
Accelerated Stochastic Greedy Coordinate Descent by Soft Thresholding Projection onto Simplex

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Abstract

In this paper we study the well-known greedy coordinate descent (GCD) algorithm to solve ℓ_1 -regularized problems and improve GCD by the two popular strategies: Nesterov’s acceleration and stochastic optimization. Firstly, based on an ℓ_1 -norm square approximation, we propose a new rule for greedy selection which is non-trivial to solve but convex; then an efficient algorithm called “SOft ThreshOlding PrOjection (SOTOPO)” is proposed to exactly solve an ℓ_1 -regularized ℓ_1 -norm square approximation problem, which is induced by the new rule. Based on the new rule and the SOTOPO algorithm, the Nesterov’s acceleration and stochastic optimization strategies are then successfully applied to the GCD algorithm. The resulted algorithm called accelerated stochastic greedy coordinate descent (ASGCD) has the optimal convergence rate $O(\sqrt{1/\epsilon})$; meanwhile, it reduces the iteration complexity of greedy selection up to a factor of sample size. Both theoretically and empirically, we show that ASGCD has better performance for high-dimensional and dense problems with sparse solutions.

1 Introduction

In large-scale convex optimization, first-order methods are widely used due to their cheap iteration cost. In order to improve the convergence rate and reduce the iteration cost further, two important strategies are used in first-order methods: Nesterov’s acceleration and stochastic optimization. Nesterov’s acceleration is referred to the technique that uses some algebra trick to accelerate first-order algorithms; while stochastic optimization is referred to the method that samples one training example or one dual coordinate at random from the training data in each iteration. Assume the objective function $F(x)$ is convex and smooth. Let $F^* = \min_{x \in R^d} F(x)$ be the optimal value. In order to find an approximate solution x that satisfies $F(x) - F^* \leq \epsilon$, the vanilla gradient descent method needs $O(1/\epsilon)$ iterations. While after applying the Nesterov’s acceleration scheme [18], the resulted accelerated full gradient method (AFG) [18] only needs $O(\sqrt{1/\epsilon})$ iterations, which is optimal for first-order algorithms [18]. Meanwhile, assume $F(x)$ is also a finite sum of n sample convex functions. By sampling one training example, the resulted stochastic gradient descent (SGD) and its variants [15, 25, 1] can reduce the iteration complexity by a factor of the sample size. As an alternative of SGD, randomized coordinate descent (RCD) can also reduce the iteration complexity by a factor of the sample size [17] and obtain the optimal convergence rate $O(\sqrt{1/\epsilon})$ by Nesterov’s acceleration [16, 14]. The development of gradient descent and RCD raises an interesting problem: can the Nesterov’s acceleration and stochastic optimization strategies be used to improve other existing first-order algorithms?

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In this paper, we answer this question partly by studying coordinate descent with Gauss-Southwell selection, i.e., greedy coordinate descent (GCD). GCD is widely used for solving sparse optimization problems in machine learning [24, 11, 19]. If an optimization problem has a sparse solution, it is more suitable than its counterpart RCD. However, the theoretical convergence rate is still $O(1/\epsilon)$. Meanwhile if the iteration complexity is comparable, GCD will be preferable than RCD [19]. However in the general case, in order to do exact Gauss-Southwell selection, computing the full gradient beforehand is necessary, which causes GCD has much higher iteration complexity than RCD. To be concrete, in this paper we consider the well-known nonsmooth ℓ_1 -regularized problem:

$$\min_{x \in \mathbb{R}^d} \left\{ F(x) \stackrel{\text{def}}{=} f(x) + \lambda \|x\|_1 \stackrel{\text{def}}{=} \frac{1}{n} \sum_{j=1}^n f_j(x) + \lambda \|x\|_1 \right\}, \quad (1)$$

where $\lambda \geq 0$ is a regularization parameter, $f(x) = \frac{1}{n} \sum_{j=1}^n f_j(x)$ is a smooth convex function that is a finite average of n smooth convex function $f_j(x)$. Given samples $\{(a_1, b_1), (a_2, b_2), \dots, (a_n, b_n)\}$ with $a_j \in \mathbb{R}^d$ ($j \in [n] \stackrel{\text{def}}{=} \{1, 2, \dots, n\}$), if each $f_j(x) = f_j(a_j^T x, b_j)$, then (1) is an ℓ_1 -regularized empirical risk minimization (ℓ_1 -ERM) problem. For example, if $b_j \in \mathbb{R}$ and $f_j(x) = \frac{1}{2}(b_j - a_j^T x)^2$, (1) is Lasso; if $b_j \in \{-1, 1\}$ and $f_j(x) = \log(1 + \exp(-b_j a_j^T x))$, ℓ_1 -regularized logistic regression is obtained.

In the above nonsmooth case, the Gauss-Southwell rule has 3 different variants [19, 24]: GS- s , GS- r and GS- q . The GCD algorithm with all the 3 rules can be viewed as the following procedure: in each iteration based on a quadratic approximation of $f(x)$ in (1), one minimizes a surrogate objective function under the constraint that the direction vector used for update has at most 1 nonzero entry. The resulted problems under the 3 rules are easy to solve but are *nonconvex* due to the cardinality constraint of direction vector. While when using Nesterov's acceleration scheme, convexity is needed for the derivation of the optimal convergence rate $O(\sqrt{1/\epsilon})$ [18]. Therefore, it is impossible to accelerate GCD by the Nesterov's acceleration scheme under the 3 existing rules.

In this paper, we propose a novel variant of Gauss-Southwell rule by using an ℓ_1 -norm square approximation of $f(x)$ rather than quadratic approximation. The new rule involves an ℓ_1 -regularized ℓ_1 -norm square approximation problem, which is nontrivial to solve but is *convex*. To *exactly* solve the challenging problem, we propose an efficient SOft ThreshOlding PrOjection (SOTOPO) algorithm. The SOTOPO algorithm has $O(d + |Q| \log |Q|)$ cost, where it is often the case $|Q| \ll d$. The complexity result $O(d + |Q| \log |Q|)$ is better than $O(d \log d)$ of its counterpart SOPOPO [20], which is an Euclidean projection method.

Then based on the new rule and SOTOPO, we accelerate GCD to attain the optimal convergence rate $O(\sqrt{1/\epsilon})$ by combing a delicately selected mirror descent step. Meanwhile, we show that it is not necessary to compute full gradient beforehand: sampling one training example and computing a noisy gradient rather than full gradient is enough to perform greedy selection. This stochastic optimization technique reduces the iteration complexity of greedy selection by a factor of the sample size. The final result is an accelerated stochastic greedy coordinate descent (ASGCD) algorithm.

Assume x^* is an optimal solution of (1). Assume that each $f_j(x)$ (for all $j \in [n]$) is L_p -smooth *w.r.t.* $\|\cdot\|_p$ ($p = 1, 2$), i.e., for all $x, y \in \mathbb{R}^d$,

$$\|\nabla f_j(x) - \nabla f_j(y)\|_q \leq L_p \|x - y\|_p, \quad (2)$$

where if $p = 1$, then $q = \infty$; if $p = 2$, then $q = 2$.

In order to find an x that satisfies $F(x) - F(x^*) \leq \epsilon$, ASGCD needs $O\left(\frac{\sqrt{CL_1}\|x^*\|_1}{\sqrt{\epsilon}}\right)$ iterations (see (16)), where C is a function of d that varies slowly over d and is upper bounded by $\log^2(d)$. For high-dimensional and dense problems with sparse solutions, ASGCD has better performance than the state of the art. Experiments demonstrate the theoretical result.

Notations: Let $[d]$ denote the set $\{1, 2, \dots, d\}$. Let \mathbb{R}_+ denote the set of nonnegative real number. For $x \in \mathbb{R}^d$, let $\|x\|_p = (\sum_{i=1}^d |x_i|^p)^{\frac{1}{p}}$ ($1 \leq p < \infty$) denote the ℓ_p -norm and $\|x\|_\infty = \max_{i \in [d]} |x_i|$ denote the ℓ_∞ -norm of x . For a vector x , let $\dim(x)$ denote the dimension of x ; let x_i denote the i -th element of x . For a gradient vector $\nabla f(x)$, let $\nabla_i f(x)$ denote the i -th element of $\nabla f(x)$. For a set S , let $|S|$ denote the cardinality of S . Denote the simplex $\Delta_d = \{\theta \in \mathbb{R}_+^d : \sum_{i=1}^d \theta_i = 1\}$.

2 The SOTOPO algorithm

The proposed SOTOPO algorithm aims to solve the proposed new rule, *i.e.*, minimize the following ℓ_1 -regularized ℓ_1 -norm square approximation problem,

$$\tilde{h} \stackrel{\text{def}}{=} \arg \min_{g \in \mathbb{R}^d} \left\{ \langle \nabla f(x), g \rangle + \frac{1}{2\eta} \|g\|_1^2 + \lambda \|x + g\|_1 \right\}, \quad (3)$$

$$\tilde{x} \stackrel{\text{def}}{=} x + \tilde{h}, \quad (4)$$

where x denotes the current iteration, η a step size, g the variable to optimize, \tilde{h} the director vector for update and \tilde{x} the next iteration. The number of nonzero entries of \tilde{h} denotes how many coordinates will be updated in this iteration. Unlike the quadratic approximation used in GS- s , GS- r and GS- q rules, in the new rule the coordinate(s) to update is implicitly selected by the sparsity-inducing property of the ℓ_1 -norm square $\|g\|_1^2$ rather than using the cardinality constraint $\|g\|_0 \leq 1$ (*i.e.*, g has at most 1 nonzero element) [19, 24]. By [8, §9.4.2], when the nonsmooth term $\lambda \|x + g\|_1$ in (1) does not exist, the minimizer of the ℓ_1 -norm square approximation (*i.e.*, ℓ_1 -norm steepest descent) is equivalent to GCD. When $\lambda \|x + g\|_1$ exists, generally, there may be one or more coordinates to update in this new rule. Because of the sparsity-inducing property of $\|g\|_1^2$ and $\|x + g\|_1$, both the direction vector \tilde{h} and the iterative solution \tilde{x} are sparse. In addition, (3) is an unconstrained problem and thus is feasible.

2.1 A variational reformulation and its properties

(3) involves the nonseparable, nonsmooth term $\|g\|_1^2$ and the nonsmooth term $\|x + g\|_1$. Because there are two nonsmooth terms, it seems difficult to solve (3) directly. While by the variational identity $\|g\|_1^2 = \inf_{\theta \in \Delta_d} \sum_{i=1}^d \frac{g_i^2}{\theta_i}$ in [5]², in Lemma 1, it is shown that we can transform the original nonseparable and nonsmooth problem into a separable and smooth optimization problem on a simplex.

Lemma 1. *By defining*

$$J(g, \theta) \stackrel{\text{def}}{=} \langle \nabla f(x), g \rangle + \frac{1}{2\eta} \sum_{i=1}^d \frac{g_i^2}{\theta_i} + \lambda \|x + g\|_1, \quad (5)$$

$$\tilde{g}(\theta) \stackrel{\text{def}}{=} \arg \min_{g \in \mathbb{R}^d} J(g, \theta), \quad J(\theta) \stackrel{\text{def}}{=} J(\tilde{g}(\theta), \theta), \quad (6)$$

$$\tilde{\theta} \stackrel{\text{def}}{=} \arg \inf_{\theta \in \Delta_d} J(\theta), \quad (7)$$

where $\tilde{g}(\theta)$ is a vector function. Then the minimization problem to find \tilde{h} in (3) is equivalent to the problem (7) to find $\tilde{\theta}$ with the relation $\tilde{h} = \tilde{g}(\tilde{\theta})$. Meanwhile, $\tilde{g}(\theta)$ and $J(\theta)$ in (6) are both coordinate separable with the expressions

$$\forall i \in [d], \tilde{g}_i(\theta) = \tilde{g}_i(\theta_i) \stackrel{\text{def}}{=} \text{sign}(x_i - \theta_i \eta \nabla_i f(x)) \cdot \max\{0, |x_i - \theta_i \eta \nabla_i f(x)| - \theta_i \eta \lambda\} - x_i, \quad (8)$$

$$J(\theta) = \sum_{i=1}^d J_i(\theta_i), \quad \text{where} \quad J_i(\theta_i) \stackrel{\text{def}}{=} \nabla_i f(x) \cdot \tilde{g}_i(\theta_i) + \frac{1}{2\eta} \sum_{i=1}^d \frac{\tilde{g}_i^2(\theta_i)}{\theta_i} + \lambda |x_i + \tilde{g}_i(\theta_i)|. \quad (9)$$

In Lemma 1, (8) is obtained by the iterative soft thresholding operator [7]. By Lemma 1, we can reformulate (3) into the problem (5), which is about two parameters g and θ . Then by the joint convexity, we swap the optimization order of g and θ . Fixing θ and optimizing with respect to (*w.r.t.*) g , we can get a closed form of $\tilde{g}(\theta)$, which is a vector function about θ . Substituting $\tilde{g}(\theta)$ into $J(g, \theta)$, we get the problem (7) about θ . Finally, the optimal solution \tilde{h} in (3) can be obtained by $\tilde{h} = \tilde{g}(\tilde{\theta})$.

The explicit expression of each $J_i(\theta_i)$ can be given by substituting (8) into (9). Because $\theta \in \Delta_d$, we have for all $i \in [d]$, $0 \leq \theta_i \leq 1$. In the following Lemma 2, it is observed that the derivate $J'_i(\theta_i)$ can be a constant or have a piecewise structure, which is the key to deduce the SOTOPO algorithm.

²The infima can be replaced by minimization if the convention “ $0/0 = 0$ ” is used.

Lemma 2. Assume that for all $i \in [d]$, $J'_i(0)$ and $J'_i(1)$ have been computed. Denote $r_{i1} \stackrel{\text{def}}{=} \frac{|x_i|}{\sqrt{-2\eta J'_i(0)}}$ and $r_{i2} \stackrel{\text{def}}{=} \frac{|x_i|}{\sqrt{-2\eta J'_i(1)}}$, then $J'_i(\theta_i)$ belongs to one of the 4 cases,

$$\begin{aligned} (\text{case a}) : J'_i(\theta_i) &= 0, \quad 0 \leq \theta_i \leq 1, & (\text{case b}) : J'_i(\theta_i) &= J'_i(0) < 0, \quad 0 \leq \theta_i \leq 1, \\ (\text{case c}) : J'_i(\theta_i) &= \begin{cases} J'_i(0), & 0 \leq \theta_i \leq r_{i1} \\ -\frac{x_i^2}{2\eta\theta_i^2}, & r_{i1} < \theta_i \leq 1 \end{cases}, & (\text{case d}) : J'_i(\theta_i) &= \begin{cases} J'_i(0), & 0 \leq \theta_i \leq r_{i1} \\ -\frac{x_i^2}{2\eta\theta_i^2}, & r_{i1} < \theta_i < r_{i2} \\ J'_i(1), & r_{i2} \leq \theta_i \leq 1 \end{cases}. \end{aligned}$$

Although the formulation of $J'_i(\theta_i)$ is complicated, by summarizing the property of the 4 cases in Lemma 2, we have Corollary 1.

Corollary 1. For all $i \in [d]$ and $0 \leq \theta_i \leq 1$, if the derivate $J'_i(\theta_i)$ is not always 0, then $J'_i(\theta_i)$ is a non-decreasing, continuous function with value always less than 0.

Corollary 1 shows that except the trivial (case a), for all $i \in [d]$, whichever $J'_i(\theta_i)$ belong to (case b), (case c) or case (d), they all share the same group of properties, which makes a consistent iterative procedure possible for all the cases. The different formulations in the four cases mainly have impact about the stopping criteria of SOTOPO.

2.2 The property of the optimal solution

The Lagrangian of the problem (7) is

$$\mathcal{L}(\theta, \gamma, \zeta) \stackrel{\text{def}}{=} J(\theta) + \gamma \left(\sum_{i=1}^d \theta_i - 1 \right) - \langle \zeta, \theta \rangle, \quad (10)$$

where $\gamma \in \mathbb{R}$ is a Lagrange multiplier and $\zeta \in \mathbb{R}_+^d$ is a vector of non-negative Lagrange multipliers. Due to the coordinate separable property of $J(\theta)$ in (9), it follows that $\frac{\partial J(\theta)}{\partial \theta_i} = J'_i(\theta_i)$. Then the KKT condition of (10) can be written as

$$\forall i \in [d], \quad J'_i(\theta_i) + \gamma - \zeta_i = 0, \quad \zeta_i \theta_i = 0, \quad \text{and} \quad \sum_{i=1}^d \theta_i = 1. \quad (11)$$

By reformulating the KKT condition (11), we have Lemma 3.

Lemma 3. If $(\tilde{\gamma}, \tilde{\theta}, \tilde{\zeta})$ is a stationary point of (10), then $\tilde{\theta}$ is an optimal solution of (7). Meanwhile, denote $S \stackrel{\text{def}}{=} \{i : \tilde{\theta}_i > 0\}$ and $T \stackrel{\text{def}}{=} \{j : \tilde{\theta}_j = 0\}$, then the KKT condition can be formulated as

$$\begin{cases} \sum_{i \in S} \tilde{\theta}_i = 1; \\ \text{for all } j \in T, \quad \tilde{\theta}_j = 0; \\ \text{for all } i \in S, \quad \tilde{\gamma} = -J'_i(\tilde{\theta}_i) \geq \max_{j \in T} -J'_j(0). \end{cases} \quad (12)$$

By Lemma 3, if the set S in Lemma 3 is known beforehand, then we can compute $\tilde{\theta}$ by simply applying the equations in (12). Therefore finding the optimal solution $\tilde{\theta}$ is equivalent to finding the set of the nonzero elements of $\tilde{\theta}$.

2.3 The soft thresholding projection algorithm

In Lemma 3, for each $i \in [d]$ with $\tilde{\theta}_i > 0$, it is shown that the negative derivate $-J'_i(\tilde{\theta}_i)$ is equal to a single variable $\tilde{\gamma}$. Therefore, a much simpler problem can be obtained if we know the coordinates of these positive elements. At first glance, it seems difficult to identify these coordinates, because the number of potential subsets of coordinates is clearly exponential on the dimension d . However, the property clarified by Lemma 2 enables an efficient procedure for identifying the nonzero elements of $\tilde{\theta}$. Lemma 4 is a key tool in deriving the procedure for identifying the non-zero elements of $\tilde{\theta}$.

Lemma 4 (Nonzero element identification). Let $\tilde{\theta}$ be an optimal solution of (7). Let s and t be two coordinates such that $J'_s(0) < J'_t(0)$. If $\tilde{\theta}_s = 0$, then $\tilde{\theta}_t$ must be 0 as well; equivalently, if $\tilde{\theta}_t > 0$, then $\tilde{\theta}_s$ must be greater than 0 as well.

Lemma 4 shows that if we sort $u \stackrel{\text{def}}{=} -\nabla J(0)$ such that $u_{i_1} \geq u_{i_2} \geq \dots \geq u_{i_d}$, where $\{i_1, i_2, \dots, i_d\}$ is a permutation of $[d]$, then the set S in Lemma 3 is of the form $\{i_1, i_2, \dots, i_\varrho\}$, where $1 \leq \varrho \leq d$. If ϱ is obtained, then we can use the fact that for all $j \in [\varrho]$,

$$-J'_{i_j}(\tilde{\theta}_{i_j}) = \tilde{\gamma} \quad \text{and} \quad \sum_{j=1}^{\varrho} \tilde{\theta}_{i_j} = 1 \quad (13)$$

to compute $\tilde{\gamma}$. Therefore, by Lemma 4, we can efficiently identify the nonzero elements of the optimal solution $\tilde{\theta}$ after a sort operation, which costs $O(d \log d)$. However based on Lemmas 2 and 3, the sort cost $O(d \log d)$ can be further reduced by the following Lemma 5.

Lemma 5 (Efficient identification). *Assume $\tilde{\theta}$ and S are given in Lemma 3. Then for all $i \in S$,*

$$-J'_i(0) \geq \max_{j \in [d]} \{-J'_j(1)\}. \quad (14)$$

By Lemma 5, before ordering u , we can filter out all the coordinates i 's that satisfy $-J'_i(0) < \max_{j \in [d]} -J'_j(1)$. Based on Lemmas 4 and 5, we propose the SOft ThreshOlding PRojection (SOTOPO) algorithm in Alg. 1 to efficiently obtain an optimal solution $\tilde{\theta}$. In the step 1, by Lemma 5, we find the quantity v_m, i_m and Q . In the step 2, by Lemma 4, we sort the elements $\{-J'_i(0) \mid i \in Q\}$. In the step 3, because S in Lemma 3 is of the form $\{i_1, i_2, \dots, i_\varrho\}$, we search the quantity ρ from 1 to $|Q| + 1$ until a stopping criteria is met. In Alg. 1, the number of nonzero elements of $\tilde{\theta}$ is ρ or $\rho - 1$. In the step 4, we compute the $\tilde{\gamma}$ in Lemma 3 according to the conditions. In the step 5, the optimal $\tilde{\theta}$ and the corresponding \tilde{h}, \tilde{x} are given.

Algorithm 1 $\tilde{x} = \text{SOTOPO}(\nabla f(x), x, \lambda, \eta)$

1. Find

$$(v_m, i_m) \stackrel{\text{def}}{=} (\max_{i \in [d]} \{-J'_i(1)\}, \arg \max_{i \in [d]} \{-J'_i(1)\}), Q \stackrel{\text{def}}{=} \{i \in [d] \mid -J'_i(0) > v_m\}.$$

2. Sort $\{-J'_i(0) \mid i \in Q\}$ such that $-J'_{i_1}(0) \geq -J'_{i_2}(0) \geq \dots \geq -J'_{i_{|Q|}}(0)$, where $\{i_1, i_2, \dots, i_{|Q|}\}$ is a permutation of the elements in Q . Denote

$$v \stackrel{\text{def}}{=} (-J'_{i_1}(0), -J'_{i_2}(0), \dots, -J'_{i_{|Q|}}(0), v_m), \quad \text{and} \quad i_{|Q|+1} \stackrel{\text{def}}{=} i_m, v_{|Q|+1} \stackrel{\text{def}}{=} v_m.$$

3. For $j \in [|Q| + 1]$, denote $R_j = \{i_k \mid k \in [j]\}$. Search from 1 to $|Q| + 1$ to find the quantity

$$\rho \stackrel{\text{def}}{=} \min \{j \in [|Q| + 1] \mid J'_{i_j}(0) = J'_{i_j}(1) \text{ or } \sum_{l \in R_j} |x_l| \geq \sqrt{2\eta v_j} \text{ or } j = |Q| + 1\}.$$

4. The $\tilde{\gamma}$ in Lemma 3 is given by

$$\tilde{\gamma} = \begin{cases} \left(\sum_{l \in R_{\rho-1}} |x_l| \right)^2 / (2\eta), & \text{if } \sum_{l \in R_{\rho-1}} |x_l| \geq \sqrt{2\eta v_\rho}; \\ v_\rho, & \text{otherwise.} \end{cases}$$

5. Then the $\tilde{\theta}$ in Lemma 3 and its corresponding \tilde{h}, \tilde{x} in (3) and (4) are obtained by

$$(\tilde{\theta}_l, \tilde{h}_l, \tilde{x}_l) = \begin{cases} \left(\frac{|x_l|}{\sqrt{2\eta\tilde{\gamma}}}, -x_l, 0 \right), & \text{if } l \in R_\rho \setminus \{i_\rho\}; \\ \left(1 - \sum_{k \in R_\rho \setminus \{i_\rho\}} \tilde{\theta}_k, \tilde{g}_l(\tilde{\theta}_l), x_l + \tilde{g}_l(\tilde{\theta}_l) \right), & \text{if } l = i_\rho; \\ (0, 0, x_l), & \text{if } l \in [d] \setminus R_\rho. \end{cases}$$

In Theorem 1, we give the main result about the SOTOPO algorithm.

Theorem 1. *The SOTOPO algorithm in Alg. 1 can get the exact minimizer \tilde{h}, \tilde{x} of the ℓ_1 -regularized ℓ_1 -norm square approximation problem in (3) and (4).*

The SOTOPO algorithm seems complicated but is indeed efficient. The dominant operations in Alg. 1 are steps 1 and 2 with the total cost $O(d + |Q| \log |Q|)$. To show the effect of the complexity reduction by Lemma 5, we give the following fact.

Proposition 1. For the optimization problem defined in (5)-(7), where λ is the regularization parameter of the original problem (1), we have that

$$0 \leq \max_{i \in [d]} \left\{ \sqrt{\frac{-2J'_i(0)}{\eta}} \right\} - \max_{j \in [d]} \left\{ \sqrt{\frac{-2J'_j(1)}{\eta}} \right\} \leq 2\lambda. \quad (15)$$

Assume v_m is defined in the step 1 of Alg. 1. By Proposition 1, for all $i \in Q$,

$$\sqrt{\frac{-2J'_i(0)}{\eta}} \leq \max_{k \in [d]} \left\{ \sqrt{\frac{-2J'_k(0)}{\eta}} \right\} \leq \max_{j \in [d]} \left\{ \sqrt{\frac{-2J'_j(1)}{\eta}} \right\} + 2\lambda = \sqrt{\frac{2v_m}{\eta}} + 2\lambda,$$

Therefore at least the coordinates j 's that satisfy $\sqrt{\frac{-2J'_j(0)}{\eta}} > \sqrt{\frac{2v_m}{\eta}} + 2\lambda$ will be not contained in Q . In practice, it can considerably reduce the sort complexity.

Remark 1. SOTOPO can be viewed as an extension of the SOPOPO algorithm [20] by changing the objective function from Euclidean distance to a more general function $J(\theta)$ in (9). It should be noted that Lemma 5 does not have a counterpart in the case that the objective function is Euclidean distance [20]. In addition, some extension of randomized median finding algorithm [12] with linear time in our setting is also deserved to research. Due to the limited space, it is left for further discussion.

3 The ASGCD algorithm

Now we can come back to our motivation, *i.e.*, accelerating GCD to obtain the optimal convergence rate $O(1/\sqrt{\epsilon})$ by Nesterov's acceleration and reducing the complexity of greedy selection by stochastic optimization. The main idea is that although like any (block) coordinate descent algorithm, the proposed new rule, *i.e.*, minimizing the problem in (3), performs update on one or several coordinates, it is a generalized proximal gradient descent problem based on ℓ_1 -norm. Therefore this rule can be applied into the existing Nesterov's acceleration and stochastic optimization framework "Katyusha" [1] if it can be solved efficiently. The final result is the accelerated stochastic greedy coordinate descent (ASGCD) algorithm, which is described in Alg. 2.

Algorithm 2 ASGCD

$$\delta = \log(d) - 1 - \sqrt{(\log(d) - 1)^2 - 1};$$

$$p = 1 + \delta, q = \frac{p}{p-1}, C = \frac{d^{2\delta}}{\delta};$$

$$z_0 = y_0 = \tilde{x}_0 = \vartheta_0 = 0;$$

$$\tau_2 = \frac{1}{2}, m = \lceil \frac{n}{b} \rceil, \eta = \frac{1}{(1+2\frac{n-b}{b(n-1)})L_1};$$

for $s = 0, 1, 2, \dots, S - 1$, **do**

1. $\tau_{1,s} = \frac{2}{s+4}, \alpha_s = \frac{\eta}{\tau_{1,s}C};$

2. $\mu_s = \nabla f(\tilde{x}_s);$

3. **for** $l = 0, 1, \dots, m - 1$, **do**

- (a) $k = (sm) + l;$

- (b) randomly sample a mini batch \mathcal{B} of size b from $\{1, 2, \dots, n\}$ with equal probability;

- (c) $x_{k+1} = \tau_{1,s}z_k + \tau_2\tilde{x}_s + (1 - \tau_{1,s} - \tau_2)y_k;$

- (d) $\tilde{\nabla}_{k+1} = \mu_s + \frac{1}{b} \sum_{j \in \mathcal{B}} (\nabla f_j(x_{k+1}) - \nabla f_j(\tilde{x}_s));$

- (e) $y_{k+1} = \text{SOTOPO}(\tilde{\nabla}_{k+1}, x_{k+1}, \lambda, \eta);$

- (f) $(z_{k+1}, \vartheta_{k+1}) = \text{pCOMID}(\tilde{\nabla}_{k+1}, \vartheta_k, q, \lambda, \alpha_s);$

- end for**

4. $\tilde{x}_{s+1} = \frac{1}{m} \sum_{l=1}^m y_{sm+l};$

end for

Output: \tilde{x}_S

Algorithm 3 $(\tilde{x}, \tilde{\vartheta}) = \text{pCOMID}(g, \vartheta, q, \lambda, \alpha)$

1. $\forall i \in [d], \tilde{\vartheta}_i = \text{sign}(\vartheta_i - \alpha g_i) \cdot \max\{0, |\vartheta_i - \alpha g_i| - \alpha \lambda\};$
 2. $\forall i \in [d], \tilde{x}_i = \frac{\text{sign}(\tilde{\vartheta}_i) |\tilde{\vartheta}_i|^{q-1}}{\|\tilde{\vartheta}\|_q^{q-2}};$
 3. **Output:** $\tilde{x}, \tilde{\vartheta}.$
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In Alg. 2, the gradient descent step 3(e) is solved by the proposed SOTOPO algorithm, while the mirror descent step 3(f) is solved by the COMID algorithm with p -norm divergence [13, Sec. 7.2]. We denote the mirror descent step as pCOMID in Alg. 3. All other parts are standard steps in the Katyusha framework except some parameter settings. For example, instead of the custom setting $p = 1 + 1/\log(d)$ [21, 13], a particular choice $p = 1 + \delta$ (δ is defined in Alg. 2) is used to minimize the $C = \frac{d^{2\delta}}{\delta}$. C varies slowly over d and is upper bounded by $\log^2(d)$. Meanwhile, α_{k+1} depends on the extra constant C . Furthermore, the step size $\eta = \frac{1}{(1+2\frac{n-b}{b(n-1)})L_1}$ is used, where L_1 is defined in (2). Finally, unlike [1, Alg. 2], we let the batch size b as an algorithm parameter to cover both the stochastic case $b < n$ and the deterministic case $b = n$. To the best of our knowledge, the existing GCD algorithms are deterministic, therefore by setting $b = n$, we can compare with the existing GCD algorithms better.

Based on the efficient SOTOPO algorithm, ASGCD has nearly the same iteration complexity with the standard form [1, Alg. 2] of Katyusha. Meanwhile we have the following convergence rate.

Theorem 2. *If each $f_j(x)$ ($j \in [n]$) is convex, L_1 -smooth in (2) and x^* is an optimum of the ℓ_1 -regularized problem (1), then ASGCD satisfies*

$$\mathbb{E}[F(\tilde{x}^S)] - F(x^*) \leq \frac{4}{(S+3)^2} \left(1 + \frac{1+2\beta(b)}{2m} C\right) L_1 \|x^*\|_1^2 = O\left(\frac{CL_1 \|x^*\|_1^2}{S^2}\right), \quad (16)$$

where $\beta(b) = \frac{n-b}{b(n-1)}$, S , b , m and C are given in Alg. 2. In other words, ASGCD achieves an ϵ -additive error (i.e., $\mathbb{E}[F(\tilde{x}^S)] - F(x^*) \leq \epsilon$) using at most $O\left(\frac{\sqrt{CL_1} \|x^*\|_1}{\sqrt{\epsilon}}\right)$ iterations.

In Table 1, we give the convergence rate of the existing algorithms and ASGCD to solve the ℓ_1 -regularized problem (1). In the first column, ‘‘Acc’’ and ‘‘Non-Acc’’ denote the corresponding algorithms are Nesterov’s accelerated or not respectively, ‘‘Primal’’ and ‘‘Dual’’ denote the corresponding algorithms solves the primal problem (1) and its regularized dual problem [22] respectively. ℓ_2 -norm and ℓ_1 -norm denote the theoretical guarantee is based on ℓ_2 -norm and ℓ_1 -norm respectively. In terms of ℓ_2 -norm based guarantee, Katyusha and APPROX give the state of the art convergence rate $O\left(\frac{\sqrt{L_2} \|x^*\|_2}{\sqrt{\epsilon}}\right)$. In terms of ℓ_1 -norm based guarantee, GCD gives the state of the art convergence rate $O\left(\frac{L_1 \|x^*\|_1^2}{\epsilon}\right)$, which is only applicable for the smooth case $\lambda = 0$ in (1). When $\lambda > 0$, the generalized GS- r , GS- s and GS- q rules generally have worse theoretical guarantee than GCD [19]. While the bound of ASGCD in this paper is $O\left(\frac{\sqrt{L_1} \|x\|_1 \log d}{\sqrt{\epsilon}}\right)$, which can be viewed as an accelerated version of the ℓ_1 -norm based guarantee $O\left(\frac{L_1 \|x\|_1^2}{\epsilon}\right)$. Meanwhile, because the bound depends on $\|x^*\|_1$ rather than $\|x^*\|_2$ and on L_1 rather than L_2 (L_1 and L_2 are defined in (2)), for the ℓ_1 -ERM problem, if the samples are high-dimensional, dense and the regularization parameter λ is relatively large, then it is possible that $L_1 \ll L_2$ (in the extreme case, $L_2 = dL_1$ [11]) and $\|x^*\|_1 \approx \|x^*\|_2$. In this case, the ℓ_1 -norm based guarantee $O\left(\frac{\sqrt{L_1} \|x\|_1 \log d}{\sqrt{\epsilon}}\right)$ of ASGCD is better than the ℓ_2 -norm based guarantee $O\left(\frac{\sqrt{L_2} \|x^*\|_2}{\sqrt{\epsilon}}\right)$ of Katyusha and APPROX. Finally, whether the $\log d$ factor in the bound of ASGCD (which also appears in the COMID [13] analysis) is necessary deserves further research.

Remark 2. *When the batch size $b = n$, ASGCD is a deterministic algorithm. In this case, we can use a better smooth constant T_1 that satisfies $\|\nabla f(x) - \nabla f(y)\|_\infty \leq T_1 \|x - y\|_1$ rather than L_1 [1].*

Remark 3. *The necessity of computing the full gradient beforehand is the main bottleneck of GCD in applications [19]. There exists some work [11] to avoid the computation of full gradient by performing some approximate greedy selection. While the method in [11] needs preprocessing,*

Table 1: Convergence rate on ℓ_1 -regularized empirical risk minimization problems. (For GCD, the convergence rate is applied for $\lambda = 0$.)

ALGORITHM TYPE	PAPER	CONVERGENCE RATE
NON-ACC, PRIMAL, ℓ_2 -NORM	SAGA [10]	$O\left(\frac{L_2 \ x^*\ _2^2}{\epsilon}\right)$
ACC, PRIMAL, ℓ_2 -NORM	KATYUSHA [1]	$O\left(\frac{\sqrt{L_2} \ x^*\ _2}{\sqrt{\epsilon}}\right)$
ACC, DUAL, ℓ_2 -NORM	ACC-SDCA [23] SPDC [26] APCG [16] APPROX [14]	$O\left(\frac{\sqrt{L_2} \ x^*\ _2}{\sqrt{\epsilon}} \log\left(\frac{1}{\epsilon}\right)\right)$
NON-ACC, PRIMAL, ℓ_1 -NORM	GCD [3]	$O\left(\frac{L_1 \ x^*\ _1^2}{\epsilon}\right)$
ACC, PRIMAL, ℓ_1 -NORM	ASGCD (THIS PAPER)	$O\left(\frac{\sqrt{L_1} \ x^*\ _1 \log d}{\sqrt{\epsilon}}\right)$

incoherence condition for dataset and is somewhat complicated. Contrary to [11], the proposed ASGCD algorithm reduces the complexity of greedy selection by a factor up to n in terms of the amortized cost by simply applying the existing stochastic variance reduction framework.

4 Experiments

In this section, we use numerical experiments to demonstrate the theoretical results in Section 3 and show the empirical performance of ASGCD with batch size $b = 1$ and its deterministic version with $b = n$ (In Fig. 1 they are denoted as ASGCD ($b = 1$) and ASGCD ($b = n$) respectively). In addition, following the claim to using data access rather than CPU time [21] and the recent SGD and RCD literature [15, 16, 1], we use the data access, i.e., the number of times the algorithm accesses the data matrix, to measure the algorithm performance. To show the effect of Nesterov’s acceleration, we compare ASGCD ($b = n$) with the non-accelerated greedy coordinate descent with GS- q rule, i.e., coordinate gradient descent (CGD) [24]. To show the effect of both Nesterov’s acceleration and stochastic optimization strategies, we compare ASGCD ($b = 1$) with Katyusha [1, Alg. 2]. To show the effect of the proposed new rule in Section 2, which is based on ℓ_1 -norm square approximation, we compare ASGCD ($b = n$) with the ℓ_2 -norm based proximal accelerated full gradient (AFG) implemented by the linear coupling framework [4]. Meanwhile, as a benchmark of stochastic optimization for the problems with finite-sum structure, we also show the performance of proximal stochastic variance reduced gradient (SVRG) [25]. In addition, based on [1] and our experiments, we find that “Katyusha” [1, Alg. 2] has the best empirical performance in general for the ℓ_1 -regularized problem (1). Therefore other well-known state-of-art algorithms, such as APCG [16] and accelerated SDCA [23], are not included in the experiments.

The datasets are obtained from LIBSVM data [9] and summarized in Table 2. All the algorithms are used to solve the following lasso problem

$$\min_{x \in \mathbb{R}^d} \{f(x) + \lambda \|x\|_1 = \frac{1}{2n} \|b - Ax\|_2^2 + \lambda \|x\|_1\} \quad (17)$$

on the 3 datasets, where $A = (a_1, a_2, \dots, a_n)^T = (h_1, h_2, \dots, h_d) \in \mathbb{R}^{n \times d}$ with each $a_j \in \mathbb{R}^d$ representing a sample vector and $h_i \in \mathbb{R}^n$ representing a feature vector, $b \in \mathbb{R}^n$ is the prediction vector.

Table 2: Characteristics of three real datasets.

DATASET NAME	# SAMPLES n	# FEATURES d
LEUKEMIA	38	7129
GISETTE	6000	5000
MNIST	60000	780

For ASGCD ($b = 1$) and Katyusha [1, Alg. 2], we can use the tight smooth constant $L_1 = \max_{j \in [n], i \in [d]} |a_{j,i}^2|$ and $L_2 = \max_{j \in [n]} \|a_j\|_2^2$ respectively in their implementation. While for AS-

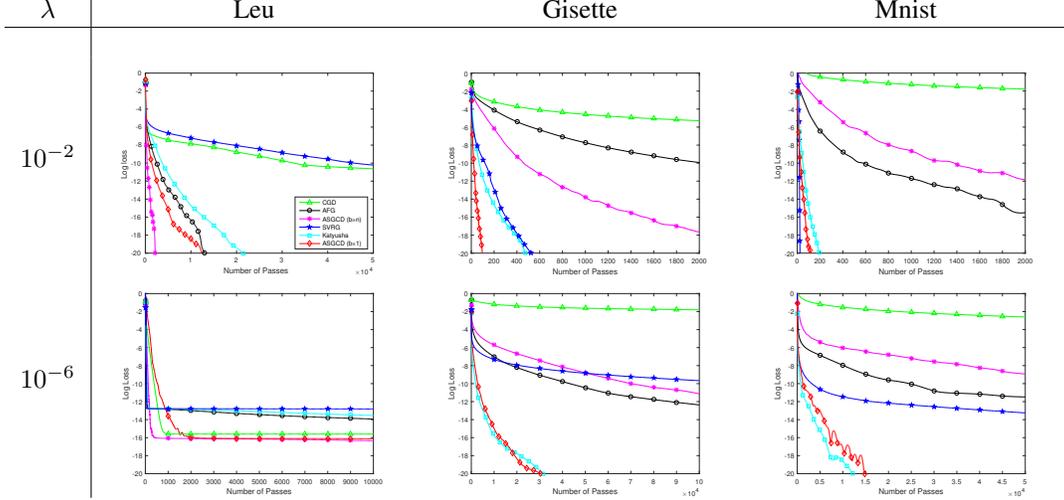


Figure 1: Comparing AGCD ($b = 1$) and ASGCD ($b = n$) with CGD, SVRG, AFG and Katyusha on Lasso.

GCD ($b = n$) and AFG, the better smooth constant $T_1 = \frac{\max_{i \in [d]} \|h_i\|_2^2}{n}$ and $T_2 = \frac{\|A\|^2}{n}$ are used respectively. The learning rate of CGD and SVRG are tuned in $\{10^{-6}, 10^{-5}, 10^{-4}, 10^{-3}, 10^{-2}, 10^{-1}\}$.

Table 3: Factor rates of for the 6 cases

λ	LEU	GISETTE	MNIST
10^{-2}	(0.85, 1.33)	(0.88, 0.74)	(5.85, 3.02)
10^{-6}	(1.45, 2.27)	(3.51, 2.94)	(5.84, 3.02)

We use $\lambda = 10^{-6}$ and $\lambda = 10^{-2}$ in the experiments. In addition, for each case (Dataset, λ), AFG is used to find an optimum x^* with enough accuracy.

The performance of the 6 algorithms is plotted in Fig. 1. We use Log loss $\log(F(x_k) - F(x^*))$ in the y -axis. x -axis denotes the number that the algorithm access the data matrix A . For example, ASGCD ($b = n$) accesses A once in each iteration, while ASGCD ($b = 1$) accesses A twice in an entire outer iteration. For each case (Dataset, λ), we compute the rate $(r_1, r_2) = \left(\frac{\sqrt{CL_1}\|x^*\|_1}{\sqrt{L_2}\|x^*\|_2}, \frac{\sqrt{CT_1}\|x^*\|_1}{\sqrt{T_2}\|x^*\|_2} \right)$ in Table 3. First, because of the acceleration effect, ASGCD ($b = n$) are always better than the non-accelerated CGD algorithm; second, by comparing ASGCD($b = 1$) with Katyusha and ASGCD ($b = n$) with AFG, we find that for the cases (Leu, 10^{-2}), (Leu, 10^{-6}) and (Gisette, 10^{-2}), ASGCD ($b = 1$) dominates Katyusha [1, Alg.2] and ASGCD ($b = n$) dominates AFG. While the theoretical analysis in Section 3 shows that if r_1 is relatively small such as around 1, then ASGCD ($b = 1$) will be better than [1, Alg.2]. For the other 3 cases, [1, Alg.2] and AFG are better. The consistency between Table 3 and Fig. 1 demonstrates the theoretical analysis.

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A Proof of Lemma 1

Proof. By using the variational identity $\|g\|_1^2 = \inf_{\theta \in \Delta_d} \sum_{i=1}^d \frac{g_i^2}{\theta_i}$ in [5] and the definition of $J(g, \theta)$ in (3), it follows that (3) can be rewritten as

$$\tilde{h} = \arg \min_{g \in \mathbb{R}^d} \left\{ \inf_{\theta \in \Delta_d} J(g, \theta) \right\}.$$

By the joint convexity of $J(g, \theta)$, we can find the minimizer \tilde{h} by swapping the optimization order of g and θ , which is to say based on the definition of $\tilde{g}(\theta)$, $J(\theta)$ and $\tilde{\theta}$, we have

$$\tilde{h} = \tilde{g}(\tilde{\theta}).$$

Therefore, the minimization problem to find \tilde{h} in (3) can be equivalently transformed to the problem (7). Meanwhile, it is observed that $J(g, \theta)$ is coordinate separable, *i.e.*,

$$J(g, \theta) = \sum_{i=1}^d J_i(g_i, \theta_i), \text{ where } J_i(g_i, \theta_i) \stackrel{\text{def}}{=} \nabla_i f(x) g_i + \frac{1}{2\eta} \frac{g_i^2}{\theta_i} + \lambda |x_i + g_i|. \quad (\text{A.1})$$

By the definition of $\tilde{g}(\theta)$ in (6), $\tilde{g}(\theta)$ is also coordinate separable, *i.e.* for all $i \in [d]$,

$$\tilde{g}_i(\theta) = \tilde{g}_i(\theta_i) \stackrel{\text{def}}{=} \arg \min_{g_i \in \mathbb{R}} \left\{ \nabla_i f(x) g_i + \frac{1}{2\eta} \frac{g_i^2}{\theta_i} + \lambda |x_i + g_i| \right\}.$$

By using the iterative soft thresholding (IST) operator [7], for all $i \in [d]$,

$$\tilde{g}_i(\theta_i) = \text{sign}(x_i - \theta_i \eta \nabla_i f(x)) \cdot \max\{0, |x_i - \theta_i \eta \nabla_i f(x)| - \theta_i \eta \lambda\} - x_i. \quad (\text{A.2})$$

Then it implies that $J(\theta)$ is also coordinate separable, *i.e.*,

$$J(\theta) = \sum_{i=1}^d J_i(\theta_i), \text{ where } J_i(\theta_i) \stackrel{\text{def}}{=} J_i(\tilde{g}_i(\theta_i), \theta_i). \quad (\text{A.3})$$

□

B Proofs of Lemma 2, Corollary 1 and Proposition 1

Proof of Lemma 2. For all $i \in [d]$, due to $\theta \in \Delta_d$, we have $0 \leq \theta_i \leq 1$. By substituting (8) into (9), we get the expression of $J_i(\theta_i)$. Taking the derivate of $J_i(\theta_i)$ and setting $\theta_i = 0, 1$ respectively, then we get the expressions of $J'_i(0), J'_i(1)$ as follows.

For all $i \in [d]$ and $\theta_i \geq 0$, denote

$$\nu_i \stackrel{\text{def}}{=} -\frac{(\max\{|\nabla_i f(x)| - \lambda, 0\})^2 \eta}{2}, \quad \chi_i(\theta_i) \stackrel{\text{def}}{=} -\frac{(\text{sign}(x_i - \theta_i \eta \nabla_i f(x)) \lambda + \nabla_i f(x))^2 \eta}{2}, \quad (\text{B.1})$$

then the derivate $J'_i(\theta_i)$ at $\theta_i = 0, 1$ are

$$J'_i(0) = \begin{cases} \nu_i, & x_i = 0, \\ \chi_i(0), & x_i \neq 0 \end{cases}, \quad J'_i(1) = \begin{cases} -\frac{x_i^2}{2\eta}, & |x_i - \eta \nabla_i f(x)| \leq \eta \lambda \\ \chi_i(1), & |x_i - \eta \nabla_i f(x)| > \eta \lambda \end{cases}. \quad (\text{B.2})$$

For all $i \in [d]$, according to the values of x_i and $\nabla_i f(x)$, by classified discussion, we can show that $J'_i(\theta_i)$ belongs to one of the 4 cases in Lemma 2. Assume that r_{i1} and r_{i2} have been defined in

Lemma 2. Firstly, we denote

$$O \stackrel{\text{def}}{=} \{i | 0 \in \nabla_i f(x) + \lambda \partial |x_i|\}, \quad (\text{B.3})$$

$$\begin{aligned} U \stackrel{\text{def}}{=} & \{i \in [d] | x_i \geq 0, \nabla_i f(x) < -\lambda\} \\ & \cup \{i \in [d] | x_i \leq 0, \nabla_i f(x) > \lambda\} \\ & \cup \{i \in [d] | x_i > 0, \nabla_i f(x) > -\lambda, r_{i1} \geq 1\} \\ & \cup \{i \in [d] | x_i < 0, \nabla_i f(x) < \lambda, r_{i1} \geq 1\}, \end{aligned} \quad (\text{B.4})$$

$$\begin{aligned} V \stackrel{\text{def}}{=} & \{i \in [d] | x_i > 0, -\lambda < \nabla_i f(x) \leq \lambda, r_{i1} < 1\} \\ & \cup \{i \in [d] | x_i < 0, -\lambda \leq \nabla_i f(x) < \lambda, r_{i1} < 1\} \\ & \cup \{i \in [d] | x_i > 0, \nabla_i f(x) > \lambda, r_{i2} \geq 1\} \\ & \cup \{i \in [d] | x_i < 0, \nabla_i f(x) < -\lambda, r_{i2} \geq 1\}, \end{aligned} \quad (\text{B.5})$$

$$\begin{aligned} W \stackrel{\text{def}}{=} & \{i \in [d] | x_i > 0, \nabla_i f(x) > \lambda, r_{i2} < 1\} \\ & \cup \{i \in [d] | x_i < 0, \nabla_i f(x) < -\lambda, r_{i2} < 1\}. \end{aligned} \quad (\text{B.6})$$

Then based on the expressions of $J_i(\theta_i)$ in (A.1), (A.1) and (A.3), we can summarize the results as follows

- If $i \in O$, then $J'_i(\theta_i)$ belongs to the (case a) in Lemma 2.
- If $i \in U$, then $J'_i(\theta_i)$ belongs to the (case b) in Lemma 2.
- If $i \in V$, then $J'_i(\theta_i)$ belongs to the (case c) in Lemma 2.
- If $i \in W$, then $J'_i(\theta_i)$ belongs to the (case d) in Lemma 2.

□

Proof of Corollary 1. Corollary 1 can be obtained by simply summarizing the 4 cases in Lemma 2. □

Proof of Proposition 1. Assume that $\chi_i(\theta_i)$ is defined in (B.1), O, U, V and W are defined in (B.3)-(B.6).

Firstly, by checking $i \in O, U, V$ or W orderly and using the expression of $J'_i(0)$ and $J'_i(1)$, it follows that

- For $i \in O \cup U$, $J'_i(0) = J'_i(1)$. Therefore $0 \leq \sqrt{\frac{-2J'_i(0)}{\eta}} - \sqrt{\frac{-2J'_i(1)}{\eta}} \leq 2\lambda$.
- For $i \in V$, by the definition of V and Lemma 2, it follows that $J'_i(0) = \chi_i(0)$ and $J'_i(1) = -\frac{x_i^2}{2\eta}$.

– If $-\lambda \leq \nabla_i f(x) \leq \lambda$, then

$$\sqrt{\frac{-2J'_i(0)}{\eta}} - \sqrt{\frac{-2J'_i(1)}{\eta}} \leq \sqrt{\frac{-2J'_i(0)}{\eta}} = |\text{sign}(x_i)\lambda + \nabla_i f(x)| \leq \lambda + |\nabla_i f(x)| \leq 2\lambda.$$

- If $\nabla_i f(x) > \lambda$, then analyzing the expressions of $J_i(\theta_i)$ in this case, there exists $\hat{\theta}_i > 1$ such that $J'_i(\hat{\theta}_i) = \chi_i(\hat{\theta}_i)$. By the non-decreasing property of $J'_i(\theta_i)$ in Corollary 1 (which can be extended to $\theta_i > 1$ trivially), $J'_i(1) \leq J'_i(\hat{\theta}_i)$. Then

$$\begin{aligned}
& \sqrt{\frac{-2J'_i(0)}{\eta}} - \sqrt{\frac{-2J'_i(1)}{\eta}} \leq \sqrt{\frac{-2J'_i(0)}{\eta}} - \sqrt{\frac{-2J'_i(\hat{\theta}_i)}{\eta}} \\
& = |\text{sign}(x_i)\lambda + \nabla_i f(x)| - |\text{sign}(x_i - \hat{\theta}_i \eta \nabla_i f(x))\lambda + \nabla_i f(x)| \\
& \stackrel{\textcircled{1}}{=} (\text{sign}(x_i)\lambda + \nabla_i f(x)) - (\text{sign}(x_i - \hat{\theta}_i \eta \nabla_i f(x))\lambda + \nabla_i f(x)) \\
& = (\text{sign}(x_i) - \text{sign}(x_i - \hat{\theta}_i \eta \nabla_i f(x)))\lambda \\
& \leq 2\lambda,
\end{aligned}$$

where $\textcircled{1}$ is by the fact that $\nabla_i f(x) \geq \lambda$.

– If $\nabla_i f(x) < -\lambda$, we can give a similar analysis as the case $\nabla_i f(x) > \lambda$.

- If $i \in W$, by the definition of W and Lemma 2, it follows that $J'_i(0) = \chi_i(0)$ and $J'_i(1) = \chi_i(1)$. Because if $i \in W$, then $|\nabla_i f(x)| \geq \lambda$, we can give a similar analysis as the case in $i \in V$.

By the above analysis, it follows that

$$\max_{i \in [d]} \left\{ \sqrt{\frac{-2J'_i(0)}{\eta}} \right\} - \max_{j \in [d]} \left\{ \sqrt{\frac{-2J'_j(1)}{\eta}} \right\} \leq \max_{i \in [d]} \left\{ \sqrt{\frac{-2J'_i(0)}{\eta}} - \sqrt{\frac{-2J'_i(1)}{\eta}} \right\} \leq 2\lambda.$$

In addition, by Corollary 1,

$$0 \leq \max_{i \in [d]} \left\{ \sqrt{\frac{-2J'_i(0)}{\eta}} \right\} - \max_{j \in [d]} \left\{ \sqrt{\frac{-2J'_j(1)}{\eta}} \right\}.$$

Proposition 1 is proved. \square

C Proof of Lemma 3

The Lagrangian of the problem (7) is (10). By the property of KKT condition, if $(\tilde{\gamma}, \tilde{\theta}, \tilde{\zeta})$ is a stationary point of the problem (10), then $\tilde{\theta}$ is an optimal solution of (7). Based on the value of $\tilde{\theta}$, one can divide $[d]$ into two disjoint parts S and T , where

$$S = \{i : \tilde{\theta}_i > 0\} \text{ and } T = \{j : \tilde{\theta}_j = 0\}.$$

Then $\forall i \in S$, by the complementary slackness $\tilde{\zeta}_i \tilde{\theta}_i = 0$, one has $\tilde{\zeta}_i = 0$ and $\tilde{\gamma} = -J'_i(\tilde{\theta}_i) \geq 0$; $\forall j \in T$, similarly, one has $\tilde{\zeta}_j \geq 0$ and $\tilde{\gamma} \geq -J'_j(\tilde{\theta}_j) \geq 0$. Thus the KKT condition can be equivalently written as (12).

D Proof of Lemma 4

Proof of Lemma 4. By Lemma 2, it follows that $J'_t(0) \leq 0$. Combing with the condition $J'_s(0) < J'_t(0)$, we have $J'_s(0) < 0$ and thus $J'_s(\theta_s)$ belongs to (case b), (case c) or (case d). Denote

$$r_s = \begin{cases} 1, & J'_s(\theta_s) \text{ belongs to (case b);} \\ r_{s1}, & J'_s(\theta_s) \text{ belongs to (case c) or (case d),} \end{cases} \quad (\text{D.1})$$

where by the definition of r_{i1} in Lemma 2, $r_{s1} = \frac{|x_s|}{\sqrt{-2\eta J'_s(0)}}$.

Assume by contradiction that $\tilde{\theta}_s = 0$ yet $\tilde{\theta}_t > 0$. Let $\hat{\theta}$ be a vector of which the elements are equal to the elements of $\tilde{\theta}$ except that

$$\hat{\theta}_s = \min\{\tilde{\theta}_t, r_s\}; \quad (\text{D.2})$$

$$\hat{\theta}_t = \max\{0, \tilde{\theta}_s - r_s\}. \quad (\text{D.3})$$

By the definition of $\hat{\theta}_s, \hat{\theta}_t$ in (D.2) and (D.3), it follows that

$$\forall \theta_s \in [\tilde{\theta}_s, \hat{\theta}_s], \quad J'_s(\theta_s) = J'_s(0) \quad (\text{D.4})$$

$$\forall \theta_t \in [\hat{\theta}_t, \tilde{\theta}_t], \quad J'_t(\theta_t) \geq J'_t(0). \quad (\text{D.5})$$

Then

$$\begin{aligned} J(\tilde{\theta}) - J(\hat{\theta}) &= J_s(0) + \int_0^{\tilde{\theta}_s} J'_s(\theta_s) d\theta_s + J_t(0) + \int_0^{\tilde{\theta}_t} J'_t(\theta_t) d\theta_t \\ &\quad - J_s(0) - \int_0^{\hat{\theta}_s} J'_s(\theta_s) d\theta_s - J_t(0) - \int_0^{\hat{\theta}_t} J'_t(\theta_t) d\theta_t \\ &= \int_{\hat{\theta}_s}^{\tilde{\theta}_s} J'_s(\theta_s) d\theta_s + \int_{\hat{\theta}_t}^{\tilde{\theta}_t} J'_t(\theta_t) d\theta_t \\ &\geq \int_{\hat{\theta}_s}^{\tilde{\theta}_s} J'_s(0) d\theta_s + \int_{\hat{\theta}_t}^{\tilde{\theta}_t} J'_t(0) d\theta_t \\ &\geq J'_s(0)(\tilde{\theta}_s - \hat{\theta}_s) + J'_t(0)(\tilde{\theta}_t - \hat{\theta}_t) \end{aligned}$$

Then by the expressions of $\hat{\theta}_s, \hat{\theta}_t$ in (D.2) and (D.3),

$$J(\tilde{\theta}) - J(\hat{\theta}) = \begin{cases} (J'_t(0) - J'_s(0)) \cdot \tilde{\theta}_t, & \tilde{\theta}_t < r_s \\ (J'_t(0) - J'_s(0)) \cdot r_s, & \tilde{\theta}_t \geq r_s \end{cases}. \quad (\text{D.6})$$

By the assumption $J'_s(0) < J'_t(0)$, $J(\tilde{\theta}) - J(\hat{\theta}) > 0$, which contradicts the fact that $\tilde{\theta}$ is the optimal solution. \square

E Proof of Lemma 5

Proof of Lemma 5. By the KKT condition (12) in Lemma 3, it follows that for all $i \in S$, $-J'_i(\tilde{\theta}_i) \geq \max_{j \in T} -J'_j(0)$; meanwhile by Corollary 1, for all $i \in [d]$, $-J'_i(\tilde{\theta}_i)$ is a non-increasing function. Therefore combing the KKT condition (12), we have

$$\forall i \in S, \quad -J'_i(0) \geq -J'_i(\tilde{\theta}_i) \geq \max_{j \in T} \{-J'_j(0)\} \geq \max_{j \in T} \{-J'_j(1)\}. \quad (\text{E.1})$$

In addition, by the KKT condition (12), for all $i_1 \in S, i_2 \in S$, $-J'_{i_1}(\tilde{\theta}_{i_1}) = -J'_{i_2}(\tilde{\theta}_{i_2})$. Because by Corollary 1, for all $i \in [d]$, $-J'_i(\tilde{\theta}_i)$ is a non-increasing function, therefore

$$\forall i_1 \in S, i_2 \in S, \quad -J'_{i_1}(0) \geq -J'_{i_1}(\tilde{\theta}_{i_1}) = -J'_{i_2}(\tilde{\theta}_{i_2}) \geq -J'_{i_2}(1).$$

Therefore it follows that

$$\forall i \in S, \quad -J'_i(0) \geq \max_{j \in S} -J'_j(1). \quad (\text{E.2})$$

By combing (E.1) and (E.2), we get

$$-J'_i(0) \geq \max_{j \in [d]} -J'_j(1). \quad (\text{E.3})$$

\square

F Proof of Theorem 1

Proof. To prove Theorem 1, by Lemma 1, we only need to show $\tilde{\theta}$ in Alg. 1 is the optimal solution of the problem (7). By Lemma 3, to prove the optimality of $\tilde{\theta}$ in Alg. 1, we only need to show the $\tilde{\gamma}, \tilde{\theta}$ in Alg. 1 satisfy the KKT condition in Lemma 3. Equivalently, we rewrite the KKT condition as follows,

$$\begin{cases} \sum_{i \in [d]} \tilde{\theta}_i = 1, & (\text{F.1a}) \end{cases}$$

$$\begin{cases} \text{for all } i \in [d], \quad \tilde{\theta}_i \geq 0, & (\text{F.1b}) \end{cases}$$

$$\begin{cases} \text{for all } i \in S, \quad \tilde{\gamma} = -J'_i(\tilde{\theta}_i) \geq \max_{j \in T} -J'_j(0), & (\text{F.1c}) \end{cases}$$

where as in Lemma 3, $S = \{i \in [d] | \tilde{\theta}_i > 0\}$, $T = \{i \in [d] | \tilde{\theta}_i = 0\}$.

By checking the step 5 in Alg. 1, it is found that (F.1a) is already satisfied. Meanwhile, for $i \in R_\rho \setminus \{i_\rho\}$ and $i \in [d] \setminus R_\rho$, $\theta_i \geq 0$. So the remaining work is to show that the two conditions $\tilde{\theta}_{i_\rho} \geq 0$ and (F.1c) can be satisfied, which is given in Lemmas 6, 7 and 8.

Lemma 6. Assume $\tilde{\gamma}$ is defined in the step 4 of Alg. 1, then $\tilde{\gamma} \geq \max_{j \in [d]} -J'_j(0)$.

Lemma 7. As in the step 5 of Alg. 1, for all $l \in R_\rho \setminus \{i_\rho\}$, by setting $\tilde{\theta}_l = \frac{|x_l|}{\sqrt{2\eta\tilde{\gamma}}}$, it follows that $\tilde{\gamma} = -J'_l(\tilde{\theta}_l)$.

Lemma 8. As in the step 5 of Alg. 1, by setting $\tilde{\theta}_{i_\rho} = 1 - \sum_{k \in R_\rho \setminus \{i_\rho\}} \tilde{\theta}_k$, then it follows that $\tilde{\theta}_{i_\rho} \geq 0$ always hold. Meanwhile, if $\tilde{\theta}_{i_\rho} > 0$, then $\tilde{\gamma} = -J'_{i_\rho}(\tilde{\theta}_{i_\rho})$.

By Lemma 8, $\tilde{\theta}_{i_\rho} \geq 0$. To show (F.1c) can be satisfied, we give the following discussion.

- If $\tilde{\theta}_{i_\rho} = 0$, let $S = R_\rho \setminus \{i_\rho\}$ and $T = [d] \setminus (R_\rho \setminus \{i_\rho\})$, then by Lemmas 6, 7 and 8, (F.1c) is satisfied.
- If $\tilde{\theta}_{i_\rho} > 0$, let $S = R_\rho$ and $T = [d] \setminus R_\rho$, then by Lemmas 6, 7 and 8, (F.1c) is satisfied.

Therefore Theorem 1 is proved. □

G Technical Lemmas and Proofs of Lemmas 6, 7 and 8

The main difficulty in the proof of Lemmas 7 and 8 comes from the fact that by Lemma 2, for all $i \in [d]$ and $0 \leq \theta_i \leq 1$, the expression of $J'_i(\theta_i)$ has 4 different cases. Here we give Lemma 9 to show an equivalence relation between the expression of $J'_i(\theta_i)$ and the relation of $J'_i(0)$ and $J'_i(1)$.

Lemma 9. For all $i \in [d]$ and $0 \leq \theta_i \leq 1$, $J'_i(\theta_i)$ belongs to the (case a) or (case b) in Lemma 2 if and only if $J'_i(0) = J'_i(1)$; $J'_i(\theta_i)$ belongs to the (case c) or (case d) in Lemma 2 if and only if $J'_i(0) \neq J'_i(1)$.

Proof of Lemma 9. For all $i \in [d]$ and $0 \leq \theta_i \leq 1$, by observing the (case a), (case b), (case c) and (case d) of $J'_i(\theta_i)$ in Lemma 2, it follows that $J'_i(\theta_i)$ belongs to the (case a) or (case b) if and only if it is a constant function, which implies $J'_i(0) = J'_i(1)$. $J'_i(\theta_i)$ belongs to the (case c) or (case d) if and only if it is a piecewise function, which implies $J'_i(0) \neq J'_i(1)$. □

By Lemma 9, the condition $J'_{i_j}(0) = J'_{i_j}(1)$ in the step 3 of Alg. 1 is used to identify which case $J'_{i_j}(\theta_{i_j})$ belongs to. Lemma 10 introduces an implied result of the conditions in the step 3 of Alg. 1. (In the following lemmas, we assume that r_{i1}, r_{i2} ($i \in [d]$) have been defined in Lemma 2 and $i_m, v_m, Q, v, \rho, \tilde{\gamma}, i_j, R_j$ ($j \in [|Q| + 1]$) have been defined in Alg. 1.)

Lemma 10. For all $j \in [\rho - 1]$, it follows that all the following conditions

$$\begin{cases} J'_{i_j}(0) \neq J'_{i_j}(1), & \text{(G.1a)} \\ \sum_{l \in R_j} |x_l| < \sqrt{2\eta v_j}, & \text{(G.1b)} \\ j < |Q| + 1, & \text{(G.1c)} \end{cases}$$

must be satisfied.

Proof of Lemma 10. By the step 3 in Alg. 1, ρ is the minimal index that can satisfy one of the 3 following conditions

$$\begin{cases} J'_{i_j}(0) = J'_{i_j}(1), \\ \sum_{l \in R_\rho} |x_l| \geq \sqrt{2\eta v_\rho}, \\ \rho = |Q| + 1, \end{cases}$$

which implies that for all $j \in [\rho - 1]$, j satisfies all the 3 conditions in Lemma 10. \square

By Lemma 10, $j \in [\rho - 1]$ shares the 3 common properties in (G.1a)-(G.1c), which is important for the proof of the subsequent lemmas about $j \in [\rho - 1]$. In Lemma 11, we can find useful inequalities.

Lemma 11. For all $j \in [\rho - 1]$, $v_j \geq \tilde{\gamma} \geq v_\rho \geq v_m \geq -J'_{i_j}(1)$.

Proof of Lemma 11. By the step 4 in Alg. 1, $\tilde{\gamma}$ has two possible values.

If $\tilde{\gamma} = (\sum_{k \in R_{\rho-1}} |x_k|)^2 / (2\eta)$, it follows that

- In Lemma 10, let $j = \rho - 1$, we have $\sum_{k \in R_{\rho-1}} |x_k| < \sqrt{2\eta v_{\rho-1}}$. Then $v_{\rho-1} \geq \tilde{\gamma}$. Then by the definition of v , $v_1 \geq v_2 \geq \dots \geq v_{\rho-1}$. Thus for all $j \in [\rho - 1]$, we have $v_j \geq \tilde{\gamma}$.
- In the step 4, when $\tilde{\gamma} = (\sum_{k \in R_{\rho-1}} |x_k|)^2 / (2\eta)$, by the condition $\sum_{k \in R_{\rho-1}} |x_k| \geq \sqrt{2\eta v_\rho}$, we have $\tilde{\gamma} \geq v_\rho$.

If $\tilde{\gamma} = v_\rho$, by the definition of v , $v_1 \geq v_2 \geq \dots \geq v_\rho$. Then it follows that for all $j \in [\rho - 1]$, $v_j \geq \tilde{\gamma} \geq v_\rho$.

By the definition of v , $v_1 \geq v_2 \geq \dots \geq v_\rho \geq \dots \geq v_m$. By the definition of v_m , for all $j \in [\rho - 1]$, $v_m = \max_{i \in [d]} -J'_i(1) \geq -J'_{i_j}(1)$. Therefore $v_\rho \geq v_m \geq -J'_{i_j}(1)$.

Combining the above analysis, Lemma 11 is proved. \square

Lemma 11 gives $\tilde{\gamma}$ both lower and upper bounds, which then further bounds the range of $\tilde{\theta}_l = \frac{|x_l|}{\sqrt{2\eta\tilde{\gamma}}}$.

Before continue, we show the relation between $R_{\rho-1}$ and $R_\rho \setminus \{i_\rho\}$.

Lemma 12. If $\rho < |Q| + 1$, then $R_{\rho-1} = R_\rho \setminus \{i_\rho\}$; if $v_{|Q|+1} = -J'_{i_{|Q|+1}}(1)$, then $R_{|Q|} = R_{|Q|+1} \setminus \{i_{|Q|+1}\}$; if $v_{|Q|+1} \neq -J'_{i_{|Q|+1}}(1)$, then $R_{|Q|} = R_{|Q|+1}$.

Proof of Lemma 12.

- If $\rho < |Q| + 1$, then by the step 2 in Alg. 1, $i_1, i_2, \dots, i_{|Q|}$ are different coordinates. Thus $R_{\rho-1} = R_\rho \setminus \{i_\rho\}$.
- If $v_{|Q|+1} = -J'_{i_{|Q|+1}}(1)$, by the definition of Q in the step 1 of Alg. 1, $i_{|Q|+1} \notin Q$. Therefore $R_{|Q|} = R_{|Q|+1} \setminus \{i_{|Q|+1}\}$.
- If $v_{|Q|+1} \neq -J'_{i_{|Q|+1}}(1)$, then by Lemma 9, $J'_{i_{|Q|+1}}(\theta_{i_{|Q|+1}})$ belongs to (case c) or (case d). It follows that $-J'_{i_{|Q|+1}}(0) > -J'_{i_{|Q|+1}}(1) = v_m$. Then by the definition of $|Q|$ in the step 1 of Alg. 1, $i_{|Q|+1} \in Q$. Therefore $R_{|Q|} = R_{|Q|+1}$.

\square

Based on the above technical lemmas, we can prove Lemmas 6, 7 and 8.

G.1 Proofs of Lemma 6, 7 and 8

Proof of Lemma 6. By the step 4 of Alg. 1, if $\sum_{l \in R_{\rho-1}} |x_l| \geq \sqrt{2\eta v_\rho}$, then $\tilde{\gamma} = (\sum_{k \in R_{\rho-1}} |x_k|)^2 / (2\eta)$ which means $\tilde{\gamma} \geq v_\rho$; otherwise, $\tilde{\gamma} = v_\rho$.

In addition by the definition of v , it follows that $v_1 \geq v_2 \geq \dots \geq v_\rho \geq \dots \geq v_m = \max_{j \in [d]} -J'_j(0)$.

Therefore $\tilde{\gamma} \geq v_\rho \geq \max_{j \in [d]} -J'_j(0)$. Lemma 6 is proved. \square

Proof of Lemma 7. By Lemma 12, $R_\rho \setminus \{i_\rho\} \subset R_{\rho-1}$. Then by Lemma 10, for all $l \in R_\rho \setminus \{i_\rho\}$, $J'_l(0) \neq J'_l(1)$. Then by Lemma 9, for all $l \in R_\rho \setminus \{i_\rho\}$ and $0 \leq \theta_l \leq 1$, $J'_l(\theta_l)$ belongs to (case c) or (case d) in Lemma 2.

Assume that for $i \in [d]$, $r_{i1} = \frac{|x_i|}{\sqrt{-2\eta J'_i(0)}}$ and $r_{i2} = \frac{|x_i|}{\sqrt{-2\eta J'_i(1)}}$ have been defined in Lemma 2. By Lemma 11, we have $-J'_{i_j}(0) = v_j \geq \tilde{\gamma} \geq -J'_{i_j}(1)$. Then $r_{i1} \leq \tilde{\theta}_l = \frac{|x_l|}{\sqrt{2\eta\tilde{\gamma}}} \leq r_{i2}$.

In addition, if $\tilde{\gamma} = (\sum_{k \in R_{\rho-1}} |x_k|)^2 / (2\eta)$, when $\tilde{\theta}_l = \frac{|x_l|}{\sum_{k \in R_{\rho-1}} |x_k|} \leq 1$; if $\tilde{\gamma} = v_\rho$, by the step 4 in Alg. 1, the condition $\sum_{k \in R_{\rho-1}} |x_k| \leq \sqrt{2\eta v_\rho}$ holds. Then $\tilde{\theta}_l = \frac{|x_k|}{\sqrt{2\eta\tilde{\gamma}}} = \frac{|x_k|}{\sqrt{2\eta v_\rho}} \leq 1$.

Therefore $r_{i1} \leq \tilde{\theta}_l \leq \min\{r_{i2}, 1\}$. By the form of (case c) and (case d) in Lemma 2, we can find that for all $l \in R_\rho \setminus \{i_\rho\}$, $J'_l(\tilde{\theta}_l) = -\frac{|x_l|^2}{2\eta\tilde{\theta}_l^2} = -\tilde{\gamma}$.

□

To prove the KKT condition (F.1c), besides Lemma 7 for the case $j \in [\rho - 1]$, we also need Lemma 8 for the case $j = \rho$.

Proof of Lemma 8. By the step 4 in Alg. 1, $\tilde{\gamma}$ has two possible values.

If $\tilde{\gamma} = (\sum_{k \in R_{\rho-1}} |x_k|)^2 / (2\eta)$, by analyzing the following 3 cases, we can show that $\tilde{\theta}_{i_\rho} \geq 0$ always holds and if $\tilde{\theta}_{i_\rho} > 0$, then $J'_{i_\rho}(\tilde{\theta}_{i_\rho}) = -\tilde{\gamma}$.

- If $\rho < |Q| + 1$, then by Lemma 10, $R_\rho \setminus \{i_\rho\} = R_{\rho-1}$. By the step 5 in Alg. 1 and $\tilde{\gamma} = (\sum_{k \in R_{\rho-1}} |x_k|)^2 / (2\eta)$, for all $l \in R_\rho \setminus \{i_\rho\}$, $\tilde{\theta}_l = \frac{|x_l|}{\sqrt{2\eta\tilde{\gamma}}} = \frac{|x_l|}{\sum_{k \in R_{\rho-1}} |x_k|}$. Then $\sum_{l \in R_\rho \setminus \{i_\rho\}} \tilde{\theta}_l = \sum_{l \in R_{\rho-1}} \tilde{\theta}_l = 1$. Therefore $\tilde{\theta}_{i_\rho} = 1 - \sum_{k \in R_\rho \setminus \{i_\rho\}} \tilde{\theta}_k = 0$.
- If $\rho = |Q| + 1$ and $J'_{i_\rho}(0) = J'_{i_\rho}(1)$, then by Lemma 10, it follows that $R_\rho \setminus \{i_\rho\} = R_{\rho-1}$. Therefore, by the same analysis in the case $\rho < |Q| + 1$, $\tilde{\theta}_{i_\rho} = 1 - \sum_{k \in R_\rho \setminus \{i_\rho\}} \tilde{\theta}_k = 0$.
- If $\rho = |Q| + 1$ and $J'_{i_\rho}(0) \neq J'_{i_\rho}(1)$, then by Lemma 10, $R_{\rho-1} = R_\rho = (R_\rho \setminus \{i_\rho\}) \cup \{i_\rho\}$. Therefore $0 \leq \tilde{\theta}_{i_\rho} = 1 - \sum_{k \in R_\rho \setminus \{i_\rho\}} \tilde{\theta}_k = \frac{|x_{i_\rho}|}{\sum_{k \in R_{\rho-1}} |x_k|} \leq 1$. Meanwhile, by Lemma 9, $J'_{i_\rho}(\tilde{\theta}_{i_\rho})$ belongs to (case c) or (case d) in Lemma 2. Due to $\tilde{\gamma} = (\sum_{k \in R_{\rho-1}} |x_k|)^2 / (2\eta)$, then the condition $\sum_{k \in R_{\rho-1}} |x_k| \geq \sqrt{2\eta v_\rho}$ holds. Thus

$$\tilde{\theta}_{i_\rho} \leq \frac{|x_{i_\rho}|}{\sqrt{2\eta v_\rho}} = \frac{|x_{i_\rho}|}{\sqrt{2\eta v_{|Q|+1}}} = \frac{|x_{i_\rho}|}{\sqrt{-2\eta J'_{i_\rho}(1)}} = r_{i_\rho 2},$$

where $r_{i_\rho 2}$ is defined in Lemma 2. Meanwhile, in Lemma 10, let $j = \rho - 1$, then $\sum_{k \in R_{\rho-1}} |x_k| < \sqrt{2\eta v_{\rho-1}}$. Thus $\tilde{\theta}_{i_\rho} > \frac{|x_{i_\rho}|}{\sqrt{2\eta v_{\rho-1}}} = \frac{|x_{i_\rho}|}{\sqrt{2\eta v_{|Q|}}}$. In addition, due to $i_\rho \in Q$, we have $v_{i_\rho} \geq v_{|Q|}$. Thus

$$\tilde{\theta}_{i_\rho} > \frac{|x_{i_\rho}|}{\sqrt{2\eta v_{i_\rho}}} = \frac{|x_{i_\rho}|}{\sqrt{-2\eta J'_{i_\rho}(0)}} = r_{i_\rho 1},$$

where $r_{i_\rho 1}$ is defined in Lemma 2. Combing the above analyses, we have

$$0 < r_{i_\rho 1} < \tilde{\theta}_{i_\rho} \leq \min\{1, r_{i_\rho 2}\}.$$

Therefore by the form of (case c) and (case d) in Lemma 2, $J'_{i_\rho}(\tilde{\theta}_{i_\rho}) = -\frac{x_{i_\rho}^2}{2\eta\tilde{\theta}_{i_\rho}} = -\tilde{\gamma}$.

If $\tilde{\gamma} = v_\rho$, according to the condition in the step 4 of Alg. 1, $\sum_{l \in R_{\rho-1}} |x_l| < \sqrt{2\eta v_\rho}$. Meanwhile by Lemma 12, we have $R_{i_\rho} \setminus \{i_\rho\} \subset R_{i_{\rho-1}}$. Then $\sum_{l \in R_{i_\rho} \setminus \{i_\rho\}} \tilde{\theta}_l = \sum_{l \in R_{i_\rho} \setminus \{i_\rho\}} \frac{|x_l|}{\sqrt{2\eta\tilde{\gamma}}} \leq \sum_{l \in R_{i_{\rho-1}}} \frac{|x_l|}{\sqrt{2\eta\tilde{\gamma}}} < 1$. Therefore $\tilde{\theta}_{i_\rho} = 1 - \sum_{l \in R_{i_\rho} \setminus \{i_\rho\}} \tilde{\theta}_l > 0$. We can give the analyses by discussing the following 3 cases.

- If $v_\rho = -J'_{i_\rho}(1)$, then by Lemma 9, $J_{i_\rho}(\theta_{i_\rho})$ belongs to (case a) or (case b). Therefore $J'_{i_\rho}(\tilde{\theta}_{i_\rho}) = J'_{i_\rho}(0) = -v_\rho = -\tilde{\gamma}$.
- If $v_\rho \neq -J'_{i_\rho}(1)$ and $\sum_{l \in R_\rho} |x_l| \geq \sqrt{2\eta v_\rho}$, then by Lemma 9, $J'_{i_\rho}(\theta_{i_\rho})$ belongs to (case c) or (case d) in Lemma 2. Meanwhile for $l \in R_\rho \setminus \{i_\rho\}$, by $\tilde{\theta}_l = \frac{|x_l|}{\sqrt{2\eta\tilde{\gamma}}} = \frac{|x_l|}{\sqrt{2\eta v_\rho}}$, we have

$$\tilde{\theta}_{i_\rho} = 1 - \sum_{l \in R_\rho \setminus \{i_\rho\}} \tilde{\theta}_l = 1 - \sum_{l \in R_\rho \setminus \{i_\rho\}} \frac{|x_l|}{\sqrt{2\eta v_\rho}} \leq \frac{|x_{i_\rho}|}{\sqrt{2\eta v_\rho}} = r_{i_\rho 1},$$

where $r_{i_\rho 1}$ is defined in Lemma 2 and the last inequality is due to the condition $\sum_{l \in R_\rho} |x_l| \geq \sqrt{2\eta v_\rho}$. Therefore by the form of (case c) and (case d) in Lemma 2, it follows that $J'_{i_\rho}(\tilde{\theta}_{i_\rho}) = J'_{i_\rho}(0) = -v_\rho = -\tilde{\gamma}$.

- If $v_\rho \neq -J'_{i_\rho}(1)$, $\sum_{l \in R_\rho} |x_l| < \sqrt{2\eta v_\rho}$ and $\rho = |Q| + 1$, then by $v_\rho \neq -J'_{i_\rho}(1)$ and Lemma 9, $J'_{i_\rho}(\theta_{i_\rho})$ belongs to (case c) or (case d) in Lemma 2. By $\sum_{l \in R_\rho} |x_l| < \sqrt{2\eta v_\rho}$, we have that the $r_{i_\rho 2}$ in Lemma 2 satisfies

$$r_{i_\rho 2} = \frac{|x_{i_\rho}|}{\sqrt{-2\eta J'_{i_\rho}(1)}} = \frac{|x_{i_\rho}|}{\sqrt{2\eta v_{|Q|+1}}} = \frac{|x_{i_\rho}|}{\sqrt{2\eta v_\rho}} < \frac{\sum_{l \in R_\rho} |x_l|}{\sqrt{2\eta v_\rho}} < 1,$$

therefore $J'_{i_\rho}(\theta_{i_\rho})$ belongs to the (case d) in Lemma 2. For all $l \in R_\rho \setminus \{i_\rho\}$, by $\tilde{\theta}_l = \frac{|x_l|}{\sqrt{2\eta\tilde{\gamma}}} = \frac{|x_l|}{\sqrt{2\eta v_\rho}}$ and $\sum_{l \in R_\rho} |x_l| < \sqrt{2\eta v_\rho}$, we have

$$\tilde{\theta}_{i_\rho} = 1 - \sum_{k \in R_\rho \setminus \{i_\rho\}} \tilde{\theta}_k \geq \frac{|x_{i_\rho}|}{\sqrt{2\eta v_\rho}} = r_{i_\rho 2}. \quad (\text{G.3})$$

Therefore, by the form of the (case d) in Lemma 2, $J'_{i_\rho}(\tilde{\theta}_{i_\rho}) = J'_{i_\rho}(1) = -v_\rho = -\tilde{\gamma}$.

Summarizing the above analyses, we prove Lemma 8. □

H Proof of Theorem 2

H.1 Some necessary Lemmas and Definitions

For $1 < p < \infty$ and the ℓ_p -norm $\|\cdot\|_p$, we denote its dual norm as $\|x\|_q = \max_{\|y\|_p \leq 1} x^T y = (\sum_{i=1}^d |x_i|^q)^{\frac{1}{q}}$, where $\frac{1}{p} + \frac{1}{q} = 1$. For $p = 1$, by the definition of dual norm, the dual norm of ℓ_1 -norm is ℓ_∞ -norm. In Lemma 13 and 14, some classical results are described.

Lemma 13. (I8, §3.1.9) *If $\forall x, y \in \mathbb{R}^d$, $1 \leq p \leq \infty$, $\frac{1}{p} + \frac{1}{q} = 1$, and $\eta > 0$, then $|\langle x, y \rangle| \leq \|x\|_q \|y\|_p \leq \frac{1}{2\eta} \|x\|_q^2 + \frac{\eta}{2} \|y\|_p^2$.*

Lemma 14. *If $\forall x \in \mathbb{R}^d$, $1 \leq p \leq \infty$, then $\|x\|_p \leq \|x\|_1 \leq n^{1-\frac{1}{p}} \|x\|_p$.*

For a continuous differentiable function $f(x)$, we give the following definitions.

Definition 1. *$f(x)$ is L_p -smooth ($1 \leq p \leq \infty$) w.r.t. $\|\cdot\|_p$ if $\forall x, y \in \mathbb{R}^d$ and $\frac{1}{p} + \frac{1}{q} = 1$, $\|\nabla f(x) - \nabla f(y)\|_q \leq L_p \|x - y\|_p$.*

By Definition 1, we have Lemma 15.

Lemma 15. *If $f(x)$ is L_p -smooth ($1 \leq p \leq \infty$) w.r.t. $\|\cdot\|_p$ and $\frac{1}{p} + \frac{1}{q} = 1$, then*

$$\frac{1}{2L_p} \|\nabla f(x) - \nabla f(y)\|_q^2 \leq f(y) - f(x) - \langle \nabla f(x), y - x \rangle \leq \frac{L_p}{2} \|x - y\|_p^2. \quad (\text{H.1})$$

Proof. Firstly it is showed that $f(x)$ being L_p -smooth w.r.t. $\|\cdot\|_p$ implies that $\forall x, y \in \mathbb{R}^d$

$$f(y) \leq f(x) + \langle \nabla f(x), y - x \rangle + \frac{L_p}{2} \|y - x\|_p^2.$$

Consider the function $g(\tau) = f(x + \tau(y - x))$ with $\tau \in \mathbb{R}$. Then

$$\begin{aligned} f(y) - f(x) - \langle \nabla f(x), y - x \rangle &= g(1) - g(0) - \langle \nabla f(x), y - x \rangle \\ &= \int_0^1 \left(\frac{dg(\tau)}{d\tau} - \langle \nabla f(x), y - x \rangle \right) d\tau \\ &= \int_0^1 (\langle \nabla f(x + \tau(y - x)), y - x \rangle - \langle \nabla f(x), y - x \rangle) d\tau \\ &= \int_0^1 \langle \nabla f(x + \tau(y - x)) - \nabla f(x), y - x \rangle d\tau \\ &\leq \int_0^1 \|\nabla f(x + \tau(y - x)) - \nabla f(x)\|_q \|y - x\|_p d\tau \\ &\leq \int_0^1 L_p \tau \|y - x\|_p^2 d\tau \\ &= \frac{L_p}{2} \tau^2 \|y - x\|_p^2 \Big|_0^1 \\ &= \frac{L_p}{2} \|y - x\|_p^2. \end{aligned}$$

To subsequently show $\frac{1}{2L_p} \|\nabla f(x) - \nabla f(y)\|_q^2 \leq f(y) - f(x) - \langle \nabla f(x), y - x \rangle$, fix $x \in \mathbb{R}^d$ and consider the function

$$\phi(y) = f(y) - \langle \nabla f(x), y \rangle,$$

which is convex on \mathbb{R}^d and also has an L_p -Lipschitz continuous gradient w.r.t. $\|\cdot\|_p$, as

$$\begin{aligned} \|\phi'(y) - \phi'(x)\|_q &= \|(\nabla f(y) - \nabla f(x)) - (\nabla f(x) - \nabla f(x))\|_q \\ &= \|\nabla f(y) - \nabla f(x)\|_q \\ &\leq L_p \|y - x\|_p. \end{aligned}$$

As the minimizer of ϕ is x (i.e., $\phi'(x) = 0$), for any $y \in \mathbb{R}^d$, we have

$$\begin{aligned} \phi(x) &= \min_v \phi(v) \leq \min_v \left\{ \phi(y) + \langle \phi'(y), v - y \rangle + \frac{L_p}{2} \|v - y\|_p^2 \right\} \\ &= \phi(y) - \sup_v \left\{ \langle -\phi'(y), v - y \rangle - \frac{L_p}{2} \|v - y\|_p^2 \right\} \\ &= \phi(y) - \frac{1}{2L_p} \|\phi'(y)\|_q^2. \end{aligned}$$

Substituting in the definition of ϕ , we have

$$\begin{aligned} f(x) - \langle \nabla f(x), x \rangle &\leq f(y) - \langle \nabla f(x), y \rangle - \frac{1}{2L_p} \|\nabla f(y) - \nabla f(x)\|_q^2 \\ \iff \frac{1}{2L_p} \|\nabla f(y) - \nabla f(x)\|_q^2 &\leq f(y) - f(x) - \langle \nabla f(x), y - x \rangle. \end{aligned}$$

□

Definition 2. $f(x)$ is σ_p -strongly convex ($1 \leq p \leq \infty$) w.r.t. $\|\cdot\|_p$ if $\forall x, y \in \mathbb{R}^d$ and $\frac{1}{p} + \frac{1}{q} = 1$, $f(y) - f(x) - \langle \nabla f(x), y - x \rangle \geq \frac{\sigma_p}{2} \|x - y\|_p^2$.

Taking $\frac{1}{2}\|x\|_p^2$ ($1 < p \leq 2$) as an example. It is known that $\frac{1}{2}\|x\|_p^2$ is $(p-1)$ -strongly convex w.r.t. $\|\cdot\|_p$ [6]. Based on $\frac{1}{2}\|x\|_p^2$ ($1 < p \leq 2$), one can define p -Bregman divergence

$$B_p(y, x) = \frac{1}{2}\|y\|_p^2 - \frac{1}{2}\|x\|_p^2 - \langle \nabla \frac{1}{2}\|x\|_p^2, y - x \rangle. \quad (\text{H.2})$$

Lemma 16 ([6, 2]). For $x, y \in \mathbb{R}^d$, $1 < p \leq 2$, $B_p(y, x) = \frac{1}{2}\|y\|_p^2 - \frac{1}{2}\|x\|_p^2 - \langle \nabla \frac{1}{2}\|x\|_p^2, y - x \rangle$ satisfies the 3 properties.

- $B_p(y, x) \geq \frac{p-1}{2}\|y - x\|_p^2$;
- $B_p(y, x) = 0$ if and only if $y = x$;
- $B_p(x, y) + B_p(y, z) = B_p(x, z) + \langle \frac{1}{2}\nabla\|z\|_p^2 - \frac{1}{2}\nabla\|y\|_p^2, x - y \rangle$.

H.2 Proof of Theorem 2

Theorem 2 is proved by following the steps of the proof in [1]. First, in Section H.2.1, ASGCD is analyzed for the fixed k -th iteration. In the one-iteration analysis, y_k, z_k and x_{k+1} are assumed to be fixed and thus the selection of the mini batch \mathcal{B} in the k -th iteration is the only source of randomness. For simplicity, let $\tilde{x} = \tilde{x}^s$, $\tau_1 = \tau_{1,s}$, $\alpha = \alpha_s$ where $s = \lfloor \frac{k}{m} \rfloor$ is the epoch corresponding to k . Let $\beta(b) \stackrel{\text{def}}{=} \frac{n-b}{b(n-1)}$ and denote $\sigma_{k+1}^2 \stackrel{\text{def}}{=} \|\nabla f(x_{k+1}) - \tilde{\nabla}_{k+1}\|_\infty^2$. Then $\mathbb{E}[\sigma_{k+1}^2]$ is the variance measured by $\|\cdot\|_\infty$ of the gradient estimator $\tilde{\nabla}_{k+1}$ in this iteration. Second, in Section H.2.2, Theorem 2 is proven by combing the one-iteration analysis in Section H.2.1 into the outer-iteration analysis in Section H.2.2.

There are 3 differences from the analysis in [1]. First, the analysis is used for the specific ASGCD algorithm that combines SOTOPO and pCOMID and thus the value of the parameter α_s is different from the setting in [1]. Second, the analysis is given under the mini batch selection setting and $\|\cdot\|_\infty$ rather than one sample selection setting and $\|\cdot\|_2$. Third, we use a different way to represent the convergence result for ℓ_1 -regularized problems (1).

H.2.1 One-iteration analysis

Lemma 17 (SOTOPO). If

$$\begin{aligned} y_{k+1} &= \text{SOTOPO}(\tilde{\nabla}_{k+1}, x_{k+1}, \lambda, \eta), \quad \text{and} \\ \text{Prog}(x_{k+1}) &= -\min_{y \in \mathbb{R}^d} \left\{ \frac{(1+2\beta(b))L_1}{2} \|y - x_{k+1}\|_1^2 + \langle \tilde{\nabla}_{k+1}, y - x_{k+1} \rangle + \lambda \|y\|_1 - \lambda \|x_{k+1}\|_1 \right\} \geq 0, \end{aligned}$$

it follows that if $b < n$ (where the expectation is only over the randomness of $\tilde{\nabla}_{k+1}$), then

$$F(x_{k+1}) - \mathbb{E}[F(y_{k+1})] \geq \mathbb{E}[\text{Prog}(x_{k+1})] - \frac{1}{4\beta(b)L_1} \mathbb{E}[\sigma_{k+1}^2]; \quad (\text{H.3})$$

if $b = n$ (no randomness exists), then

$$F(x_{k+1}) - \mathbb{E}[F(y_{k+1})] \geq \mathbb{E}[\text{Prog}(x_{k+1})]. \quad (\text{H.4})$$

Proof. If $b < n$, it follows that

$$\begin{aligned}
\text{Prog}(x_{k+1}) &= - \min_{y \in \mathbb{R}^d} \left\{ \frac{(1 + 2\beta(b)) L_1}{2} \|y - x_{k+1}\|_1^2 + \langle \tilde{\nabla}_{k+1}, y - x_{k+1} \rangle + \lambda \|y\|_1 - \lambda \|x_{k+1}\|_1 \right\} \\
&\stackrel{\textcircled{1}}{=} - \left(\frac{(1 + 2\beta(b)) L_1}{2} \|y_{k+1} - x_{k+1}\|_1^2 + \langle \tilde{\nabla}_{k+1}, y_{k+1} - x_{k+1} \rangle + \lambda \|y_{k+1}\|_1 - \lambda \|x_{k+1}\|_1 \right) \\
&= - \left(\frac{L_1}{2} \|y_{k+1} - x_{k+1}\|_1^2 + \langle \nabla f(x_{k+1}), y_{k+1} - x_{k+1} \rangle + \lambda \|y_{k+1}\|_1 - \lambda \|x_{k+1}\|_1 \right) \\
&\quad + \left(\langle \nabla f(x_{k+1}) - \tilde{\nabla}_{k+1}, y_{k+1} - x_{k+1} \rangle - \beta(b) L_1 \|y_{k+1} - x_{k+1}\|_1^2 \right) \\
&\stackrel{\textcircled{2}}{\leq} - (f(y_{k+1}) - f(x_{k+1})) + \lambda \|y_{k+1}\|_1 - \lambda \|x_{k+1}\|_1 + \frac{1}{4\beta(b)L_1} \|\nabla f(x_{k+1}) - \tilde{\nabla}_{k+1}\|_\infty^2,
\end{aligned}$$

where $\textcircled{1}$ is by Theorem 1, $\textcircled{2}$ is by the smoothness assumption (2), Lemma 13 and 15. Taking expectation on both sides, (H.3) is obtained.

If $b = n$, then $\beta(b) = 0$. By using a similar analysis as the case $b < n$, (H.4) is obtained. \square

Lemma 18. (variance upper bound). If $b < n$, then

$$\mathbb{E}[\|\tilde{\nabla}_{k+1} - \nabla f(x_{k+1})\|_\infty^2] \leq 2\beta(b)L_1(f(\tilde{x}) - f(x_{k+1}) - \langle \nabla f(x_{k+1}), \tilde{x} - x_{k+1} \rangle). \quad (\text{H.5})$$

Proof. Before the proof, it should be noted that the variance upper bound measured by $\|\cdot\|_2$ of mini-batch selection has been proved in [27]. The variance in our case is measured by $\|\cdot\|_\infty$. Because some properties of $\|\cdot\|_2$ such as $\mathbb{E}[\|x - \mathbb{E}[x]\|_2^2] = \mathbb{E}[\|x\|_2^2] - \|\mathbb{E}[x]\|_2^2$ and $\|\sum_i x_i\|_2^2 = \sum_{i,j} x_i^T x_j$ can't be generalized to $\|\cdot\|_\infty$ directly, the proof is slightly different from the proof in [27].

Let $\phi_j = (\nabla f_j(x_{k+1}) - \nabla f_j(\tilde{x})) - (\nabla f(x_{k+1}) - \nabla f(\tilde{x}))$ and $\phi_j^i = (\nabla_i f_j(x_{k+1}) - \nabla_i f_j(\tilde{x})) - (\nabla_i f(x_{k+1}) - \nabla_i f(\tilde{x}))$. Denote $i_{\max} = \arg \max_i |\nabla_i f(x_{k+1}) - \nabla_{k+1,i}|$. It follows that

$$\begin{aligned}
\mathbb{E} \left[\left\| \frac{1}{b} \sum_{j \in \mathcal{B}} \phi_j^{i_{\max}} \right\|_\infty^2 \right] &= \frac{1}{b^2} \mathbb{E} \left[\sum_{j_1, j_2 \in \mathcal{B}} \phi_{j_1}^{i_{\max}} \phi_{j_2}^{i_{\max}} \right] \\
&= \frac{1}{b^2} \mathbb{E} \left[\sum_{j_1 \neq j_2 \in \mathcal{B}} \phi_{j_1}^{i_{\max}} \phi_{j_2}^{i_{\max}} \right] + \frac{1}{b} \mathbb{E} [(\phi_j^{i_{\max}})^2] \\
&= \frac{b-1}{bn(n-1)} \sum_{j_1 \neq j_2 \in [n]} \phi_{j_1}^{i_{\max}} \phi_{j_2}^{i_{\max}} + \frac{1}{b} \mathbb{E} [(\phi_j^{i_{\max}})^2] \\
&= \frac{b-1}{bn(n-1)} \sum_{j_1, j_2 \in [n]} \phi_{j_1}^{i_{\max}} \phi_{j_2}^{i_{\max}} - \frac{b-1}{b(n-1)} \mathbb{E} [(\phi_j^{i_{\max}})^2] + \frac{1}{b} \mathbb{E} [(\phi_j^{i_{\max}})^2] \\
&= \frac{b-1}{bn(n-1)} \sum_{j_1, j_2 \in [n]} \phi_{j_1}^{i_{\max}} \phi_{j_2}^{i_{\max}} - \beta(b) \mathbb{E} [(\phi_j^{i_{\max}})^2] \\
&\stackrel{\textcircled{1}}{=} \beta(b) \mathbb{E} [(\phi_j^{i_{\max}})^2] \\
&\stackrel{\textcircled{2}}{\leq} \beta(b) \mathbb{E} [\|\phi_j\|_\infty^2], \tag{H.6}
\end{aligned}$$

where ① is using the fact $\sum_{j=1}^n \phi_j^{i_{\max}} = 0$, ② is by definition of $\|\cdot\|_\infty$. Denote $i_{j_{\max}} = \arg \max_i |\phi_j^i|$. Hence

$$\begin{aligned}
& \mathbb{E} \left[\left\| \nabla f(x_{k+1}) - \tilde{\nabla}_{k+1} \right\|_\infty^2 \right] \\
&= \mathbb{E} \left[\left\| \frac{1}{b} \sum_{j \in \mathcal{B}} (\nabla f_j(x_{k+1}) - \nabla f_j(\tilde{x})) - (\nabla f(x_{k+1}) - \nabla f(\tilde{x})) \right\|_\infty^2 \right] \\
&\stackrel{\text{①}}{\leq} \beta(b) \mathbb{E} \left[\left\| \nabla f_j(x_{k+1}) - \nabla f_j(\tilde{x}) - (\nabla f(x_{k+1}) - \nabla f(\tilde{x})) \right\|_\infty^2 \right] \\
&\stackrel{\text{②}}{=} \beta(b) \mathbb{E} \left[(\nabla_{i_{j_{\max}}} f_j(x_{k+1}) - \nabla_{i_{j_{\max}}} f_j(\tilde{x})) - \nabla_{i_{j_{\max}}} f(x_{k+1}) + \nabla_{i_{j_{\max}}} f(\tilde{x}) \right]^2 \\
&= \beta(b) \mathbb{E} \left[(\nabla_{i_{j_{\max}}} f_j(x_{k+1}) - \nabla_{i_{j_{\max}}} f_j(\tilde{x}))^2 - (\nabla_{i_{j_{\max}}} f(x_{k+1}) - \nabla_{i_{j_{\max}}} f(\tilde{x}))^2 \right] \\
&\leq \beta(b) \mathbb{E} \left[(\nabla_{i_{j_{\max}}} f_j(x_{k+1}) - \nabla_{i_{j_{\max}}} f_j(\tilde{x}))^2 \right] \\
&\stackrel{\text{③}}{\leq} \beta(b) \mathbb{E} \left[\left\| \nabla f_j(x_{k+1}) - \nabla f_j(\tilde{x}) \right\|_\infty^2 \right] \\
&\leq 2\beta(b)L_1 \mathbb{E} [f_j(\tilde{x}) - f_j(x_{k+1}) - \langle \nabla f_j(x_{k+1}), \tilde{x} - x_{k+1} \rangle] \\
&\stackrel{\text{④}}{=} 2\beta(b)L_1 (f(\tilde{x}) - f(x_{k+1}) - \langle \nabla f(x_{k+1}), \tilde{x} - x_{k+1} \rangle),
\end{aligned}$$

where ① is by (H.6), ② is using the fact $\mathbb{E}[(x - \mathbb{E}[x])^2] = \mathbb{E}[x^2] - (\mathbb{E}[x])^2$, ④ is by the definition of $\|\cdot\|_\infty$, ③ is by Lemma 15. □

Lemma 19 (pCOMID). *Fixing $\tilde{\nabla}_{k+1}$ and letting*

$$(z_{k+1}, \theta_{k+1}) = pCOMID(\tilde{\nabla}_{k+1}, \theta_k, q, \lambda, \alpha), \quad (\text{H.7})$$

it satisfies for all $u \in \mathbb{R}^d$,

$$\alpha \langle \tilde{\nabla}_{k+1}, z_{k+1} - u \rangle + \alpha \lambda \|z_{k+1}\|_1 - \alpha \lambda \|u\|_1 \leq -B_p(z_{k+1}, z_k) + B_p(u, z_k) - B_p(u, z_{k+1}). \quad (\text{H.8})$$

Proof. From [13], we know that pCOMID exactly solves the following mirror descent problem,

$$z_{k+1} = \arg \min_z \{ \langle \tilde{\nabla}_{k+1}, z - z_k \rangle + \frac{1}{\alpha} B_p(z, z_k) + \lambda \|z\|_1 \}. \quad (\text{H.9})$$

By the optimality condition of z_{k+1} , it follows that

$$\nabla \frac{1}{2} \|z_{k+1}\|_p^2 - \nabla \frac{1}{2} \|z_k\|_p^2 + \alpha \tilde{\nabla}_{k+1} + \alpha g = 0,$$

where $g \in \partial \lambda \|z_{k+1}\|_1$. Then the equality

$$\langle \nabla \frac{1}{2} \|z_{k+1}\|_p^2 - \nabla \frac{1}{2} \|z_k\|_p^2 + \alpha \tilde{\nabla}_{k+1} + \alpha g, z_{k+1} - u \rangle = 0$$

holds. In addition by Lemma 16, it follows that $\langle \nabla \frac{1}{2} \|z_{k+1}\|_p^2 - \nabla \frac{1}{2} \|z_k\|_p^2, z_{k+1} - u \rangle = B_p(z_{k+1}, z_k) - B_p(u, z_k) + B_p(u, z_{k+1})$. By the convexity of $\lambda \|z\|_1$, $\langle g, z_{k+1} - u \rangle \geq \lambda \|z_{k+1}\|_1 - \lambda \|u\|_1$. Therefore, we can write

$$\begin{aligned}
& \alpha \langle \tilde{\nabla}_{k+1}, z_{k+1} - u \rangle + \alpha \lambda \|z_{k+1}\|_1 - \alpha \lambda \|u\|_1 \\
&= -\langle \nabla \frac{1}{2} \|z_{k+1}\|_p^2 - \nabla \frac{1}{2} \|z_k\|_p^2, z_{k+1} - u \rangle - \langle \alpha g, z_{k+1} - u \rangle + \alpha \lambda \|z_{k+1}\|_1 - \alpha \lambda \|u\|_1 \\
&\leq -B_p(z_{k+1}, z_k) + B_p(u, z_k) - B_p(u, z_{k+1}).
\end{aligned}$$

□

Lemma 20 (Coupling step 1). *If $x_{k+1} = \tau_1 z_k + \tau_2 \tilde{x} + (1 - \tau_1 - \tau_2)y_k$, where $\tau_1 \leq \frac{1}{(1+2\beta(b))\alpha L_1}$ and $\tau_2 = \frac{1}{2}$,*

$$\begin{aligned} & \alpha \langle \nabla f(x_{k+1}), z_k - u \rangle - \alpha \lambda \|u\|_1 \\ \leq & \frac{\alpha}{\tau_1} (F(x_{k+1}) - \mathbb{E}[F(y_{k+1})] + \tau_2 F(\tilde{x}) - \tau_2 \mathbb{E}[F(x_{k+1})] - \tau_2 \langle \nabla f(x_{k+1}), \tilde{x} - x_{k+1} \rangle) \\ & + B_p(u, z_k) - \mathbb{E}[B_p(u, z_{k+1})] + \frac{\alpha(1 - \tau_1 - \tau_2)}{\tau_1} \lambda \|y_k\|_1 - \frac{\alpha}{\tau_1} \lambda \|x_{k+1}\|_1. \end{aligned}$$

Proof. It follows that

$$\begin{aligned} & \alpha \langle \tilde{\nabla}_{k+1}, z_k - u \rangle + \alpha \lambda \|z_{k+1}\|_1 - \alpha \lambda \|u\|_1 \\ = & \alpha \langle \tilde{\nabla}_{k+1}, z_k - z_{k+1} \rangle + \alpha \langle \tilde{\nabla}_{k+1}, z_{k+1} - u \rangle + \alpha \lambda \|z_{k+1}\|_1 - \alpha \lambda \|u\|_1 \\ \stackrel{\textcircled{1}}{\leq} & \alpha \langle \tilde{\nabla}_{k+1}, z_k - z_{k+1} \rangle - B_p(z_{k+1}, z_k) + B_p(u, z_k) - B_p(u, z_{k+1}) \\ \stackrel{\textcircled{2}}{\leq} & \alpha \langle \tilde{\nabla}_{k+1}, z_k - z_{k+1} \rangle - \frac{p-1}{2} \|z_{k+1} - z_k\|_p^2 + B_p(u, z_k) - B_p(u, z_{k+1}) \\ \stackrel{\textcircled{3}}{\leq} & \alpha \langle \tilde{\nabla}_{k+1}, z_k - z_{k+1} \rangle - \frac{p-1}{2} d^{-(1-\frac{1}{p})} \|z_{k+1} - z_k\|_1^2 + B_p(u, z_k) - B_p(u, z_{k+1}), \\ \stackrel{\textcircled{4}}{=} & \alpha \langle \tilde{\nabla}_{k+1}, z_k - z_{k+1} \rangle - \frac{1}{2C} \|z_{k+1} - z_k\|_1^2 + B_p(u, z_k) - B_p(u, z_{k+1}), \end{aligned} \quad (\text{H.10})$$

where $\textcircled{1}$ is by Lemma 19, $\textcircled{2}$ is by Lemma 16, $\textcircled{3}$ is by Lemma 14 and $\textcircled{4}$ is by the setting $C = \frac{d^{\frac{2\delta}{1+\delta}}}{\delta}$ and $p = 1 + \delta$ in Alg. 2.

By defining $v \stackrel{\text{def}}{=} \tau_1 z_{k+1} + \tau_2 \tilde{x} + (1 - \tau_1 - \tau_2)y_k$, we have $x_{k+1} - v = \tau_1(z_k - z_{k+1})$ and therefore

$$\begin{aligned} & \mathbb{E}[\alpha \langle \tilde{\nabla}_{k+1}, z_k - z_{k+1} \rangle - \frac{1}{2C} \|z_{k+1} - z_k\|_1^2] = \mathbb{E}[\frac{\alpha}{\tau_1} \langle \tilde{\nabla}_{k+1}, x_{k+1} - v \rangle - \frac{1}{2C\tau_1^2} \|x_{k+1} - v\|_1^2] \\ = & \mathbb{E} \left[\frac{\alpha}{\tau_1} \left(\langle \tilde{\nabla}_{k+1}, x_{k+1} - v \rangle - \frac{1}{2C\alpha\tau_1} \|x_{k+1} - v\|_1^2 - \lambda \|v\|_1 + \lambda \|x_{k+1}\|_1 \right) \right. \\ & \left. + \frac{\alpha}{\tau_1} (\lambda \|v\|_1 - \lambda \|x_{k+1}\|_1) \right] \\ \stackrel{\textcircled{1}}{=} & \mathbb{E} \left[\frac{\alpha}{\tau_1} \left(\langle \tilde{\nabla}_{k+1}, x_{k+1} - v \rangle - \frac{(1+2\beta(b))L_1}{2} \|x_{k+1} - v\|_1^2 - \lambda \|v\|_1 + \lambda \|x_{k+1}\|_1 \right) \right. \\ & \left. + \frac{\alpha}{\tau_1} (\lambda \|v\|_1 - \lambda \|x_{k+1}\|_1) \right] \\ \stackrel{\textcircled{2}}{\leq} & \mathbb{E} \left[\frac{\alpha}{\tau_1} (F(x_{k+1}) - F(y_{k+1})) + \frac{1}{4\beta(b)L_1} \sigma_{k+1}^2 + \frac{\alpha}{\tau_1} (\lambda \|v\|_1 - \lambda \|x_{k+1}\|_1) \right] \\ \stackrel{\textcircled{3}}{\leq} & \mathbb{E} \left[\frac{\alpha}{\tau_1} \left(F(x_{k+1}) - F(y_{k+1}) + \frac{1}{2} (f(\tilde{x}) - f(x_{k+1}) - \langle \nabla f(x_{k+1}), \tilde{x} - x_{k+1} \rangle) \right) \right. \\ & \left. + \frac{\alpha}{\tau_1} (\tau_1 \lambda \|z_{k+1}\|_1 + \tau_2 \lambda \|\tilde{x}\|_1 + (1 - \tau_1 - \tau_2) \lambda \|y_k\|_1 - \lambda \|x_{k+1}\|_1) \right], \end{aligned} \quad (\text{H.11})$$

where $\textcircled{1}$ is by the setting $\alpha\tau_1 = \frac{1}{(1+2\beta(b))CL_1}$, $\textcircled{2}$ is by Lemma 17, $\textcircled{3}$ is by Lemma 18 and the convexity $\|v\|_1 = \|\tau_1 z_{k+1} + \tau_2 \tilde{x} + (1 - \tau_1 - \tau_2)y_k\|_1 \leq \tau_1 \|z_{k+1}\|_1 + \tau_2 \|\tilde{x}\|_1 + (1 - \tau_1 - \tau_2) \|y_k\|_1$. Then, it is showed that $\mathbb{E}[\langle \tilde{\nabla}_{k+1}, z_k - u \rangle] = \langle \nabla f(x_{k+1}), z_k - u \rangle$ and $\tau_2 = \frac{1}{2}$. Combing (H.10) and (H.11), Lemma 20 is obtained. \square

Lemma 21 (Coupling step 2). *Under the same choices of τ_1, τ_2 as in Lemma 20, one has*

$$\begin{aligned} 0 \leq & \frac{\alpha(1 - \tau_1 - \tau_2)}{\tau_1} (F(y_k) - F(x^*)) - \frac{\alpha}{\tau_1} (\mathbb{E}[F(y_{k+1})] - F(x^*)) + \frac{\alpha\tau_2}{\tau_1} (F(\tilde{x}) - F(x^*)) \\ & + B_p(x^*, z_k) - \mathbb{E}[B_p(x^*, z_{k+1})]. \end{aligned}$$

Proof. It follows that

$$\begin{aligned}
& \alpha(f(x_{k+1}) - f(u)) \stackrel{\textcircled{1}}{\leq} \alpha \langle \nabla f(x_{k+1}), x_{k+1} - u \rangle \\
& = \alpha \langle \nabla f(x_{k+1}), x_{k+1} - z_k \rangle + \alpha \langle \nabla f(x_{k+1}), z_k - u \rangle \\
& \stackrel{\textcircled{2}}{=} \frac{\alpha\tau_2}{\tau_1} \langle \nabla f(x_{k+1}), \tilde{x} - x_{k+1} \rangle + \frac{\alpha(1 - \tau_1 - \tau_2)}{\tau_1} \langle \nabla f(x_{k+1}), y_k - x_{k+1} \rangle + \alpha \langle \nabla f(x_{k+1}), z_k - u \rangle \\
& \stackrel{\textcircled{3}}{\leq} \frac{\alpha\tau_2}{\tau_1} \langle \nabla f(x_{k+1}), \tilde{x} - x_{k+1} \rangle + \frac{\alpha(1 - \tau_1 - \tau_2)}{\tau_1} (f(y_k) - f(x_{k+1})) + \alpha \langle f(x_{k+1}), z_k - u \rangle, \text{(H.12)}
\end{aligned}$$

where $\textcircled{1}$ is by the convexity of $f(x)$, $\textcircled{2}$ is by the convex combination $x_{k+1} = \tau_1 z_k + \tau_2 \tilde{x} + (1 - \tau_1 - \tau_2)y_k$, $\textcircled{3}$ is again by the convexity of $f(x)$. Applying Lemma 20 to (H.12), it follows that

$$\begin{aligned}
& \alpha(f(x_{k+1}) - F(u)) \leq \frac{\alpha(1 - \tau_1 - \tau_2)}{\tau_1} (F(y_k) - f(x_{k+1})) \\
& + \frac{\alpha}{\tau_1} (F(x_{k+1}) - \mathbb{E}[F(y_{k+1})]) + \tau_2 F(\tilde{x}) - \tau_2 f(x_{k+1}) + B_p(u, z_k) - \mathbb{E}[B_p(u, z_{k+1})] - \frac{\alpha}{\tau_1} \lambda \|x_{k+1}\|_1,
\end{aligned}$$

which implies

$$\begin{aligned}
& \alpha(F(x_{k+1}) - F(u)) \leq \frac{\alpha(1 - \tau_1 - \tau_2)}{\tau_1} (F(y_k) - F(x_{k+1})) \\
& + \frac{\alpha}{\tau_1} (F(x_{k+1}) - \mathbb{E}[F(y_{k+1})]) + \tau_2 F(\tilde{x}) - \tau_2 F(x_{k+1}) + B_p(u, z_k) - \mathbb{E}[B_p(u, z_{k+1})].
\end{aligned}$$

After arrangement and setting u to some minimizer x^* , Lemma 21 is obtained. \square

H.2.2 Proof of Theorem 2

Proof. Assume the parameter $\tau_{1,s}$ and α_s satisfies the assumption $\tau_{1,s}\alpha_s = \frac{1}{(1+2\beta(b))CL_1}$ in Lemma 20. Let $D_k \stackrel{\text{def}}{=} F(y_k) - F(x^*)$ and $\tilde{D}^s \stackrel{\text{def}}{=} F(\tilde{x}^s) - F(x^*)$, Lemma 21 can be rewritten as

$$0 \leq \frac{\alpha_s(1 - \tau_{1,s} - \tau_2)}{\tau_{1,s}} D_k - \frac{\alpha_s}{\tau_{1,s}} \mathbb{E}[D_{k+1}] + \frac{\alpha_s\tau_2}{\tau_{1,s}} \tilde{D}^s + B_p(x^*, z_k) - \mathbb{E}[B_p(x^*, z_{k+1})].$$

In the s -th epoch, summing up the above inequality for all the iterations $k = sm, sm + 1, \dots, sm + m - 1$, it follows that

$$\begin{aligned}
& \mathbb{E} \left[\alpha_s \frac{1 - \tau_{1,s} - \tau_2}{\tau_{1,s}} D_{(s+1)m} + \alpha_s \frac{\tau_{1,s} + \tau_2}{\tau_{1,s}} \sum_{l=1}^m D_{sm+l} \right] \\
& \leq \alpha_s \frac{1 - \tau_{1,s} - \tau_2}{\tau_{1,s}} D_{sm} + \alpha_s \frac{\tau_2}{\tau_{1,s}} m \tilde{D}^s + B_p(x^*, z_{sm}) - \mathbb{E}[B_p(x^*, z_{(s+1)m})]. \text{(H.13)}
\end{aligned}$$

It should be noted that in (H.13), we fix all the randomness in the first $s - 1$ epochs and take expectation on the current epoch s .

By the definition $\tilde{x}^s = \frac{1}{m} \sum_{l=1}^m y_{(s-1)m+l}$ in Alg. 2 and the convexity of $F(x)$, we have $m\tilde{D}^s \leq \sum_{l=1}^m D_{(s-1)m+l}$. Then for each $s \geq 1$, by (H.13), it follows that

$$\begin{aligned}
& \mathbb{E} \left[\frac{1}{\tau_{1,s}^2} D_{(s+1)m} + \frac{\tau_{1,s} + \tau_2}{\tau_{1,s}^2} \sum_{l=1}^{m-1} D_{sm+l} \right] \\
& \leq \frac{1 - \tau_{1,s}}{\tau_{1,s}^2} D_{sm} + \frac{\tau_2}{\tau_{1,s}^2} \sum_{l=1}^{m-1} D_{(s-1)m+l} + (1 + 2\beta(b)) CL B_p(x^*, z_{sm}) - (1 + 2\beta(b)) CL \mathbb{E}[B_p(x^*, z_{(s+1)m})].
\end{aligned}$$

For $s = 0$, (H.13) can be written as

$$\begin{aligned}
& \mathbb{E} \left[\frac{1}{\tau_{1,0}^2} D_m + \frac{\tau_{1,0} + \tau_2}{\tau_{1,0}^2} \sum_{l=1}^{m-1} D_l \right] \\
& \leq \frac{1 - \tau_{1,0} - \tau_2}{\tau_{1,0}^2} D_0 + \frac{\tau_2 m}{\tau_{1,0}^2} \tilde{D}^0 + (1 + 2\beta(b)) CL_1 B_p(x^*, z_0) \\
& \quad - (1 + 2\beta(b)) CL_1 \mathbb{E}[B_p(x^*, z_m)]. \text{(H.14)}
\end{aligned}$$

Choose $\tau_{1,s} = \frac{2}{s+4} \leq \frac{1}{2}$ which satisfies

$$\frac{1}{\tau_{1,s}^2} \geq \frac{1 - \tau_{1,s+1}}{\tau_{1,s+1}^2} \text{ and } \frac{\tau_{1,s} + \tau_2}{\tau_{1,s}^2} \geq \frac{\tau_2}{\tau_{1,s+1}^2}. \quad (\text{H.15})$$

Then it follows that

$$\begin{aligned} & \mathbb{E} \left[\frac{m}{\tau_{1,S-1}^2} \tilde{D}^S + (1 + 2\beta(b)) CL_1 B_p(x^*, z_{S_m}) \right] \\ \stackrel{\textcircled{1}}{\leq} & \mathbb{E} \left[\frac{1}{\tau_{1,S-1}^2} D_{S_m} + \frac{\tau_{1,S-1}}{\tau_{1,S-1}^2} \sum_{l=1}^{m-1} D_{(S-1)m+l} + (1 + 2\beta(b)) CL_1 B_p(x^*, z_{S_m}) \right] \\ \stackrel{\textcircled{2}}{\leq} & \mathbb{E} \left[\frac{1}{\tau_{1,S-1}^2} D_{S_m} + \frac{\tau_2}{\tau_{1,S}^2} \sum_{l=1}^{m-1} D_{(S-1)m+l} + (1 + 2\beta(b)) CL_1 B_p(x^*, z_{S_m}) \right] \\ \stackrel{\textcircled{3}}{\leq} & \frac{1 - \tau_{1,0} - \tau_2}{\tau_{1,0}^2} D_0 + \frac{\tau_2 m}{\tau_{1,0}^2} \tilde{D}^0 + (1 + 2\beta(b)) CL_1 B_p(x^*, z_0) \\ \stackrel{\textcircled{4}}{=} & \frac{\tau_2 m}{\tau_{1,0}^2} \tilde{D}^0 + (1 + 2\beta(b)) CL_1 B_p(x^*, z_0), \end{aligned}$$

where $\textcircled{1}$ is by $m\tilde{D}^s \leq \sum_{l=1}^m D_{(s-1)m+l}$, $\textcircled{2}$ is by $\tau_2 \geq \tau_{1,S-1} \geq \tau_{1,S}$, $\textcircled{3}$ uses (H.15) to telescope (H.13) and (H.14) for all $s = 0, 1, \dots, S-1$ and $\textcircled{4}$ is by $\tau_{1,0} = \tau_2 = \frac{1}{2}$.

$$\begin{aligned} & \mathbb{E}[F(\tilde{x}^S) - F(x^*)] \\ &= \mathbb{E}[\tilde{D}^S] \leq \frac{\tau_{1,S-1}^2}{m} \cdot \left(\frac{\tau_2 m}{\tau_{1,0}^2} \tilde{D}^0 + (1 + 2\beta(b)) CL_1 B_p(x^*, z_0) \right) \\ &= \frac{4}{m(S+3)^2} \left(2m(F(\tilde{x}^0) - F(x^*)) + (1 + 2\beta(b)) CL_1 B_p(x^*, z_0) \right). \quad (\text{H.16}) \end{aligned}$$

Meanwhile, by setting $\tilde{x}_0 = z_0 = 0$, using the optimality condition $0 \in \nabla f(x^*) + \partial\lambda\|x^*\|_1$ we have

$$\begin{aligned} F(\tilde{x}^0) - F(x^*) &= f(0) - f(x^*) - \lambda\|x^*\|_1 \\ &\stackrel{\textcircled{1}}{\leq} \langle \nabla f(x^*), 0 - x^* \rangle + \frac{L_1}{2} \|x^*\|_1^2 - \lambda\|x^*\|_1 \\ &\stackrel{\textcircled{2}}{=} \langle -\partial\lambda\|x^*\|_1, -x^* \rangle + \frac{L_1}{2} \|x^*\|_1^2 - \lambda\|x^*\|_1 \\ &\stackrel{\textcircled{3}}{\leq} \|\partial\lambda\|x^*\|_1\|_\infty \|x^*\|_1 + \frac{L_1}{2} \|x^*\|_1^2 - \lambda\|x^*\|_1 \\ &\stackrel{\textcircled{4}}{\leq} \lambda\|x^*\|_1 + \frac{L_1}{2} \|x^*\|_1^2 - \lambda\|x^*\|_1 \\ &= \frac{L_1}{2} \|x^*\|_1^2, \quad (\text{H.17}) \end{aligned}$$

where $\textcircled{1}$ is by the smoothness assumption of $f(x)$, $\textcircled{2}$ is by selecting the subgradient of $\lambda\|x^*\|_1$ with $-\partial\lambda\|x^*\|_1 = \nabla f(x^*)$, $\textcircled{3}$ is by lemma 13, $\textcircled{4}$ is by using the property of subgradient $\|\partial\lambda\|x^*\|_1\|_\infty \leq \lambda$. In addition, for $1 < p \leq 2$,

$$\begin{aligned} B_p(x^*, z_0) &= B_p(x^*, 0) = \frac{1}{2} \|x^*\|_p^2 - \frac{1}{2} \|0\|_p^2 - \langle \nabla \frac{1}{2} \|0\|_p^2, x^* - 0 \rangle \\ &= \frac{1}{2} \|x^*\|_p^2 \leq \frac{1}{2} \|x^*\|_1^2. \quad (\text{H.18}) \end{aligned}$$

Furthermore, minimizing $C = \frac{d^{2\delta}}{\delta}$ w.r.t δ , we get $\delta = \log(d) - 1 - \sqrt{(\log(d) - 1)^2 - 1}$ and $p = 1 + \delta = \log(d) - \sqrt{(\log(d) - 1)^2 - 1} \in (1, 2]$,

Then combing (H.16), (H.17) and (H.18), we get the final result.

$$\mathbb{E}[F(\tilde{x}^S) - F(x^*)] \leq \frac{4}{(S+3)^2} \left(1 + \frac{1 + 2\beta(b)}{2m} C \right) L_1 \|x\|_1^2. \quad (\text{H.19})$$

□