\sqrt{2z\ell + z^2} - \varepsilon)^2}{\frac{1}{2} (z + \varepsilon)^2 + \frac{m\varepsilon^2}{2}}.$$

Since the limit of this expression as $\varepsilon \rightarrow 0$ is $\frac{2z\ell + z^2}{z^2}$, for sufficiently small ε this will be at least $\frac{1}{2} \frac{2z\ell + z^2}{z^2} \geq \frac{\ell}{z}$. Since $z = \sqrt{h^2 + \ell^2} - \ell$ and $h = \sqrt{\gamma m} \ell$, the ratio of costs is at least

$$\frac{\ell}{\sqrt{\gamma m \ell^2 + \ell^2} - \ell} = \frac{1}{\sqrt{\gamma m + 1} - 1} = \frac{\sqrt{\gamma m + 1} + 1}{\gamma m} \geq \frac{2}{\gamma m}.$$

Now, we describe the whole process. When $t = 1$, the hitting cost function is $f_1(x) = \frac{m}{2} \|x\|_2^2$. While OBD stays at $x = 0$, the adversary moves to the point $(\ell, 0)$; it incurs a one-time cost of



Figure 3: *Balance condition at time step t in Lemma 10. Starting from x_{t-1} , OBD picks x_t after observing the hitting cost function $f_t(x) = \frac{m}{2}(x - t)^2$.*

584 $M_1^* + H_1^* = \frac{1}{2}\ell^2 + \frac{m}{2}\ell^2$. On all subsequent steps $t = 2 \dots T$, we repeatedly apply the construction,
 585 which forces OBD to move in a circle. The one-time cost incurred by the adversary to setup the game
 586 is negligible in the limit as T is large, and the per-round ratio of costs is $\Omega(\frac{1}{\gamma m})$, so the competitive
 587 ratio is also $\Omega(\frac{1}{\gamma m})$ as claimed. \square

588 The key technical lemma used in the proof is Lemma 9, and we now provide a proof of that result.

589 *Proof of Lemma 9.* Suppose $\alpha > \frac{hm\ell^2}{\varepsilon^2}$. We first show that OBD selects the point E strictly contained
 590 by the $\frac{m}{2}\ell^2$ -level set, which is the one B lies on. First observe that the point B satisfies the balance
 591 condition: $\frac{1}{2}|AB|^2 = \gamma \frac{m}{2}|BC|^2$, because we constructed ABC so that $|AB| = h = \sqrt{\gamma m}\ell$ and
 592 $|BC| = \ell$. However, the point B is not necessarily a projection of A onto any level set of f_t . If OBD
 593 projected onto the level set which B lies on, it would incur less cost than if it moved to B ; however
 594 then the balance condition would be violated. To restore the balance condition, we must increase the
 595 movement cost while decreasing the hitting cost – which means we must move to a strictly smaller
 596 level set, say the $\frac{m}{2}l_1^2$ -level set, where $l_1 < \ell$.

597 Let E_y denote the y -coordinate of E , using the coordinate system we define in the proof of Lemma
 598 8. Notice that $E_y = \frac{g(E)}{\alpha}$, since $g(E)$ was defined to be the vertical distance to the x -axis times α .
 599 Since $g(E) \leq f_t(E)$, we see that $E_y \leq \frac{f_t(E)}{\alpha} = \frac{ml_1^2}{2\alpha} \leq \frac{ml^2}{2\alpha}$, where we used the fact that E lies on
 600 the $\frac{m}{2}l_1^2$ level set and $l_1 \leq \ell$. By the balance condition, $\frac{1}{2}|AE|^2 = \gamma \frac{m}{2}l_1^2 \leq \gamma \frac{m}{2}l^2 = \frac{1}{2}h^2$. Let G be
 601 the point with coordinates (B_x, E_y) . Applying the Pythagorean theorem successively to the right
 602 triangle BEG and the right triangle AEG , we see that

$$\begin{aligned} |EB|^2 &= |E_x - B_x|^2 + E_y^2 \leq (|AE|^2 - (|AB| - E_y)^2) + E_y^2 \\ &\leq (|AB|^2 - (|AB| - E_y)^2) + E_y^2 \leq 2h \cdot E_y \leq h \frac{ml^2}{\alpha}, \end{aligned} \quad (9)$$

603 where we used the fact that $|AB| \geq |AE|$ and $|AB| = h$. Since we picked $\alpha > \frac{hm\ell^2}{\varepsilon^2}$, we see that
 604 $|EB| < \varepsilon$. \square

606 Now we move on to the next technical lemma in the proof of Theorem 2.

607 **Lemma 10.** *When $\gamma = o(\frac{1}{m})$, the competitive ratio of OBD is $\Omega(\sqrt{\frac{\gamma}{m}})$.*

608 *Proof.* We consider a sequence of cost functions on the real line such that the OBD algorithm moves
 609 far away from the starting point, incurring significant movement costs, whereas the offline adversary
 610 could pay relatively little cost by staying at the starting point. More specifically, we consider the
 611 sequence of hitting cost functions $f_t(x) = \frac{m}{2}(x - t)^2, t = 1, 2, \dots, n$. The value of n will be picked
 612 later. We assume the starting point is at zero.

613 Notice that by the balance condition we always have $M_t = \gamma H_t$, so $\frac{1}{2}\|x_t - x_{t-1}\|^2 = \gamma \frac{m}{2}\|x_t - t\|^2$.

614 We can rearrange this expression to obtain $\frac{x_t - x_{t-1}}{t - x_{t-1}} = \sqrt{\gamma m}$. Define

$$\lambda = \frac{x_t - x_{t-1}}{t - x_{t-1}} = \frac{\sqrt{\gamma m}}{1 + \sqrt{\gamma m}}.$$

615 We obtain the recursive equation $x_t = x_{t-1} + (t - x_{t-1})\lambda$ with initial condition $x_0 = 0$. Solving
 616 this equation, we obtain $x_t = t - \frac{1-\lambda}{\lambda}(1 - (1 - \lambda)^t)$.

617 Suppose we picked n to be $\lceil \frac{1}{\lambda} \rceil$. By assumption, $\gamma = o(\frac{1}{m})$; hence in the limit as m tends to zero, λ
 618 also tends to zero. Notice that $x_n = n - \frac{1-\lambda}{\lambda}(1 - (1-\lambda)^n) \geq \frac{1}{\lambda} \frac{1}{2e} - (1 - \frac{1}{e}) \geq \frac{1}{6\lambda}$ for sufficiently
 619 small λ . Here we used the fact that $(1-\lambda)^{\frac{1}{\lambda}} \rightarrow e^{-1}$.

620 Suppose the next cost function is $f_{n+1}(x) = m'x^2$. Notice that if the offline adversary simply stays
 621 at zero throughout the game, the total cost it incurs would be

$$\text{cost}(ADV) = \frac{m}{2}(1^2 + 2^2 + \dots + n^2) \leq \frac{mn^3}{2} = \Theta\left(\frac{m}{\lambda^3}\right) = \Theta\left(\frac{1}{\sqrt{\gamma^3 m}}\right).$$

622 In the last step, we used the fact that λ tends to $\sqrt{\gamma m}$ when $\gamma = o(\frac{1}{m})$ and m tends to zero.

623 If we pick m' large enough that OBD is forced to incur movement cost at least $\frac{1}{2}(\frac{x_n}{2})^2$, the total cost
 624 incurred by OBD is

$$\text{cost}(OBD) \geq \frac{1}{2}\left(\frac{x_n}{2}\right)^2 = \Theta\left(\frac{1}{\lambda^2}\right) = \Theta\left(\frac{1}{\gamma m}\right).$$

625 Putting these facts together, we see that the competitive ratio is at least $\Theta(\sqrt{\frac{\gamma}{m}})$. \square

626 The last technical lemma used to proof Theorem 2 is the following.

627 **Lemma 11.** *When $\gamma = \Omega(\frac{1}{m})$, the competitive ratio of OBD is $\Omega(\frac{1}{m})$.*

628 *Proof.* Since $\gamma = \Omega(\frac{1}{m})$, we can assume there exists $C > 0$ such that $\gamma \geq C/m$. We again
 629 consider a situation such that the OBD algorithm moves far away from the starting point, incurring
 630 significant movement cost, whereas the offline adversary could pay relatively little cost by staying at
 631 the starting point. More specifically, suppose the starting point is zero and the first cost function is
 632 $f_1(x) = \frac{m}{2}(1-x)^2$. Suppose the adversary stays at zero. The cost incurred by the adversary will be

$$\text{cost}(ADV) = \frac{m}{2}.$$

633 Notice that by the balance condition ($M_t = \gamma H_t$), the point x_1 picked by OBD satisfies $\frac{x_1^2}{2} =$
 634 $\gamma \frac{m}{2}(1-x_1)^2$. So the cost incurred by OBD is lower bounded by

$$\text{cost}(OBD) \geq M_1 = \frac{1}{2}\left(\frac{\sqrt{\gamma m}}{1 + \sqrt{\gamma m}}\right)^2 \geq \frac{1}{2}\left(\frac{\sqrt{C}}{1 + \sqrt{C}}\right)^2.$$

635 Since C is a positive constant, the competitive ratio of OBD is lower bounded by $\frac{OBD}{ADV} = \Theta(\frac{1}{m})$. \square

636 Now we return to the proof of Theorem 2. This proof is a straightforward combination of the above
 637 lemmas. When $\gamma = o(\frac{1}{m})$, by combining Lemma 8 and Lemma 10, we know the competitive
 638 ratio is at least $\max\left(\frac{C_1}{\gamma m}, C_2 \sqrt{\frac{\gamma}{m}}\right)$ for some positive constants C_1, C_2 . Notice that function $\frac{C_1}{\gamma m}$ is
 639 monotonically decreasing in γ and $C_2 \sqrt{\frac{\gamma}{m}}$ is monotonically increasing in γ . Solving the equation
 640 $\frac{C_1}{\gamma m} = C_2 \sqrt{\frac{\gamma}{m}}$, we get $\gamma = \left(\frac{C_1}{C_2}\right)^{\frac{2}{3}} m^{-\frac{1}{3}}$. Therefore we see that

$$\max\left\{\frac{C_1}{\gamma m}, C_2 \sqrt{\frac{\gamma}{m}}\right\} \geq C_1^{\frac{1}{3}} C_2^{\frac{2}{3}} m^{-\frac{2}{3}} = \Theta(m^{-\frac{2}{3}}).$$

641 On the other hand, when $\gamma = \Omega(\frac{1}{m})$, by Lemma 11, we know the competitive ratio of OBD is lower
 642 bounded by $\Theta(\frac{1}{m})$.

643 Together, the above implies that the competitive ratio of OBD is at least $\Theta(m^{-\frac{2}{3}})$ when $m \rightarrow 0^+$.

644 C Proof of Theorem 3

645 To begin, note that it is sufficient to prove result for all positive $m \leq \frac{9}{64}$. Similarly, it also suffices
 646 to show Theorem 3 when the minimum of every hitting cost function is zero, since otherwise the
 647 competitive ratio can only improve if this is not the case.

648 Our argument makes use of the following potential function: $\phi(x_t, x_t^*) = \eta \|x_t - x_t^*\|^2$. We define
 649 $\Delta\phi = \phi(x_t, x_t^*) - \phi(x_{t-1}, x_{t-1}^*)$ and $\Delta\phi' = \phi(x'_t, x_t^*) - \phi(x_{t-1}, x_{t-1}^*)$. It suffices to show that
 650 $H_t + M_t + \Delta\phi \leq C(H_t^* + M_t^*)$, for some positive constant C . From this inequality, we can sum
 651 over all timesteps t to yield that the competitive ratio is upper bounded by C :

$$\sum_{t=0}^T H_t + M_t \leq \sum_{t=0}^T H_t + M_t + \Delta\phi \leq C \sum_{t=0}^T (H_t^* + M_t^*).$$

652 Throughout the proof, we fix $\eta = 4$ and use $\|\cdot\|$ to denote ℓ_2 norm. When we refer to generalized
 653 mean inequality, we mean

$$(a + b)^2 \leq 2a^2 + 2b^2, \forall a, b \in \mathbb{R}.$$

654 We define $H'_t := f_t(x'_t)$ and $M'_t := c(x'_t, x_{t-1}) = \frac{1}{2} \|x'_t - x_{t-1}\|_2^2$, where x'_t is the point chosen by
 655 the first OBD phase (line 3) of Algorithm 2.

656 Before we move to the main casework in the proof, we begin with a technical lemma that we use to
 657 bound the change in the potential function.

658 **Lemma 12.** Suppose the potential function $\phi : \mathbb{R}^d \times \mathbb{R}^d \rightarrow \mathbb{R}_{\geq 0}$ is defined as $\phi(a, b) = \eta \|a - b\|^2$,
 659 where $\eta > 0$. Then $\forall \lambda > 0$, the change in potential satisfies

$$\phi(a, c) - \phi(a, b) \leq (1 + \lambda^2)\phi(b, c) + \frac{1}{\lambda^2}\phi(a, b),$$

660 for all $a, b, c \in \mathbb{R}^d$.

661 *Proof.* Using the triangle inequality, we obtain

$$\|a - c\|^2 \leq (\|a - b\| + \|b - c\|)^2 = \|a - b\|^2 + \|b - c\|^2 + 2\|a - b\|\|b - c\|.$$

662 Rearranging the terms, we obtain

$$\begin{aligned} \|a - c\|^2 - \|a - b\|^2 &\leq \|b - c\|^2 + 2\|a - b\|\|b - c\| \\ &= \|b - c\|^2 + 2\left(\frac{1}{\lambda}\|a - b\|\right)(\lambda\|b - c\|) \\ &\leq (1 + \lambda^2)\|b - c\|^2 + \frac{1}{\lambda^2}\|a - b\|^2, \end{aligned}$$

663 where in the last line we use the AM-GM inequality. □

664 We are now ready to precede with the proof, which is divided up into two cases based on the
 665 relationship between the hitting cost of the algorithm and that of the adversary.

666 **Case 1:** $H'_t \leq H_t^*$

667 Since the hitting cost function satisfies $f_t(x) \geq \frac{m}{2} \|x - v_t\|^2$, by the triangle inequality, we have

$$\|x'_t - x_t^*\| \leq \|x'_t - v_t\| + \|x_t^* - v_t\| \leq \left(\sqrt{\frac{2H'_t}{m}} + \sqrt{\frac{2H_t^*}{m}} \right). \quad (10)$$

668 Thus the change in potential satisfies

$$\begin{aligned}
\frac{1}{\eta} \Delta \phi' &= \|x'_t - x_t^*\|^2 - \|x_{t-1} - x_{t-1}^*\|^2 \\
&= (\|x'_t - x_t^*\| - \|x_{t-1} - x_{t-1}^*\|)(\|x'_t - x_t^*\| + \|x_{t-1} - x_{t-1}^*\|) \\
&\leq (\|x'_t - x_{t-1}\| + \|x_t^* - x_{t-1}^*\|)(\|x'_t - x_{t-1}\| + \|x_t^* - x_{t-1}^*\| + 2\|x'_t - x_t^*\|) \quad (11a) \\
&= (\|x'_t - x_{t-1}\| + \|x_t^* - x_{t-1}^*\|)^2 + 2(\|x'_t - x_{t-1}\| + \|x_t^* - x_{t-1}^*\|)\|x'_t - x_t^*\| \\
&\leq 2\|x'_t - x_{t-1}\|^2 + 2\|x_t^* - x_{t-1}^*\|^2 + 2(\|x'_t - x_{t-1}\| + \|x_t^* - x_{t-1}^*\|)\|x'_t - x_t^*\| \quad (11b) \\
&\leq 4M'_t + 4M_t^* + 2(\sqrt{2M'_t} + \sqrt{2M_t^*}) \left(\sqrt{\frac{2H'_t}{m}} + \sqrt{\frac{2H_t^*}{m}} \right) \quad (11c) \\
&\leq 4M'_t + 4M_t^* + \sqrt{\frac{1}{m}}((\sqrt{2M'_t} + \sqrt{2M_t^*})^2 + (\sqrt{2H'_t} + \sqrt{2H_t^*})^2) \quad (11d) \\
&\leq 4M'_t + 4M_t^* + \sqrt{\frac{1}{m}}((4M'_t + 4M_t^*) + (4H'_t + 4H_t^*)) \quad (11e) \\
&= \left(4 + 4\sqrt{\frac{1}{m}}\right) M'_t + \left(4 + 4\sqrt{\frac{1}{m}}\right) M_t^* + 4\sqrt{\frac{1}{m}} H'_t + 4\sqrt{\frac{1}{m}} H_t^*,
\end{aligned}$$

669 where we use the triangle inequality in line (11a); the generalized mean inequality in lines (11b),
670 (11d) and (11e) and inequality (10) in line (11c).

671 Using the OBD's balance condition $M'_t = \gamma H'_t$ and the assumption $H'_t \leq H_t^*$ based on inequality
672 (11), we have

$$\begin{aligned}
\frac{1}{\eta} \Delta \phi' &\leq \left(4 + 4\sqrt{\frac{1}{m}}\right) \gamma H'_t + \left(4 + 4\sqrt{\frac{1}{m}}\right) M_t^* + 4\sqrt{\frac{1}{m}} H'_t + 4\sqrt{\frac{1}{m}} H_t^* \\
&\leq \left(4 + 4\sqrt{\frac{1}{m}}\right) \gamma H_t^* + \left(4 + 4\sqrt{\frac{1}{m}}\right) M_t^* + 8\sqrt{\frac{1}{m}} H_t^*.
\end{aligned}$$

673 Notice that by the triangle inequality and the generalized mean inequality, we have that

$$M_t = \frac{1}{2} \|x_t - x_{t-1}\|^2 \leq \frac{1}{2} (\|x'_t - x_{t-1}\| + \|x_t - x'_t\|)^2 \leq \frac{1}{2} (2\|x'_t - x_{t-1}\|^2 + 2\|x_t - x'_t\|^2).$$

Remember that since $\mu = 1$, we have $\|x_t - x'_t\|^2 = m \|x'_t - v_t\|^2$. Using this fact, we derive the following bound on $H_t + M_t + \Delta\phi$:

$$\begin{aligned} H_t + M_t + \Delta\phi &\leq H'_t + \frac{1}{2} \left(2 \|x'_t - x_{t-1}\|^2 + 2 \|x_t - x'_t\|^2 \right) \\ &\quad + \eta (\|x_t - x_t^*\|^2 - \|x'_t - x_t^*\|^2) + \Delta\phi' \\ &\leq H'_t + (2M'_t + m \|x'_t - v_t\|^2) \\ &\quad + \left(\eta \left(1 + \frac{1}{\sqrt{m}} \right) \|x_t - x'_t\|^2 + \eta\sqrt{m} \|x'_t - x_t^*\|^2 \right) + \Delta\phi' \end{aligned} \quad (12a)$$

$$\begin{aligned} &\leq H'_t + (2M'_t + m \|x'_t - v_t\|^2) \\ &\quad + \left(\eta \left(1 + \frac{1}{\sqrt{m}} \right) m \|x'_t - v_t\|^2 + \eta\sqrt{m} \left(2 \|x'_t - v_t\|^2 + 2 \|x_t^* - v_t\|^2 \right) \right) \\ &\quad + \Delta\phi' \end{aligned} \quad (12b)$$

$$\begin{aligned} &\leq H'_t + (2M'_t + 2H'_t) + \left(\eta \left(1 + \frac{1}{\sqrt{m}} \right) 2H'_t + \eta\sqrt{m} \left(\frac{4H'_t}{m} + \frac{4H_t^*}{m} \right) \right) + \Delta\phi' \end{aligned} \quad (12c)$$

$$\begin{aligned} &= (3 + 2\eta + \frac{6\eta}{\sqrt{m}})H'_t + 2M'_t + 4\eta \frac{H_t^*}{\sqrt{m}} + \Delta\phi' \\ &= \left(3 + 2\eta + \frac{6\eta}{\sqrt{m}} + 2\gamma \right) H'_t + 4\eta \frac{H_t^*}{\sqrt{m}} + \Delta\phi' \\ &\leq \left(3 + 2\eta + \frac{6\eta}{\sqrt{m}} + 2\gamma \right) H_t^* + 4\eta \frac{H_t^*}{\sqrt{m}} + \Delta\phi' \\ &= \left(3 + 2\eta + \frac{10\eta}{\sqrt{m}} + 2\gamma \right) H_t^* + \Delta\phi', \end{aligned} \quad (12d)$$

where we use Lemma 12 in line (12a); the triangle inequality in line (12b); m -strongly convexity of f_t in line (12c); and the assumption $H'_t \leq H_t^*$ in line (12d).

Combining inequalities (11) and (12), we obtain

$$H_t + M_t + \Delta\phi \leq (3 + 2\eta + 2\gamma + 4\eta\gamma + \frac{\eta}{\sqrt{m}}(18 + 4\gamma))H_t^* + \eta(4 + 4\sqrt{\frac{1}{m}})M_t^*. \quad (13)$$

Case 2: $H'_t \geq H_t^*$

In this case, we prove that for any $x_t^*, x_{t-1}^* \in \mathbb{R}^d$, we have

$$H_t + M_t + \Delta\phi \leq \frac{C}{\sqrt{m}}(H_t^* + M_t^*), \quad (14)$$

for some positive constant C .

In the proof, we use D_1, D_2, \dots, D_d to represent the d axes in the coordinate system.

As shown in Figure 4, without loss of generality, let $v_t = (0, 0, \dots, 0)$, $x'_t = (h_1, h_2, 0, \dots, 0)$ and $D_2 = h_2$ be the projection hyper plane, where $h_1 \geq 0, h_2 \geq 0$. And let $l = \|x_{t-1} - x'_t\| > 0$. Note that our analysis still holds in one-dimension because we can restrict ourselves to the D_2 axis.

Then we know $x_{t-1} = (h_1, h_2 + l, 0, \dots, 0)$, $x_t = (h_1(1 - \sqrt{m}), h_2(1 - \sqrt{m}), 0, \dots, 0)$. Since we know x_t^* must lie below the projection hyper plane, we can let $x_t^* = (x, h_2 - y, a_3, a_4, \dots, a_d)$, where $y > 0$.

Now we show that it suffices to prove the statement when x_{t-1}^* is on the line segment $x_t^*x_{t-1}$. Suppose x_{t-1}^* is not on the line segment $x_t^*x_{t-1}$. If $\|x_{t-1}^* - x_{t-1}\| > \|x_t^* - x_{t-1}\|$, by moving x_{t-1}^* to x_t^* , $\Delta\phi$ increases and M_t^* decreases. Otherwise, we can choose a point K on line segment $x_t^*x_{t-1}$ such that $\|K - x_{t-1}\| = \|x_{t-1}^* - x_{t-1}\|$. By moving x_{t-1}^* to K , $\Delta\phi$ remains unchanged and M_t^* decreases. Therefore if inequality (14) holds for x_{t-1}^* on the segment $x_t^*x_{t-1}$, then it must also hold for any other $x_{t-1}^* \in \mathbb{R}^d$.

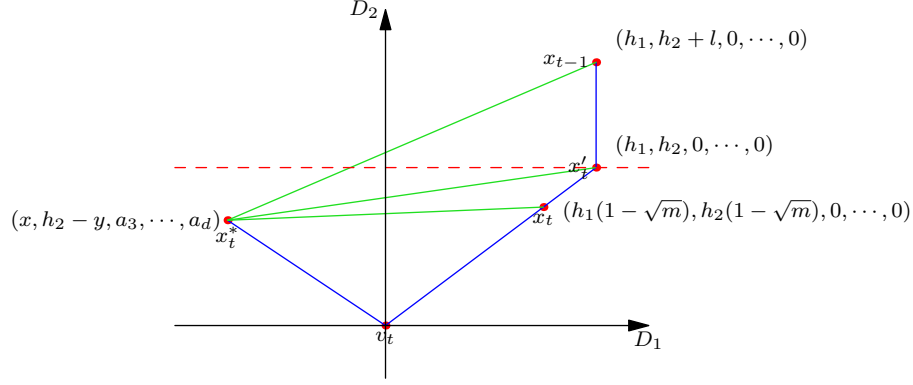


Figure 4: Starting at x_{t-1} , G -OBD first does projection on to the H'_t level set (red dashed line) in the first phase. The projection point is x'_t . Then G -OBD moves toward the minimizer to obtain point x_t in the second phase. Let the minimizer v_t be the origin. Notice that the three points x_{t-1} , x'_t , v_t defines a plane S . Without loss of generality, we can let axis D_2 be parallel to line $x'_t x_{t-1}$; and let axis D_1 be parallel to the projection hyperplane.

Now we suppose x_{t-1}^* is on the line segment $x_t^* x_{t-1}$, and $\|x_t^* - x_{t-1}^*\| = \lambda \|x_t^* - x_{t-1}\|$.

Recall that we set $\gamma = 1$, so $M'_t = \gamma H'_t = H'_t$. It follows that

$$M_t \leq l^2 + \|x_t - x'_t\|^2 = l^2 + m(h_1^2 + h_2^2) \leq l^2 + 2H'_t = l^2 + 2M'_t \leq 2l^2,$$

and

$$H_t \leq H'_t = M'_t = \frac{l^2}{2}.$$

We can separate $\Delta\phi$ into two parts:

$$\frac{\Delta\phi}{\eta} = \left(\|x_t^* - x_t\|^2 - \|x_t^* - x_{t-1}\|^2 \right) + \left(\|x_t^* - x_{t-1}\|^2 - \|x_{t-1}^* - x_{t-1}\|^2 \right).$$

For convenience, we define

$$\Delta\phi_1 := \left(\|x_t^* - x_t\|^2 - \|x_t^* - x_{t-1}\|^2 \right),$$

and

$$\Delta\phi_2 := \left(\|x_t^* - x_{t-1}\|^2 - \|x_{t-1}^* - x_{t-1}\|^2 \right).$$

We further notice that from the triangle inequality,

$$\Delta\phi_2 \leq (1 - (1 - \lambda)^2) \|x_t^* - x_{t-1}\|^2 = \lambda(2 - \lambda) \left((x - h_1)^2 + (y + l)^2 + \sum_{i=3}^d a_i^2 \right). \quad (15)$$

Now we express M_t^* and H_t^* in terms of the variables we define, which are

$$M_t^* = \frac{1}{2} (\lambda \|x_t^* - x_{t-1}\|)^2 = \frac{\lambda^2}{2} \left((x - h_1)^2 + (y + l)^2 + \sum_{i=3}^d a_i^2 \right), \quad (16)$$

and

$$H_t^* \geq \frac{m}{2} \|x_t^* - v_t\|^2 = \frac{m}{2} \left(x^2 + (h_2 - y)^2 + \sum_{i=3}^d a_i^2 \right). \quad (17)$$

We also expand $\Delta\phi_1$:

$$\begin{aligned} \Delta\phi_1 &= \|x_t^* - x_t\|^2 - \|x_t^* - x_{t-1}\|^2 \\ &= (x - h_1 + h_1\sqrt{m})^2 + (y - h_2\sqrt{m})^2 + \sum_{i=3}^d a_i^2 - (x - h_1)^2 - (y + l)^2 - \sum_{i=3}^d a_i^2 \\ &= ((x - h_1 + h_1\sqrt{m})^2 - (x - h_1)^2) + ((y - h_2\sqrt{m})^2 - (y + l)^2) \\ &= h_1\sqrt{m}(2x - 2h_1 + h_1\sqrt{m}) - (h_2\sqrt{m} + l)(2y + l - h_2\sqrt{m}) \\ &= 2xh_1\sqrt{m} - 2h_1^2\sqrt{m} + h_1^2m - 2y(h_2\sqrt{m} + l) - l^2 + h_2^2m. \end{aligned} \quad (18)$$

705 Using the condition that $m \leq \frac{9}{64} < 1$, we derive the following bound:

$$\begin{aligned}\Delta\phi_1 &\leq 2xh_1\sqrt{m} - 2h_1^2\sqrt{m} + h_1^2m - 2y(h_2\sqrt{m} + l) - l^2 + h_2^2\sqrt{m} \\ &= 2xh_1\sqrt{m} - 2h_1^2\sqrt{m} + h_1^2m + \sqrt{m}(h_2 - y)^2 - \sqrt{m}y^2 - 2yl - l^2.\end{aligned}\quad (19)$$

706 Substituting equations (16) and (17) into inequality (14), we know that it suffices to show that for
707 some constant C ,

$$M_t + H_t + \eta\Delta\phi_1 + \eta\Delta\phi_2 \leq \frac{C}{\sqrt{m}} \left(\frac{m}{2} \left(x^2 + (h_2 - y)^2 + \sum_{i=3}^d a_i^2 \right) + \frac{\lambda^2}{2} \left((x - h_1)^2 + (y + l)^2 + \sum_{i=3}^d a_i^2 \right) \right). \quad (20)$$

708 **Subcase 2.1:** $\lambda \leq \frac{\sqrt{m}}{2}$

709 We can bound equation (15) as follows:

$$\begin{aligned}\Delta\phi_2 &= \lambda(2 - \lambda) \left((x - h_1)^2 + (y + l)^2 + \sum_{i=3}^d a_i^2 \right) \\ &\leq \sqrt{m}(x - h_1)^2 + \sqrt{m}(y + l)^2 + \sqrt{m} \sum_{i=3}^d a_i^2 \\ &= \sqrt{m}x^2 - 2\sqrt{m}xh_1 + \sqrt{m}h_1^2 + \sqrt{m}y^2 + 2\sqrt{m}yl + \sqrt{m}l^2 + \sqrt{m} \sum_{i=3}^d a_i^2.\end{aligned}\quad (21)$$

710 Summing inequalities (19) and (21), we get

$$\begin{aligned}\Delta\phi_1 + \Delta\phi_2 &\leq \sqrt{m}x^2 + (-h_1^2\sqrt{m} + h_1^2m) + \sqrt{m}(h_2 - y)^2 \\ &\quad + (2\sqrt{m}yl - 2yl) + (\sqrt{m}l^2 - l^2) + \sqrt{m} \sum_{i=3}^d a_i^2 \\ &\leq \sqrt{m}x^2 + 0 + \sqrt{m}(h_2 - y)^2 + 0 - \frac{5}{8}l^2 + \sqrt{m} \sum_{i=3}^d a_i^2 \\ &\leq \sqrt{m}x^2 + \sqrt{m}(h_2 - y)^2 - \frac{5}{8}l^2 + \sqrt{m} \sum_{i=3}^d a_i^2,\end{aligned}\quad (22a)$$

711 where we use the condition that $m \leq \frac{9}{64}$ in line (22a). We further obtain

$$\begin{aligned}M_t + H_t + \eta(\Delta\phi_1 + \Delta\phi_2) &\leq 2l^2 + \frac{l^2}{2} + \eta \left(\sqrt{m}x^2 + \sqrt{m}(h_2 - y)^2 - \frac{5}{8}l^2 + \sqrt{m} \sum_{i=3}^d a_i^2 \right) \\ &= \frac{5l^2}{2} + 4 \left(\sqrt{m}x^2 + \sqrt{m}(h_2 - y)^2 - \frac{5}{8}l^2 + \sqrt{m} \sum_{i=3}^d a_i^2 \right) \\ &= 4 \left(\sqrt{m}x^2 + \sqrt{m}(h_2 - y)^2 + \sqrt{m} \sum_{i=3}^d a_i^2 \right).\end{aligned}$$

712 Therefore, for $C \geq 8$, we have

$$M_t + H_t + \eta\Delta\phi_1 + \eta\Delta\phi_2 \leq \frac{C}{\sqrt{m}} \left(\frac{m}{2} \left(x^2 + (h_2 - y)^2 + \sum_{i=3}^d a_i^2 \right) + \frac{\lambda^2}{2} \left((x - h_1)^2 + (y + l)^2 + \sum_{i=3}^d a_i^2 \right) \right),$$

713 which establishes inequality (20).

714 **Subcase 2.2:** $\lambda \geq \frac{\sqrt{m}}{2}$

715 Notice that when $C \geq 32$, we have

$$\frac{C}{2\sqrt{m}}\lambda^2 \geq \frac{16}{\sqrt{m}}\lambda^2 \geq \frac{16}{\sqrt{m}} \cdot \frac{\sqrt{m}}{2}\lambda = 8\lambda \geq 4\lambda(2-\lambda) = \eta\lambda(2-\lambda).$$

716 Substituting this inequality into equation (15), we know that for $C \geq 32$,

$$\eta\Delta\phi_2 \leq \frac{C}{\sqrt{m}} \cdot \frac{\lambda^2}{2} \left((x-h_1)^2 + (y+l)^2 + \sum_{i=3}^d a_i^2 \right). \quad (23)$$

717 We can further bound inequality (19):

$$\begin{aligned} \Delta\phi_1 &\leq 2xh_1\sqrt{m} - 2h_1^2\sqrt{m} + h_1^2m + \sqrt{m}(h_2-y)^2 - \sqrt{m}y^2 - 2yl - l^2 \\ &\leq \sqrt{m}x^2 + \sqrt{m}h_1^2 - 2h_1^2\sqrt{m} + h_1^2m + \sqrt{m}(h_2-y)^2 - l^2 \\ &\leq \sqrt{m}x^2 + \sqrt{m}(h_2-y)^2 - l^2, \end{aligned}$$

718 where we apply the AM-GM inequality in step 2 and use the condition $m < 1$ in step 3.

719 Therefore we have

$$\begin{aligned} H_t + M_t + \eta\Delta\phi_1 &\leq \frac{5l^2}{2} + 4(\sqrt{m}x^2 + \sqrt{m}(h_2-y)^2 - l^2) \\ &\leq 4(\sqrt{m}x^2 + \sqrt{m}(h_2-y)^2). \end{aligned} \quad (24)$$

720 Summing inequalities (24) and (23), we yield that for $C \geq 32$,

$$M_t + H_t + \eta\Delta\phi_1 + \eta\Delta\phi_2 \leq \frac{C}{\sqrt{m}} \left(\frac{m}{2} (x^2 + (h_2-y)^2 + \sum_{i=3}^d a_i^2) + \frac{\lambda^2}{2} ((x-h_1)^2 + (y+l)^2 + \sum_{i=3}^d a_i^2) \right),$$

721 which establishes inequality (20).

722 Combining all cases above, we conclude that G-OBd is an $O(\frac{1}{\sqrt{m}})$ -competitive algorithm.

723 **D Proof of Theorem 4**

724 To prove Theorem 4 we make use of Lemma 1 and 5.

725 Our approach is to make use of strong convexity and properties of Bregman Divergences to derive an
726 inequality in the form of $H_t + M_t + \Delta\phi \leq C(H_t^* + M_t^*)$ for some positive constant C , where $\Delta\phi$
727 is the change in potential, which we will define later. The constant C is then an upper bound for the
728 competitive ratio.

729 To begin, recall that h is assumed to be α -strongly convex and β -strongly smooth with respect to
730 norm $\|\cdot\|$. Thus we can give a trivial bound on Bregman Divergence, namely

$$\forall x, y, \frac{\alpha}{2} \|x - y\|^2 \leq D_h(x||y) \leq \frac{\beta}{2} \|x - y\|^2. \quad (25)$$

731 Recall that the update rule in Algorithm 3 can be stated as:

$$x_t = \arg \min_x f_t(x) + \lambda_1 D_h(x||x_{t-1}) + \lambda_2 D_h(x||v_t).$$

732 Since the function $f_t(x) + \lambda_1 D_h(x||x_{t-1}) + \lambda_2 D_h(x||v_t)$ is strongly convex, the minimizer x_t exists
733 and is unique. Furthermore, it must satisfy the first-order condition

$$\nabla f_t(x_t) + \lambda_1(\nabla h(x_t) - \nabla h(x_{t-1})) + \lambda_2(\nabla h(x_t) - \nabla h(v_t)) = 0.$$

734 Further, since $f_t(x)$ is m -strongly convex, we have

$$\begin{aligned} f_t(x_t^*) &\geq f_t(x_t) + \langle \nabla f_t(x_t), x_t^* - x_t \rangle + \frac{m}{2} \|x_t^* - x_t\|^2 \\ &= f_t(x_t) - \lambda_1 \langle \nabla h(x_{t-1}) - \nabla h(x_t), x_t - x_t^* \rangle \\ &\quad - \lambda_2 \langle \nabla h(v_t) - \nabla h(x_t), x_t - x_t^* \rangle + \frac{m}{2} \|x_t^* - x_t\|^2. \end{aligned} \quad (26)$$

Using Lemma 5, we obtain

$$\langle \nabla h(x_{t-1}) - \nabla h(x_t), x_t - x_t^* \rangle = D_h(x_t^* || x_{t-1}) - D_h(x_t^* || x_t) - D_h(x_t || x_{t-1}),$$

and

$$\langle \nabla h(v_t) - \nabla h(x_t), x_t - x_t^* \rangle = D_h(x_t^* || v_t) - D_h(x_t^* || x_t) - D_h(x_t || v_t).$$

Substituting the two above identities into inequality (26), we get

$$\begin{aligned} & f_t(x_t) + \lambda_1 D_h(x_t || x_{t-1}) + \lambda_2 D_h(x_t || v_t) + (\lambda_1 + \lambda_2) D_h(x_t^* || x_t) + \frac{m}{2} \|x_t^* - x_t\|^2 \\ & \leq f_t(x_t^*) + \lambda_1 D_h(x_t^* || x_{t-1}) + \lambda_2 D_h(x_t^* || v_t). \end{aligned}$$

It follows that

$$\begin{aligned} & f_t(x_t) + \lambda_1 D_h(x_t || x_{t-1}) + (\lambda_1 + \lambda_2) D_h(x_t^* || x_t) + \frac{m}{2} \|x_t^* - x_t\|^2 \\ & \leq f_t(x_t^*) + \lambda_1 D_h(x_t^* || x_{t-1}) + \lambda_2 D_h(x_t^* || v_t). \end{aligned} \quad (27)$$

We define the potential function as $\phi(x_t, x_t^*) = (\lambda_1 + \lambda_2) D_h(x_t^* || x_t) + \frac{m}{2} \|x_t^* - x_t\|^2$, and let $\Delta\phi = \phi(x_t, x_t^*) - \phi(x_{t-1}, x_{t-1}^*)$. Applying this notation to inequality (27) and rearranging terms, we obtain

$$\begin{aligned} & H_t + \lambda_1 M_t + \Delta\phi \\ & \leq (H_t^* + \lambda_2 D_h(x_t^* || v_t)) + \lambda_1 D_h(x_t^* || x_{t-1}) - (\lambda_1 + \lambda_2) D_h(x_{t-1}^* || x_{t-1}) - \frac{m}{2} \|x_{t-1}^* - x_{t-1}\|^2. \end{aligned} \quad (28)$$

Using Lemma 1, we get

$$\frac{1}{2\beta} \|\nabla h(x_{t-1}) - \nabla h(x_{t-1}^*)\|_*^2 \leq D_h(x_{t-1}^* || x_{t-1}), \quad (29)$$

and

$$\|\nabla h(x_{t-1}) - \nabla h(x_{t-1}^*)\|_* \leq \beta \|x_{t-1} - x_{t-1}^*\|. \quad (30)$$

Using Lemma 5 and the two above inequalities, we get

$$\begin{aligned} & \lambda_1 D_h(x_t^* || x_{t-1}) - (\lambda_1 + \lambda_2) D_h(x_{t-1}^* || x_{t-1}) - \frac{m}{2} \|x_{t-1}^* - x_{t-1}\|^2 \\ & = \lambda_1 (D_h(x_t^* || x_{t-1}) - D_h(x_{t-1}^* || x_{t-1})) - \lambda_2 D_h(x_{t-1}^* || x_{t-1}) - \frac{m}{2} \|x_{t-1}^* - x_{t-1}\|^2 \end{aligned} \quad (31a)$$

$$\begin{aligned} & = \lambda_1 D_h(x_t^* || x_{t-1}^*) + \lambda_1 \langle \nabla h(x_{t-1}) - \nabla h(x_{t-1}^*), x_{t-1}^* - x_t^* \rangle \\ & \quad - \lambda_2 D_h(x_{t-1}^* || x_{t-1}) - \frac{m}{2} \|x_{t-1}^* - x_{t-1}\|^2 \end{aligned} \quad (31b)$$

$$\begin{aligned} & \leq \lambda_1 D_h(x_t^* || x_{t-1}^*) + \lambda_1 \|\nabla h(x_{t-1}) - \nabla h(x_{t-1}^*)\|_* \|x_{t-1}^* - x_t^*\| \\ & \quad - \lambda_2 D_h(x_{t-1}^* || x_{t-1}) - \frac{m}{2} \|x_{t-1}^* - x_{t-1}\|^2 \end{aligned} \quad (31c)$$

$$\begin{aligned} & \leq \lambda_1 D_h(x_t^* || x_{t-1}^*) + \frac{\lambda_2 \beta + m}{2\beta^2} \|\nabla h(x_{t-1}) - \nabla h(x_{t-1}^*)\|_*^2 + \frac{\lambda_1^2 \beta^2}{2(\lambda_2 \beta + m)} \|x_{t-1}^* - x_t^*\|^2 \\ & \quad - \lambda_2 D_h(x_{t-1}^* || x_{t-1}) - \frac{m}{2} \|x_{t-1}^* - x_{t-1}\|^2 \\ & = \lambda_1 D_h(x_t^* || x_{t-1}^*) + \frac{\lambda_1^2 \beta^2}{2(\lambda_2 \beta + m)} \|x_{t-1}^* - x_t^*\|^2 \\ & \quad + \left(\frac{\lambda_2}{2\beta} \|\nabla h(x_{t-1}) - \nabla h(x_{t-1}^*)\|_*^2 - \lambda_2 D_h(x_{t-1}^* || x_{t-1}) \right) \\ & \quad + \left(\frac{m}{2\beta^2} \|\nabla h(x_{t-1}) - \nabla h(x_{t-1}^*)\|_*^2 - \frac{m}{2} \|x_{t-1}^* - x_{t-1}\|^2 \right) \end{aligned} \quad (31d)$$

$$\begin{aligned} & \leq \lambda_1 D_h(x_t^* || x_{t-1}^*) + \frac{\lambda_1^2 \beta^2}{2(\lambda_2 \beta + m)} \|x_{t-1}^* - x_t^*\|^2 \\ & \leq \lambda_1 \left(1 + \frac{\lambda_1 \beta^2}{\alpha(\lambda_2 \beta + m)} \right) D_h(x_t^* || x_{t-1}^*), \end{aligned} \quad (31e)$$

745 where we use Lemma 5 in line (31a); Cauchy-Schwartz inequality in line (31b); the AM-GM
 746 inequality in the line (31c); inequalities (29) and (30) in line (31d); and inequality (25) in line (31e).
 747 Substituting inequality (31) into inequality (28), we obtain

$$H_t + \lambda_1 M_t + \Delta\phi \leq (H_t^* + \lambda_2 D_h(x_t^* || v_t)) + \lambda_1 \left(1 + \frac{\lambda_1 \beta^2}{\alpha(\lambda_2 \beta + m)}\right) M_t^*.$$

748 Using inequality (25) and the fact that f_t is m -strongly convex, we obtain

$$\lambda_2 D_h(x_t^* || v_t) \leq \frac{\lambda_2 \beta}{2} \|x_t^* - v_t\|^2 \leq \frac{\lambda_2 \beta}{m} H_t^*.$$

749 Therefore we have

$$H_t + \lambda_1 M_t + \Delta\phi \leq (1 + \frac{\lambda_2 \beta}{m}) H_t^* + \lambda_1 \left(1 + \frac{\lambda_1 \beta^2}{\alpha(\lambda_2 \beta + m)}\right) M_t^*.$$

750 Since $0 < \lambda_1 \leq 1$, we have

$$H_t + M_t + \frac{1}{\lambda_1} \Delta\phi \leq \frac{H_t + \lambda_1 M_t + \Delta\phi}{\lambda_1} \leq \frac{m + \lambda_2 \beta}{m \lambda_1} H_t^* + \left(1 + \frac{\beta^2}{\alpha} \cdot \frac{\lambda_1}{\lambda_2 \beta + m}\right) M_t^*.$$

751 Theorem 4 follows from summing the above inequality over all timesteps t .

752 E R-OBD with Squared ℓ_2 Norm

753 When $h(x) = \frac{1}{2} \|x\|_2^2$, the Bregman Divergence $D_h(x || y)$ is equal to the squared ℓ_2 norm $\frac{1}{2} \|x - y\|_2^2$.
 754 Hence, setting $h(x) = \frac{1}{2} \|x\|_2^2$ in Algorithm 3 gives us R-OBD in the squared ℓ_2 setting. In this
 755 section, we present a separate proof of Regularized OBD with squared ℓ_2 norm, in order to remove
 756 the assumption that the hitting costs $\{f_t\}$ are differentiable.

757 **Theorem 7.** Consider hitting cost functions that are m -strongly convex with respect to ℓ_2 norm and
 758 movement costs given by $\frac{1}{2} \|x_t - x_{t-1}\|_2^2$. There exists a choice λ_1, λ_2 such that the competitive ratio
 759 of Regularized OBD matches the lower bound proved in Theorem 1, i.e. the competitive ratio is at
 760 most $\frac{1}{2} \left(1 + \sqrt{1 + \frac{4}{m}}\right)$.

761 This result follows from the more general bound in Theorem 8 below, which describes the competitive
 762 ratio of Algorithm 3 as a function of λ_1, λ_2 .

763 **Theorem 8.** Consider hitting cost functions that are m -strongly convex with respect to ℓ_2 norm and
 764 movement costs given by $\frac{1}{2} \|x_t - x_{t-1}\|_2^2$. Regularized-OBD (Algorithm 3 with $h(x) = \frac{1}{2} \|x\|_2^2$) with
 765 parameters $1 \geq \lambda_1 > 0, \lambda_2 \geq 0$ has competitive ratio at most

$$\max \left(\frac{m + \lambda_2}{\lambda_1} \cdot \frac{1}{m}, 1 + \frac{\lambda_1}{\lambda_2 + m} \right).$$

766 Notice that Theorem 7 follows immediately by setting $\frac{m + \lambda_2}{\lambda_1} = \frac{m}{2} \left(1 + \sqrt{1 + \frac{4}{m}}\right)$ in Theorem 8.

767 Before proving Theorem 8, we first prove a technical lemma which gives a lower bound of the value
 768 of hitting cost as a function of the distance to the minimizer.

769 **Lemma 13.** If $f : \mathcal{X} \rightarrow \mathbb{R}$ is a m -strongly convex function with respect to some norm $\|\cdot\|$, and v is
 770 the minimizer of f (i.e. $v = \arg \min_{x \in \mathcal{X}} f(x)$), then we have $\forall x \in \mathcal{X}$,

$$f(x) \geq f(v) + \frac{m}{2} \|x - v\|^2.$$

771 *Proof.* By the definition of m -strongly convex, we obtain that $\forall \alpha \in (0, 1)$,

$$f(\alpha x + (1 - \alpha)v) \leq \alpha f(x) + (1 - \alpha)f(v) - \frac{m}{2} \alpha(1 - \alpha) \|x - v\|^2. \quad (32)$$

772 Notice that $f(v) \leq f(\alpha x + (1-\alpha)v)$. Combining this with inequality (32), we obtain that $\forall \alpha \in (0, 1)$,

$$f(v) \leq \alpha f(x) + (1-\alpha)f(v) - \frac{m}{2}\alpha(1-\alpha)\|x-v\|^2.$$

773 Rearranging the terms, we observe that $\forall \alpha \in (0, 1)$,

$$f(x) \geq f(v) + \frac{m}{2}(1-\alpha)\|x-v\|^2.$$

774 Therefore

$$f(x) \geq \lim_{\alpha \rightarrow 0^+} \left(f(v) + \frac{m}{2}(1-\alpha)\|x-v\|^2 \right) = f(v) + \frac{m}{2}\|x-v\|^2.$$

775

□

776 Now we return to the proof of Theorem 8.

777 *Proof of Theorem 8.* In the proof, we use the property of strongly convex to derive an inequality in
 778 the form of $H_t + M_t + \Delta\phi \leq C(H_t^* + M_t^*)$, where $\Delta\phi$ is the change in potential and C is an upper
 779 bound for the competitive ratio.

780 Throughout the proof, we use $\|\cdot\|$ to denote ℓ_2 norm.

781 Notice that when $h(x) = \frac{1}{2}\|x\|^2$, the update rule in Algorithm 3 is:

$$x_t = \arg \min_x f_t(x) + \frac{\lambda_1}{2}\|x - x_{t-1}\|^2 + \frac{\lambda_2}{2}\|x - v_t\|^2.$$

782 For convenience, we define

$$F_t(x) = f_t(x) + \frac{\lambda_1}{2}\|x - x_{t-1}\|^2 + \frac{\lambda_2}{2}\|x - v_t\|^2.$$

783 Since $f_t(x)$ is m -strongly convex, $\frac{\lambda_1}{2}\|x - x_{t-1}\|^2$ is λ_1 -strongly convex, and $\frac{\lambda_2}{2}\|x - v_t\|^2$ is λ_2 -
 784 strongly convex, $F_t(x)$ is $(m + \lambda_1 + \lambda_2)$ -strongly convex. Since $x_t = \arg \min_x F_t(x)$, by Lemma
 785 13, we obtain

$$F_t(x_t^*) \geq F_t(x_t) + \frac{m + \lambda_1 + \lambda_2}{2}\|x_t^* - x_t\|^2,$$

786 which implies

$$\begin{aligned} & H_t + \lambda_1 M_t + \frac{m + \lambda_1 + \lambda_2}{2}\|x_t^* - x_t\|^2 \\ & \leq H_t + \lambda_1 M_t + \frac{\lambda_2}{2}\|x - v_t\|^2 + \frac{m + \lambda_1 + \lambda_2}{2}\|x_t^* - x_t\|^2 \\ & \leq H_t^* + \frac{\lambda_1}{2}\|x_t^* - x_{t-1}\|^2 + \frac{\lambda_2}{2}\|x_t^* - v_t\|^2. \end{aligned} \tag{33}$$

787 We define the potential function as $\phi(x_t, x_t^*) = \frac{m + \lambda_1 + \lambda_2}{2}\|x_t^* - x_t\|^2$ and $\Delta\phi = \phi(x_t, x_t^*) -$
 788 $\phi(x_{t-1}, x_{t-1}^*)$. We then can rewrite inequality (33) as

$$H_t + \lambda_1 M_t + \Delta\phi \leq \left(H_t^* + \frac{\lambda_2}{2}\|x_t^* - v_t\|^2 \right) + \frac{\lambda_1}{2}\|x_t^* - x_{t-1}\|^2 - \frac{m + \lambda_1 + \lambda_2}{2}\|x_{t-1}^* - x_{t-1}\|^2. \tag{34}$$

789 Additionally

$$\begin{aligned} & \frac{\lambda_1}{2} \|x_t^* - x_{t-1}\|^2 - \frac{m + \lambda_1 + \lambda_2}{2} \|x_{t-1}^* - x_{t-1}\|^2 \\ & \leq \frac{\lambda_1}{2} (\|x_t^* - x_{t-1}^*\| + \|x_{t-1}^* - x_{t-1}\|)^2 - \frac{m + \lambda_1 + \lambda_2}{2} \|x_{t-1}^* - x_{t-1}\|^2 \end{aligned} \quad (35a)$$

$$\begin{aligned} & = \frac{\lambda_1}{2} \|x_t^* - x_{t-1}^*\|^2 + \lambda_1 \|x_t^* - x_{t-1}^*\| \cdot \|x_{t-1}^* - x_{t-1}\| - \frac{m + \lambda_2}{2} \|x_{t-1}^* - x_{t-1}\|^2 \\ & \leq \frac{\lambda_1}{2} \|x_t^* - x_{t-1}^*\|^2 + \frac{\lambda_1^2}{2(m + \lambda_2)} \|x_t^* - x_{t-1}^*\|^2 + \frac{m + \lambda_2}{2} \|x_{t-1}^* - x_{t-1}\|^2 \\ & \quad - \frac{m + \lambda_2}{2} \|x_{t-1}^* - x_{t-1}\|^2 \end{aligned} \quad (35b)$$

$$\begin{aligned} & = \frac{\lambda_1(\lambda_1 + \lambda_2 + m)}{2(\lambda_2 + m)} \|x_{t-1}^* - x_{t-1}^*\|^2 \\ & = \lambda_1 \left(1 + \frac{\lambda_1}{\lambda_2 + m} \right) M_t^*, \end{aligned}$$

790 where we apply the triangle inequality in line (35a) and AM-GM in line (35b).

791 Combining inequalities (34) and (35), we obtain

$$H_t + \lambda_1 M_t + \Delta\phi \leq \left(H_t^* + \frac{\lambda_2}{2} \|x_t^* - v_t\|^2 \right) + \lambda_1 \left(1 + \frac{\lambda_1}{\lambda_2 + m} \right) M_t^*. \quad (36)$$

792 And since $f_t(x)$ is m -strongly convex, we have

$$\frac{\lambda_2}{2} \|x_t^* - v_t\|^2 \leq \frac{\lambda_2}{m} H_t^*.$$

793 Substituting the above identity into inequality (36) yields

$$H_t + \lambda_1 M_t + \Delta\phi \leq \frac{m + \lambda_2}{m} H_t^* + \lambda_1 \left(1 + \frac{\lambda_1}{m + \lambda_2} \right) M_t^*. \quad (37)$$

794 Using inequality (37), we obtain

$$H_t + M_t + \frac{1}{\lambda_1} \Delta\phi \leq \frac{H_t + \lambda_1 M_t + \Delta\phi}{\lambda_1} \leq \frac{m + \lambda_2}{\lambda_1 m} H_t^* + \left(1 + \frac{\lambda_1}{m + \lambda_2} \right) M_t^*.$$

795 Theorem 8 follows from summing the above inequality over all timesteps t . \square

796 **F Proof of Theorem 5**

797 In this proof, we construct counterexamples for two separate cases, based on whether λ_1 is larger or
798 smaller than m . Recall that $\lambda_2 = 0$ throughout the proof.

799 **Case 1:** $\lambda_1 > m$

800 In this case, we show the competitive ratio can be unbounded by proposing a series of identical
801 hitting cost functions on the real number line. We construct a hitting cost function f with minimizer
802 v so that there exists a fixed point $K \neq v$ (i.e. when $x_{t-1} = K$ and $f_t = f$, the algorithm selects
803 $x_t = x_{t-1}$). Since R-OBD is independent of timestep t , we can propose $f_t = f$ for $t = 1, 2, \dots, T$
804 and let $x_0 = K$. In this scenario, the total cost of R-OBD grows linearly in T . However, by choosing
805 $x_1 = x_2 = \dots = x_T = v$, the total cost incurred by the offline adversary is a constant. Therefore the
806 competitive ratio of R-OBD will be unbounded.

807 Specifically, consider the hitting cost function

$$f(x) = \begin{cases} \frac{m}{2} (1 - (x + 1)^2) & -1 \leq x \leq 0 \\ \frac{m}{2} x^2 & \text{otherwise} \end{cases}.$$

808 Suppose $x_{t-1} = -1$, then R-OBDD will choose x_t such that

$$x_t = \arg \min_x f(x) + \frac{\lambda_1}{2}(x+1)^2.$$

809 Notice that

$$f(x) + \frac{\lambda_1}{2}(x+1)^2 = \begin{cases} \frac{m}{2} + \frac{\lambda_1-m}{2}(x+1)^2 & -1 \leq x \leq 0 \\ \frac{m}{2}x^2 + \frac{\lambda_1}{2}(x+1)^2 & \text{otherwise} \end{cases}.$$

810 Since $\lambda_1 > m$, we see that the quantity above is $\geq \frac{m}{2}$ for all real x , where equality only holds when
811 $x = -1$. It follows that $x_t = x_{t-1} = -1 \neq 0 = v$. Thus $K = -1$ is a fixed point satisfying the
812 requirements described as above.

813 **Case 2:** $\lambda_1 \leq m$

814 We consider a situation such that the R-OBDD algorithm moves far away from the starting point,
815 incurring significant movement cost, whereas the offline adversary could pay relatively little cost by
816 staying at the starting point. More specifically, suppose the starting point $x_0 = 0$ and the first hitting
817 cost function is $f_1(x) = \frac{m}{2}(1-x)^2$. Consider an adversary which chooses $x_0 = x_1 = \dots = x_T$.
818 The cost incurred by the adversary is

$$\text{cost}(ADV) = \frac{m}{2}.$$

819 Using the update rule, the R-OBDD algorithm chooses

$$x_1 = \arg \min_x \frac{m}{2}(1-x)^2 + \frac{\lambda_1}{2}x^2 = \frac{m}{m+\lambda_1} \geq \frac{1}{2}.$$

820 The movement cost incurred by R-OBDD is at least

$$\text{cost}(ALG) \geq M_1 = \frac{1}{2}x_1^2 \geq \frac{1}{8}.$$

821 Thus the competitive ratio is at least

$$\frac{\text{cost}(ALG)}{\text{cost}(ADV)} \geq \frac{1}{4m}.$$

822 Theorem 5 follows from combining these two cases.

823 G Proof of Theorem 6

824 Let $\{x_t^L\}$ be the sequence of points achieving the L -constrained offline optimal. We first prove an
825 upper bound on the difference of hitting costs $f_t(x_t) - f_t(x_t^L)$, and then use this bound to prove a
826 $O(G\sqrt{TL})$ upper bound on the regret $\sum_{t=1}^T (f_t(x_t) - f_t(x_t^L) + c(x_t, x_{t-1})) - \sum_{t=1}^T c(x_t^L, x_{t-1}^L)$.

827 Since the function $f_t(x) + \lambda_1 D_h(x||x_{t-1}) + \lambda_2 D_h(x||v_t)$ is strongly convex, it has a unique mini-
828 mizer, at which point the gradient vanishes. This is the point x_t which Algorithm 3 picks in round t .
829 We can rearrange the vanishing gradient condition to obtain

$$\nabla f_t(x_t) = \lambda_1 (\nabla h(x_{t-1}) - \nabla h(x_t)) + \lambda_2 (\nabla h(v_t) - \nabla h(x_t)).$$

830 Therefore by Lemma 5, we have

$$\begin{aligned} \langle \nabla f_t(x_t), x_t - x_t^L \rangle &= \lambda_1 \langle \nabla h(x_{t-1}) - \nabla h(x_t), x_t - x_t^L \rangle + \lambda_2 \langle \nabla h(v_t) - \nabla h(x_t), x_t - x_t^L \rangle \\ &= \lambda_1 (D_h(x_t^L||x_{t-1}) - D_h(x_t^L||x_t) - D_h(x_t||x_{t-1})) \\ &\quad + \lambda_2 (D_h(x_t^L||v_t) - D_h(x_t^L||x_t) - D_h(x_t||v_t)). \end{aligned} \tag{38}$$

831 Recall that h is α -strongly convex and β -strongly smooth with respect to the norm $\|\cdot\|$, hence

$$\forall x, y, \frac{\alpha}{2} \|x - y\|^2 \leq D_h(x||y) \leq \frac{\beta}{2} \|x - y\|^2. \tag{39}$$

832 Therefore

$$D_h(x_t^L||v_t) - D_h(x_t^L||x_t) - D_h(x_t||v_t) \leq D_h(x_t^L||v_t) \leq \frac{\beta}{2} \|x_t^L - v_t\|^2 \leq \frac{\beta D^2}{2}.$$

833 In light of equation (38), we obtain

$$\langle \nabla f_t(x_t), x_t - x_t^L \rangle \leq \lambda_1 (D_h(x_t^L||x_{t-1}) - D_h(x_t^L||x_t) - D_h(x_t||x_{t-1})) + \frac{\beta D^2}{2} \cdot \lambda_2. \quad (40)$$

834 Let $q > 0$ be a parameter which we will pick later. For all $q > 0$, it holds that

$$\begin{aligned} & f_t(x_t) - f_t(x_t^L) \\ & \leq \langle \nabla f_t(x_t), x_t - x_t^L \rangle - \frac{m}{2} \|x_t - x_t^L\|^2 \end{aligned} \quad (41a)$$

$$\begin{aligned} & \leq \lambda_1 (D_h(x_t^L||x_{t-1}) - D_h(x_t^L||x_t) - D_h(x_t||x_{t-1})) - \frac{m}{2} \|x_t - x_t^L\|^2 + \frac{\beta D^2}{2} \cdot \lambda_2 \quad (41b) \\ & = (\lambda_1 + q) (D_h(x_t^L||x_{t-1}) - D_h(x_t^L||x_t)) - \lambda_1 D_h(x_t||x_{t-1}) \\ & \quad - \left(q D_h(x_t^L||x_{t-1}) - q D_h(x_t^L||x_t) + \frac{m}{2} \|x_t - x_t^L\|^2 \right) \\ & \quad + \frac{\beta D^2}{2} \cdot \lambda_2. \end{aligned}$$

835 where we apply strong convexity in line (41a), and equation (40) in line (41b). Using Lemma 5, we
836 obtain

$$\begin{aligned} & q D_h(x_t^L||x_{t-1}) - q D_h(x_t^L||x_t) + \frac{m}{2} \|x_t - x_t^L\|^2 \\ & = q D_h(x_t||x_{t-1}) + q \langle \nabla h(x_{t-1}) - \nabla h(x_t), x_t - x_t^L \rangle + \frac{m}{2} \|x_t - x_t^L\|^2 \\ & \geq q D_h(x_t||x_{t-1}) - q \|\nabla h(x_{t-1}) - \nabla h(x_t)\|_* \|x_t - x_t^L\| + \frac{m}{2} \|x_t - x_t^L\|^2 \end{aligned} \quad (42a)$$

$$\geq q D_h(x_t||x_{t-1}) - \left(\frac{q^2}{2m} \|\nabla h(x_{t-1}) - \nabla h(x_t)\|_*^2 + \frac{m}{2} \|x_t - x_t^L\|^2 \right) + \frac{m}{2} \|x_t - x_t^L\|^2 \quad (42b)$$

$$\begin{aligned} & = q D_h(x_t||x_{t-1}) - \frac{q^2}{2m} \|\nabla h(x_{t-1}) - \nabla h(x_t)\|_*^2 \\ & \geq q D_h(x_t||x_{t-1}) - \frac{\beta q^2}{m} D_h(x_t||x_{t-1}) \quad (42c) \\ & = \left(q - \frac{\beta q^2}{m} \right) D_h(x_t||x_{t-1}), \end{aligned}$$

837 where we apply the Cauchy-Schwartz inequality in line (42a), the AM-GM inequality in line (42b),
838 and Lemma 1 in line (42c).

839 In order to maximize the coefficient $\left(q - \frac{\beta q^2}{m} \right)$, we set $q = \frac{m}{2\beta}$. By substituting inequality (42) into
840 inequality (41), we obtain

$$\begin{aligned} & f_t(x_t) - f_t(x_t^L) \\ & \leq \left(\lambda_1 + \frac{m}{2\beta} \right) (D_h(x_t^L||x_{t-1}) - D_h(x_t^L||x_t)) - \left(\lambda_1 + \frac{m}{4\beta} \right) D_h(x_t||x_{t-1}) + \frac{\beta D^2}{2} \cdot \lambda_2. \end{aligned} \quad (43)$$

841 Using the condition $\lambda_1 + \frac{m}{4\beta} \geq 1$, we observe that

$$f_t(x_t) - f_t(x_t^L) + D_h(x_t||x_{t-1}) \left(\lambda_1 + \frac{m}{2\beta} \right) (D_h(x_t^L||x_{t-1}) - D_h(x_t^L||x_t)) + \frac{\beta D^2}{2} \cdot \lambda_2. \quad (44)$$

842 Notice that

$$\sum_{t=1}^T \|x_t^L - x_{t+1}^L\| \leq \sqrt{T \left(\sum_{t=1}^T \|x_t^L - x_{t+1}^L\|^2 \right)} \leq \sqrt{T \left(\sum_{t=1}^T \frac{2D_h(x_{t+1}^L||x_t^L)}{\alpha} \right)} \leq \sqrt{\frac{2TL}{\alpha}}. \quad (45)$$

843 where we use the generalized mean inequality in the first step and α -strong convexity of h in the
 844 second step (cf. equation (39)). By Lemma 6, we can give the following upper bound:

$$\begin{aligned}
 & \sum_{t=1}^T D_h(x_t^L || x_{t-1}) - D_h(x_t^L || x_t) \\
 &= \sum_{t=1}^T (D_h(0 || x_{t-1}) - D_h(0 || x_t) + \langle \nabla h(x_t) - \nabla h(x_{t-1}), x_t^L \rangle) \\
 &= D_h(0 || x_0) - D_h(0 || x_T) + \sum_{t=1}^{T-1} \langle \nabla h(x_t), x_t^L - x_{t+1}^L \rangle - \langle \nabla h(x_0), x_1^L \rangle + \langle \nabla h(x_T), x_T^L \rangle \\
 &\leq \sum_{t=1}^T \langle \nabla h(x_t), x_t^L - x_{t+1}^L \rangle \tag{46a}
 \end{aligned}$$

$$\leq \sum_{t=1}^T \|\nabla h(x_t)\|_* \|x_t^L - x_{t+1}^L\| \tag{46b}$$

$$\begin{aligned}
 &\leq G \sum_{t=1}^T \|x_t^L - x_{t+1}^L\| \\
 &\leq G \sqrt{\frac{2TL}{\alpha}}, \tag{46c}
 \end{aligned}$$

845 where we use the facts $x_0 = x_0^L = x_{T+1}^L = 0, \nabla h(0) = 0$ in line (46a), the Cauchy-Schwartz
 846 inequality in line (46b), and inequality (45) in line (46c).

847 Therefore we obtain

$$\begin{aligned}
 & \text{cost}(OBD) - \text{cost}(OPT(L)) \\
 &= \sum_{t=1}^T (f_t(x_t) + D_h(x_t || x_{t-1})) - (f_t(x_t^L) + D_h(x_t^L || x_{t-1}^L)) \tag{47a}
 \end{aligned}$$

$$\begin{aligned}
 &\leq \left(\sum_{t=1}^T f_t(x_t) - f_t(x_t^L) + D_h(x_t || x_{t-1}) \right) - L \\
 &\leq \left(\lambda_1 + \frac{m}{2\beta} \right) G \sqrt{\frac{2TL}{\alpha}} + T \cdot \frac{\beta D^2}{2} \cdot \lambda_2 - L, \tag{47b}
 \end{aligned}$$

848 where we use the definition of $OPT(L)$ in line (47a); inequalities (44) and (46) in line (47b).

849 Since by assumption we have $G < \infty, \lambda_2 = \eta(T, L, D, G) \leq \frac{KG}{D^2} \cdot \sqrt{\frac{L}{T}}$ for some constant K , by
 850 inequality (47), we obtain

$$\text{cost}(OBD) - \text{cost}(OPT(L)) = O(G\sqrt{TL}),$$

851 which completes the proof.